

Distributed Route Aggregation on the GlObal Network (DRAGON)

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Last year in the news (August 2014)

THE WALL STREET JOURNAL.  TECH

TECHNOLOGY

Echoes of Y2K: Engineers Buzz That Internet Is Outgrowing Its Gear

Routers That Send Data Online Could Become Overloaded as Number of Internet Routes Hits '512K'

By DREW FITZGERALD 

Updated Aug. 13, 2014 7:38 p.m. ET



14 August 2014 Last updated at 12:05 GMT

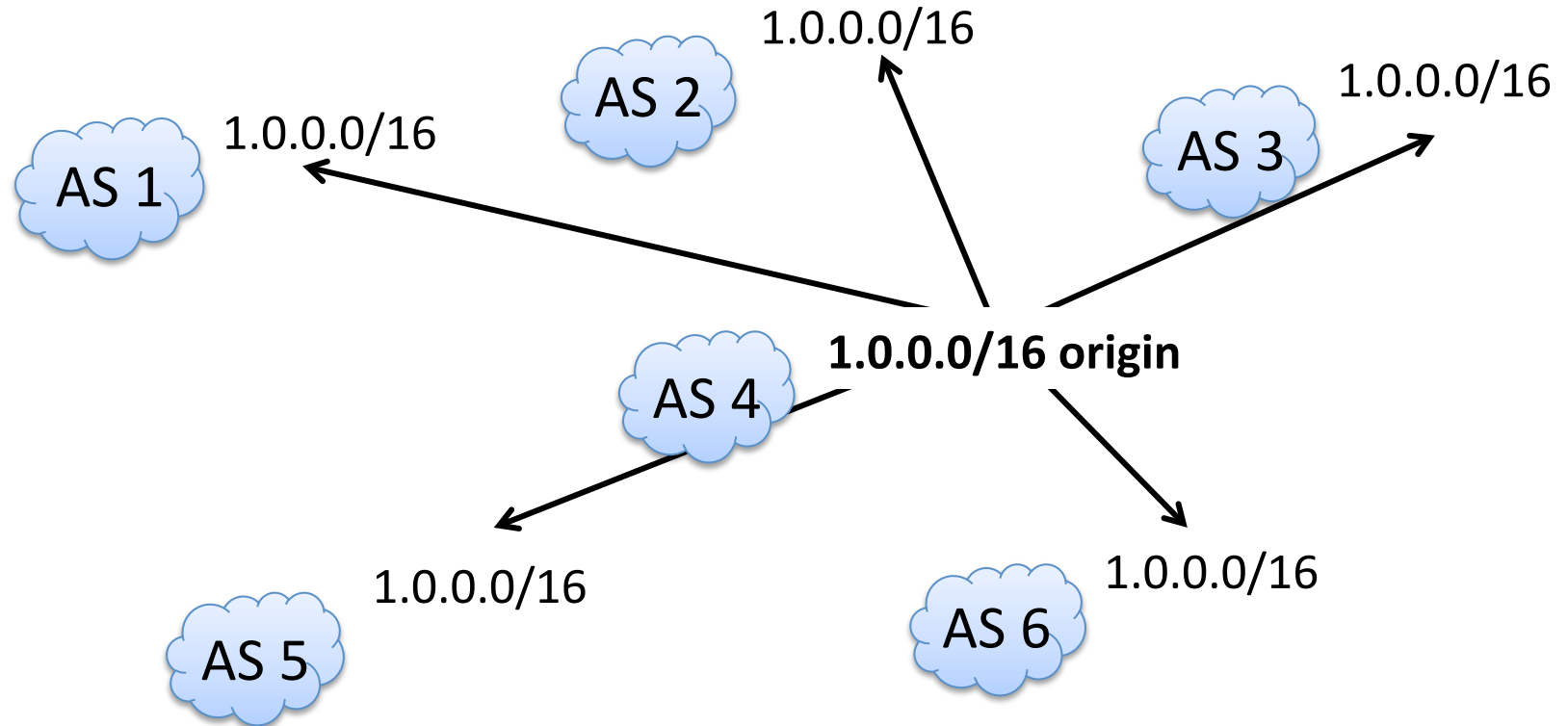
Browsing speeds may slow as net hardware bug bites

By Mark Ward

Technology correspondent, BBC News

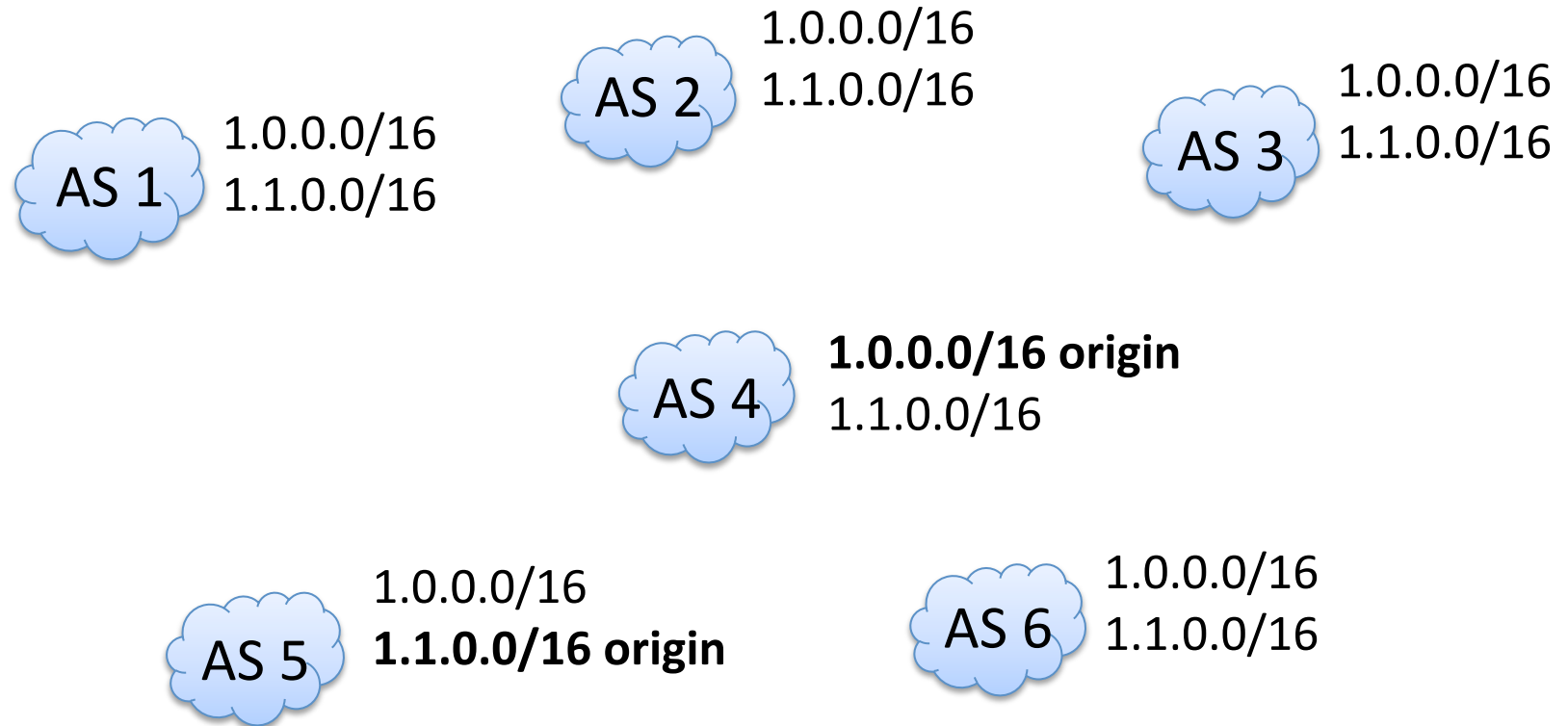
Some routers could not process the +512 K IPv4 prefixes they were learning about

Not a scalable routing system



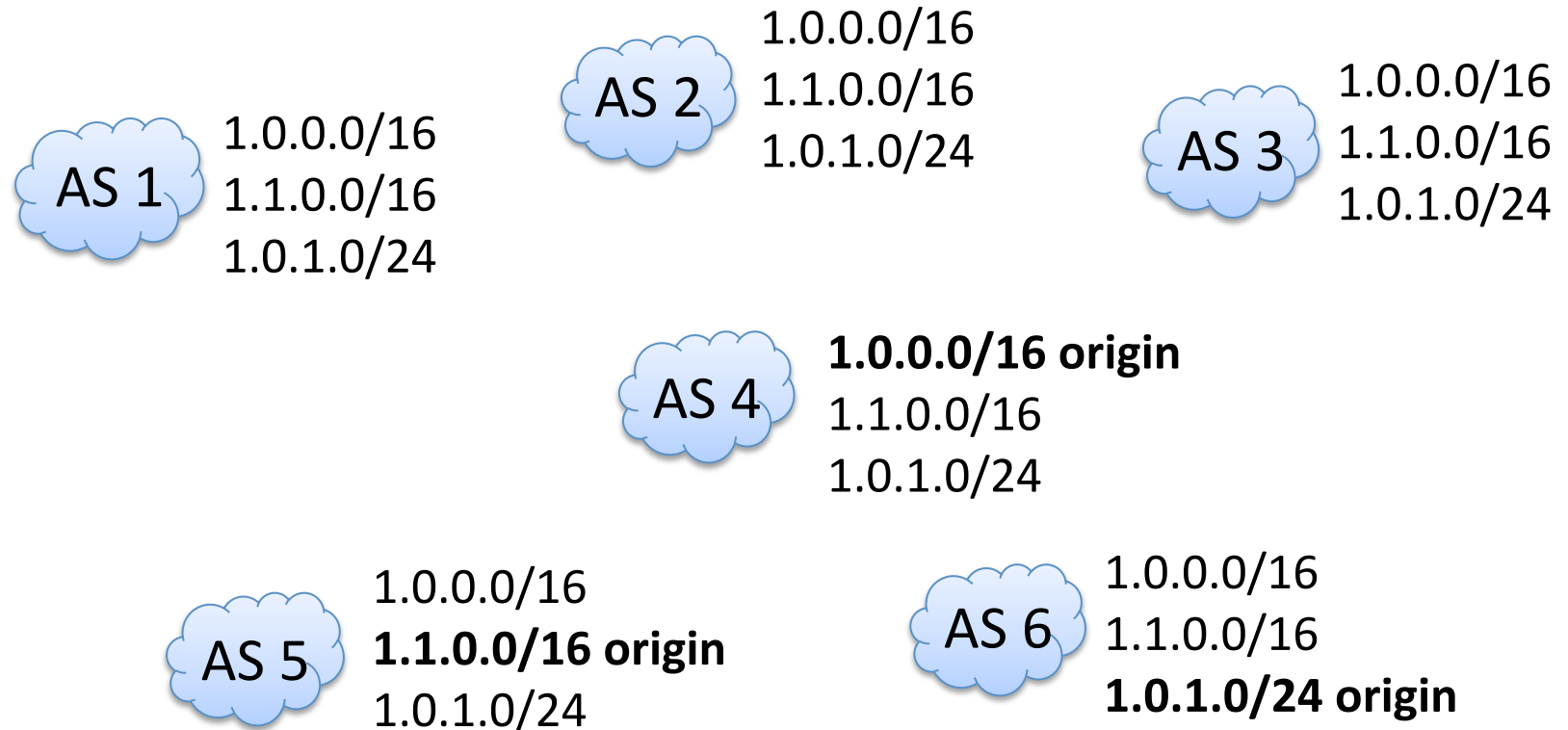
**Most of the originated prefixes are
routed globally (by BGP)**

Not a scalable routing system



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routed globally (by BGP)**

No scalability: poor performance

- Forwarding tables (FIBs) growth & address look-up time increase
- Routing tables (RIBs) growth
- BGP session set-up time increase
- Churn & convergence time increase

Further scalability concerns

- IPv6 prefixes can be formed in potentially larger numbers than IPv4 prefixes
- Secure BGP adds computational overhead to routing processes

DRAGON

Distributed solution to scale the Internet routing system

Basic DRAGON: 49% savings on routing state

Full DRAGON: 79% savings on routing state

No changes to the BGP protocol

No changes to the forwarding plane

Readily implemented with updated router software

Outline

- Scalability: global view
- DRAGON: filtering strategy
- DRAGON: aggregation strategy
- DRAGON: performance evaluation
- Conclusions

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Scalability: global view (routing)

AS 1 1.0.0.0/16

Specificity

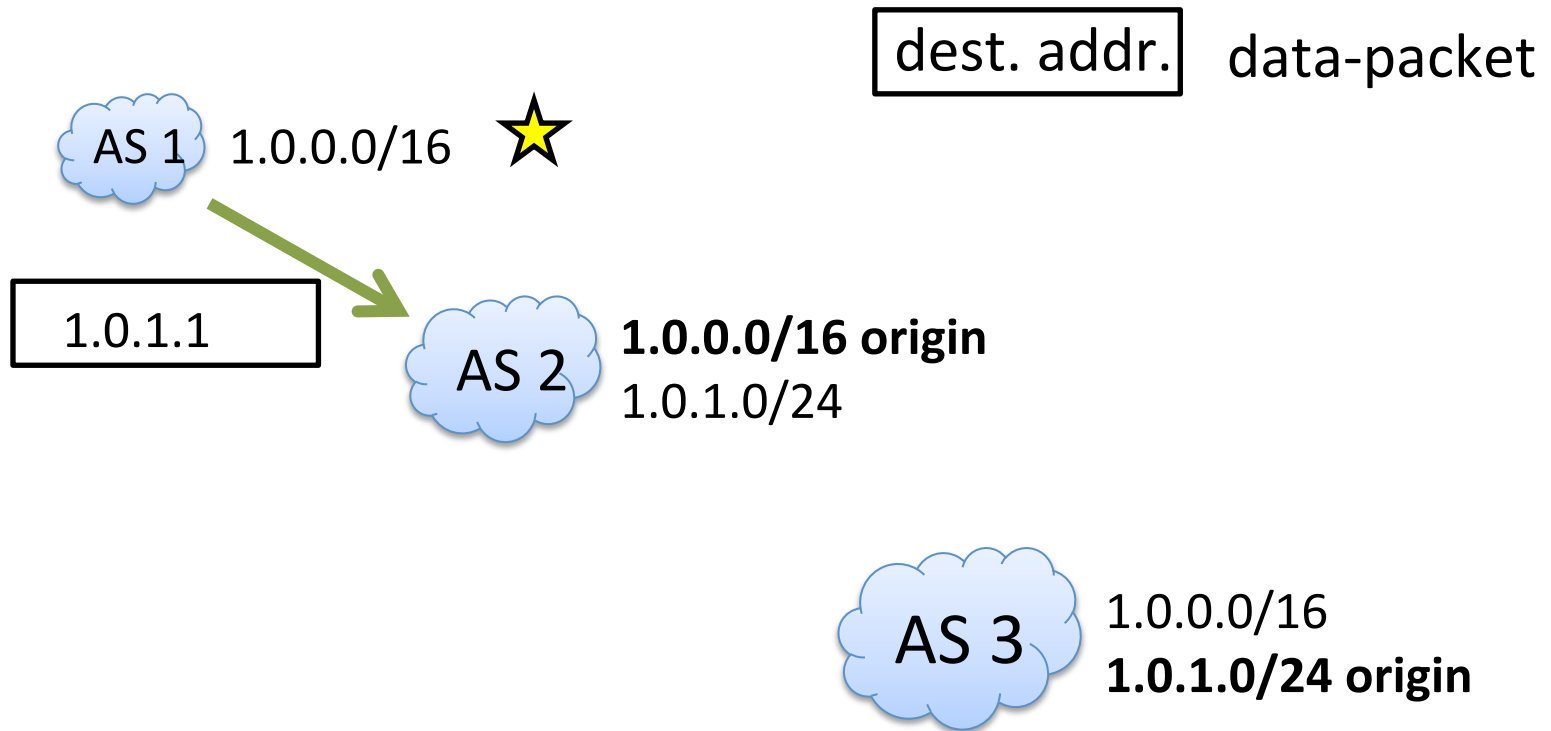
Prefix q is more specific than prefix p if the address space of q is contained in that of p

AS 2 1.0.0.0/16 origin
1.0.1.0/24

AS 3 1.0.0.0/16
1.0.1.0/24 origin

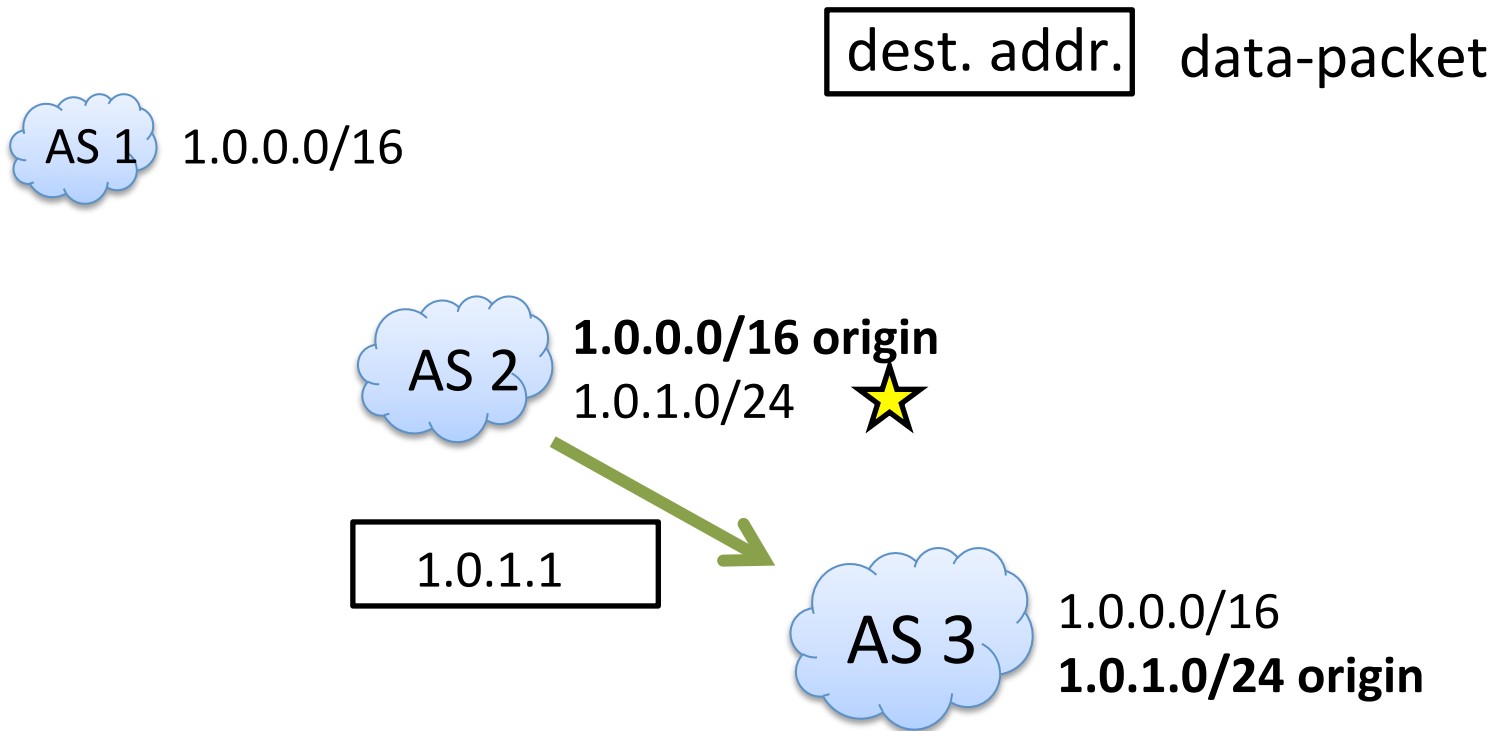
Propagation of more specific prefixes only in a small vicinity of their origin ASs

Scalability: global view (forwarding)

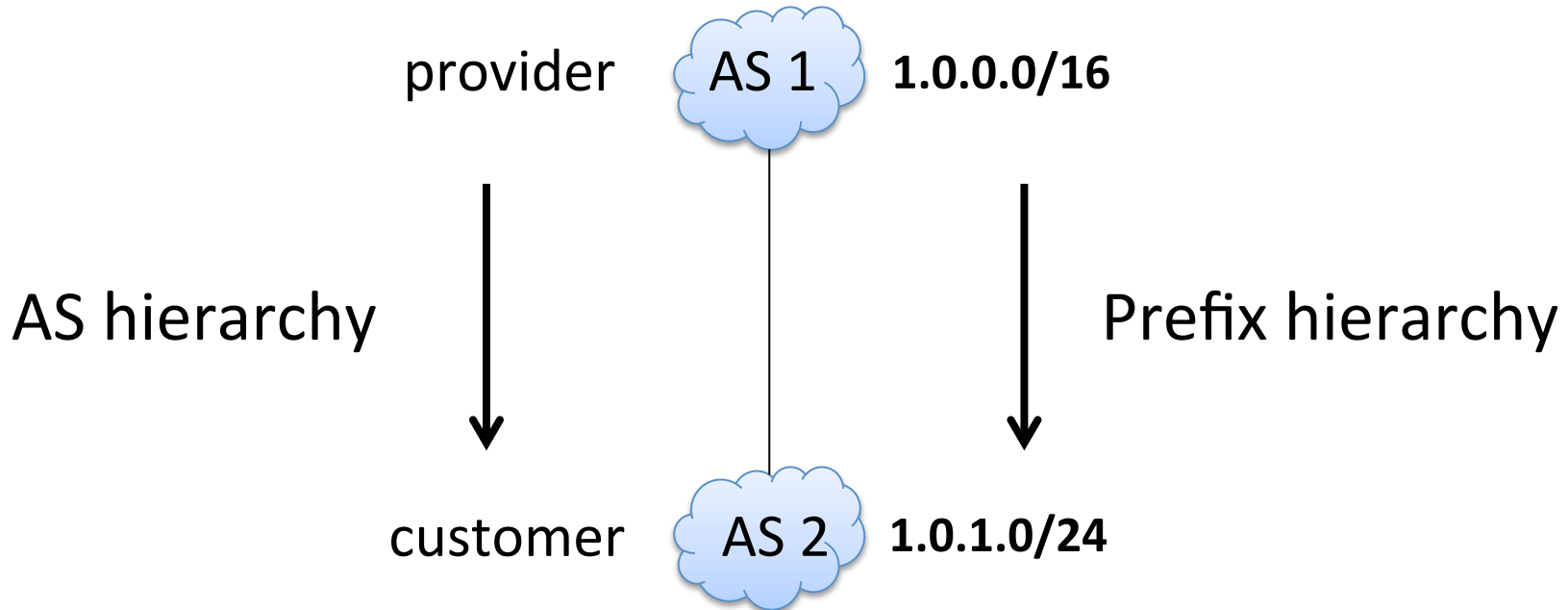


**Most ASs forward data-packets on the
(aggregated) less specific prefixes**

Scalability: global view (forwarding)



Hope for scalability? Hierarchies



AS-hierarchy aligned with prefix hierarchy

Hope for scalability? Clustering

Routing Information Registry (RIR)



$$1.0.0.0/24 + 1.0.1.0/24 + 1.0.2.0/23 = 1.0.0.0/22$$

Geography *roughly* clusters together ASs with aggregatable address space

Challenge: global vs. local

How to realize the global view through automated local routing decisions?

especially, given that the Internet routing system is as decentralized as it can be:

- each AS decides where to connect
- each AS decides where to acquire address space
- each AS sets its own routing policies

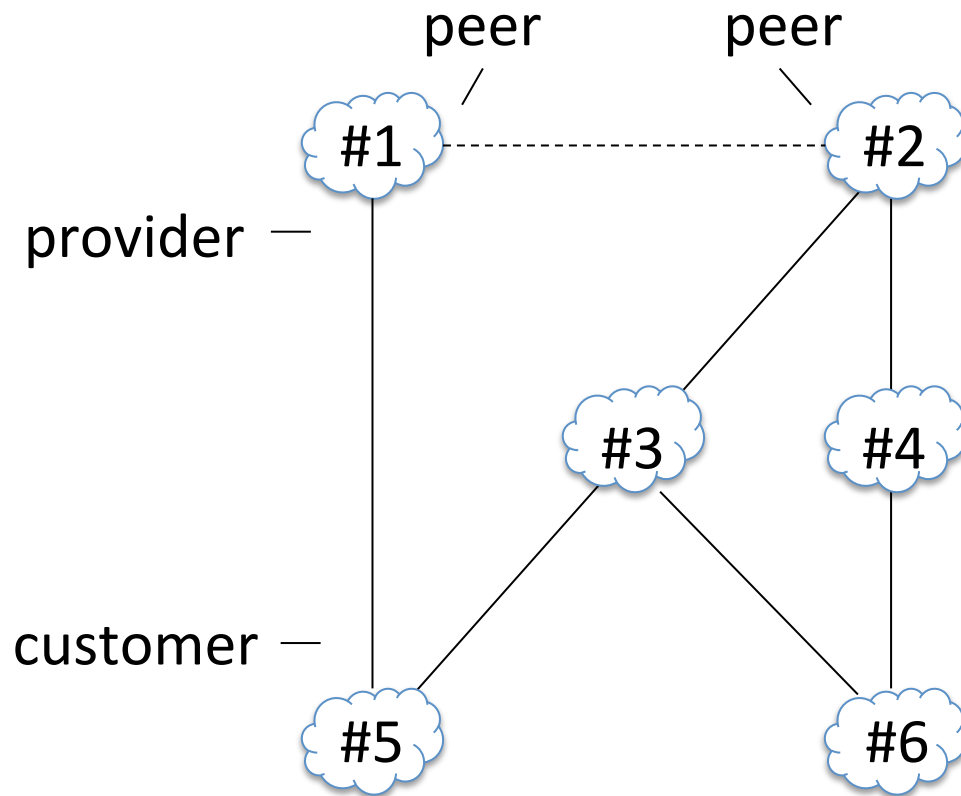
Outline

- Scalability: global view
- **DRAGON: filtering strategy**
- DRAGON: aggregation strategy
- DRAGON: performance evaluation
- Conclusions

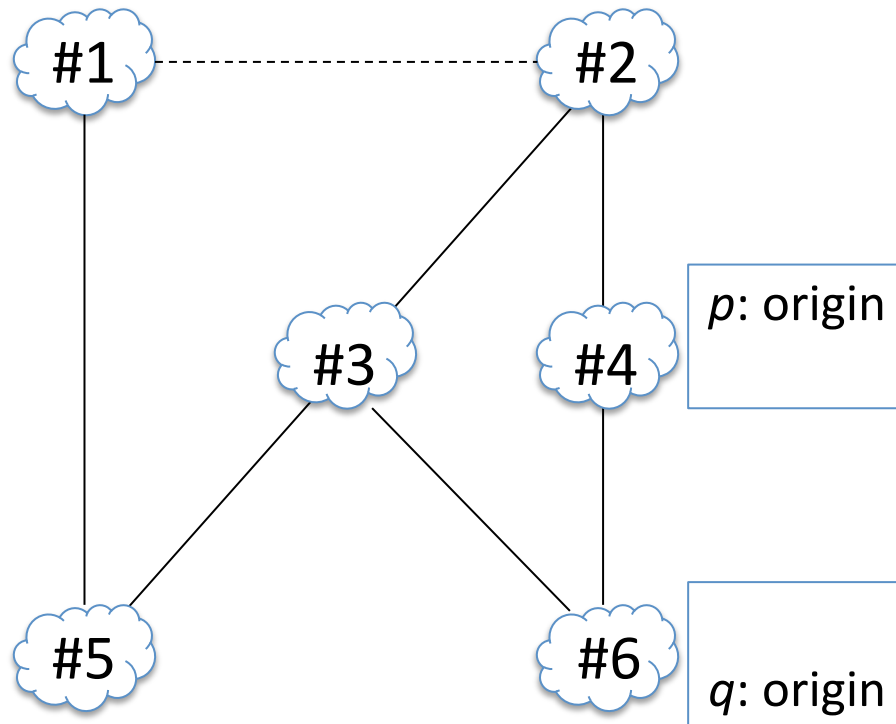
Filtering strategy

- Locally filter the more specific prefixes when possible
 - no black holes
 - respect routing policies
- Use built-in incentives to filter locally
 - save on forwarding state
 - forward along best route (dictated by routing policies)
- *Exchange routing information with standard BGP*

Providers, customers, and peers



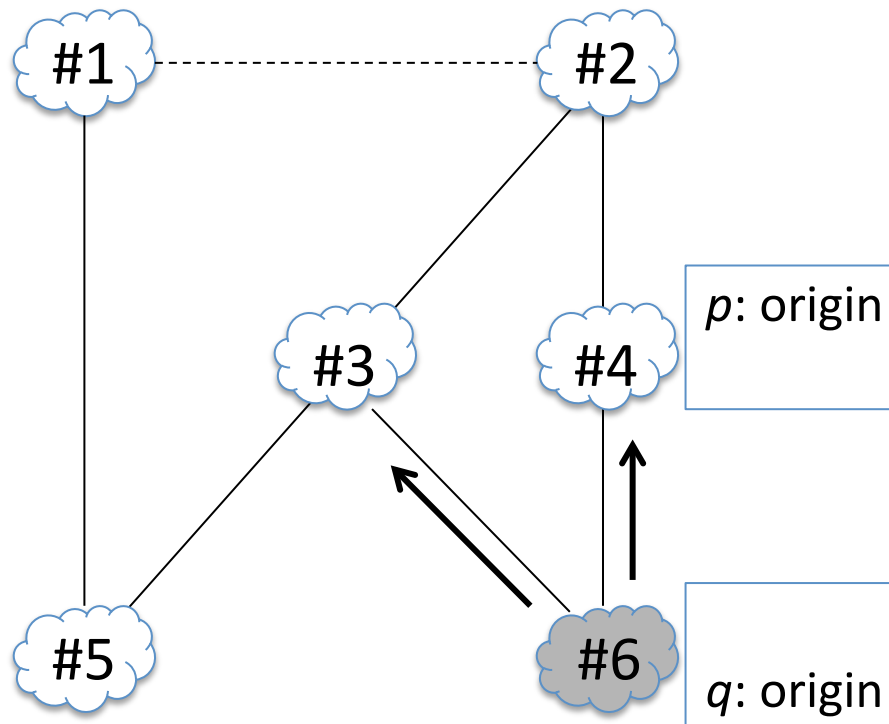
Prefixes



#6 originates *q* (1.0.0.0/24); #4 originates *p* (1.0.0.0/16)

q* more specific than *p

Routes

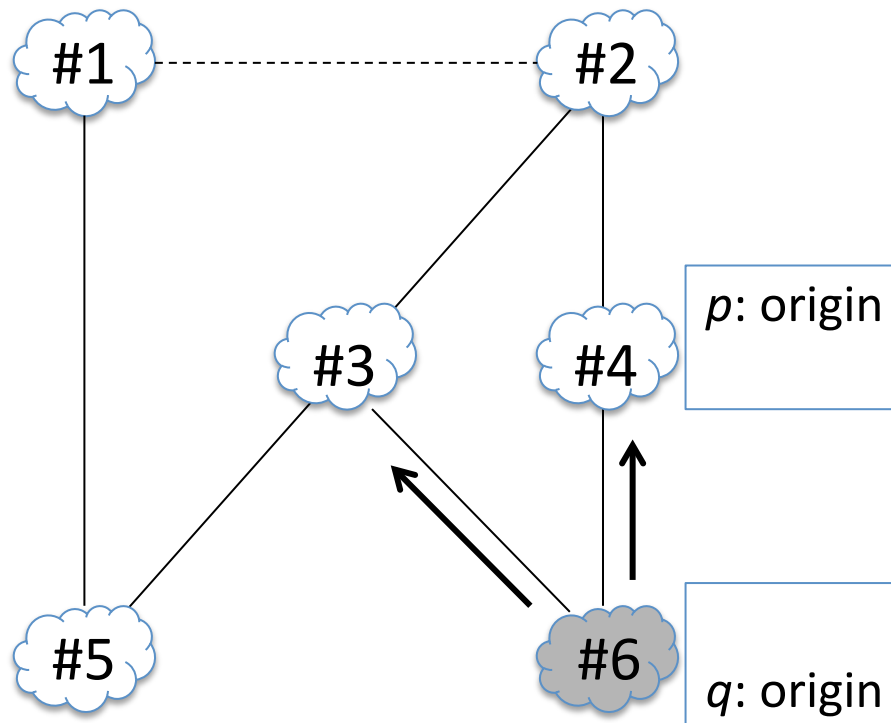


Route

Association between a prefix and an attribute, from a totally ordered set of attributes

\longrightarrow
 q -route
(route pertaining to q)

Gao-Rexford routing policies



route attributes:

“customer”

“peer”

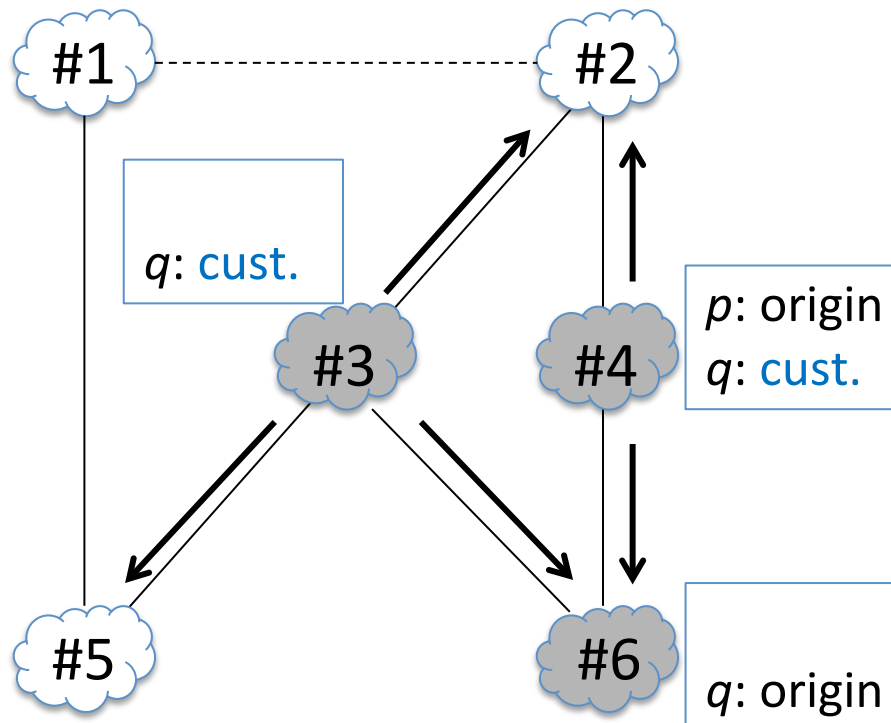
“provider”

→
q-route

preferences: customer then peer then provider

exportations: all routes from customers; all routes to customers

Gao-Rexford routing policies



route attributes:

“customer”

“peer”

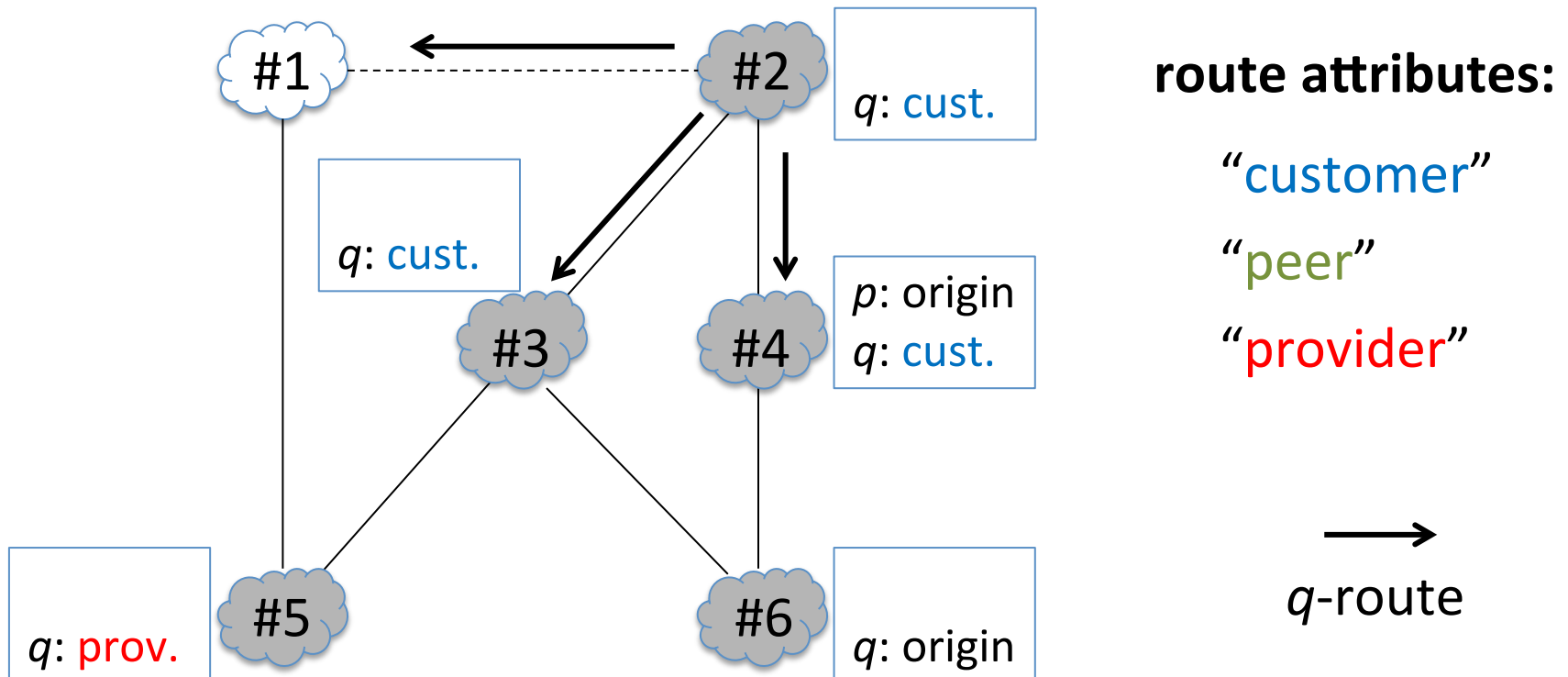
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q-route

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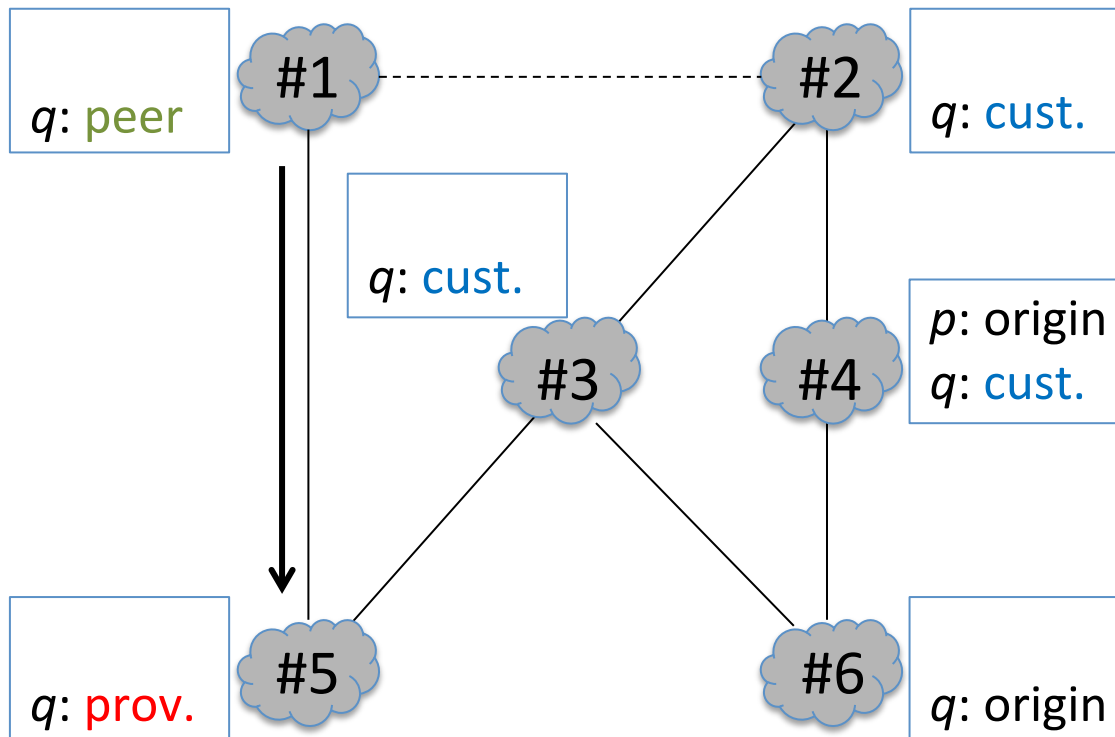
Gao-Rexford routing policies



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Gao-Rexford routing policies



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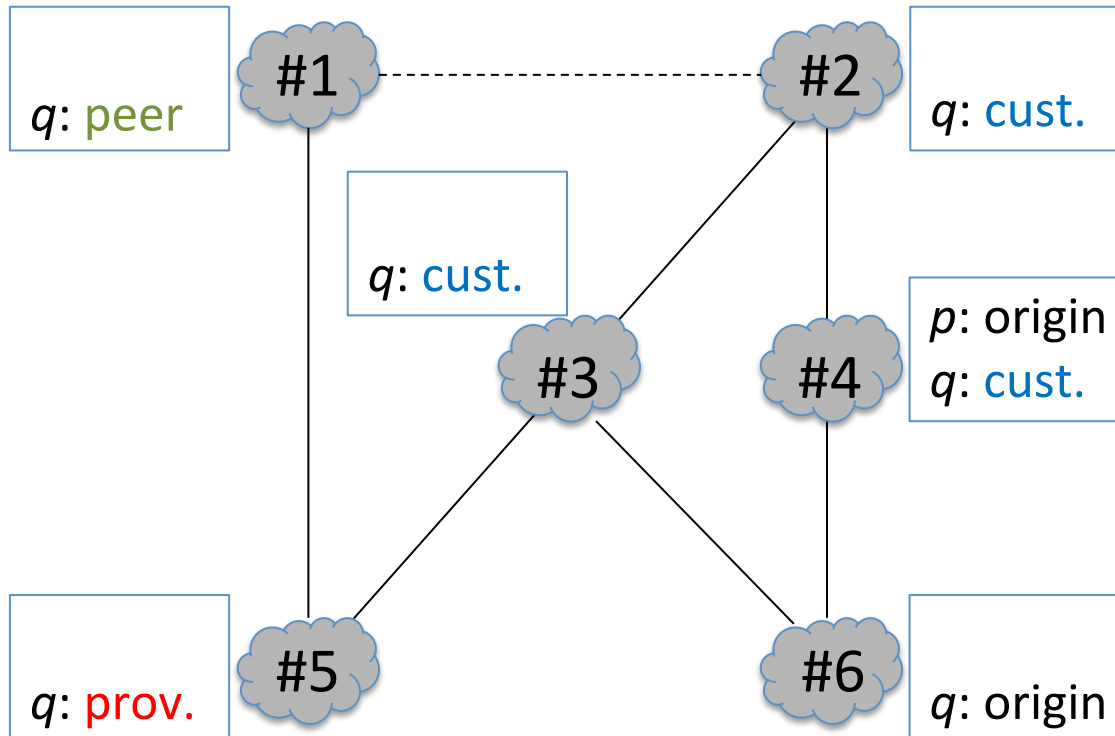
“provider”

→
 $q\text{-route}$

preferences: customer then peer then provider

exportations: all routes from customers; all routes to customers

Final state for prefix q



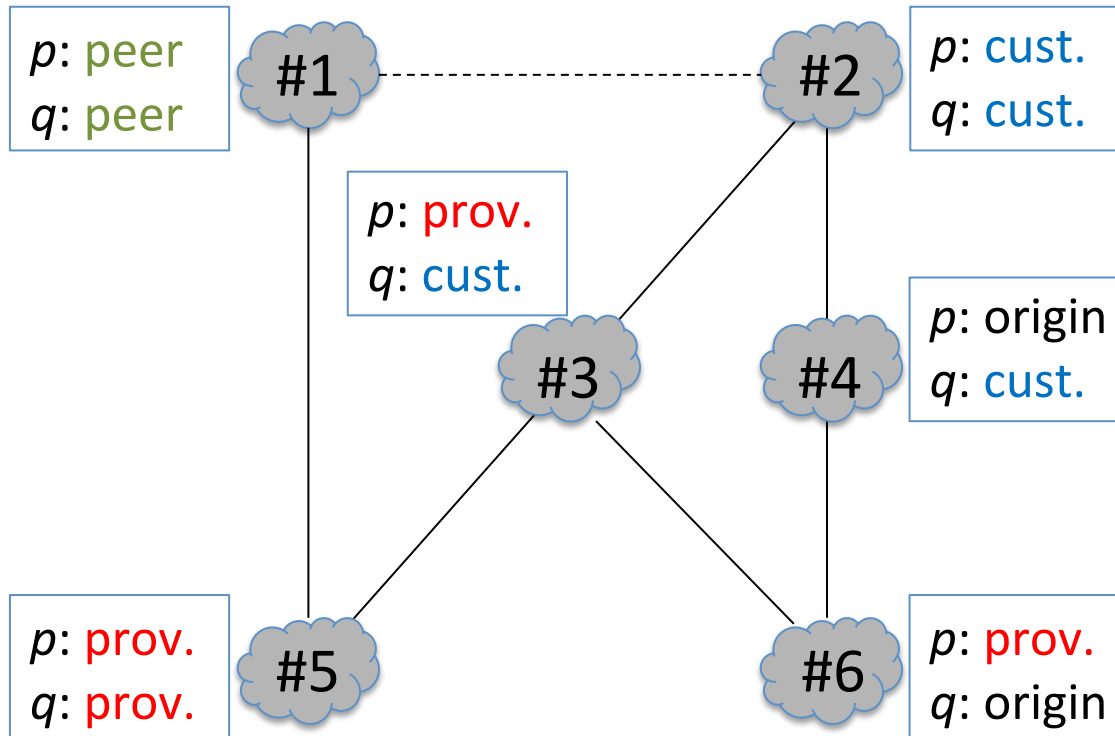
route attributes:

“customer”

“peer”

“provider”

Final state for prefixes q and p



route attributes:

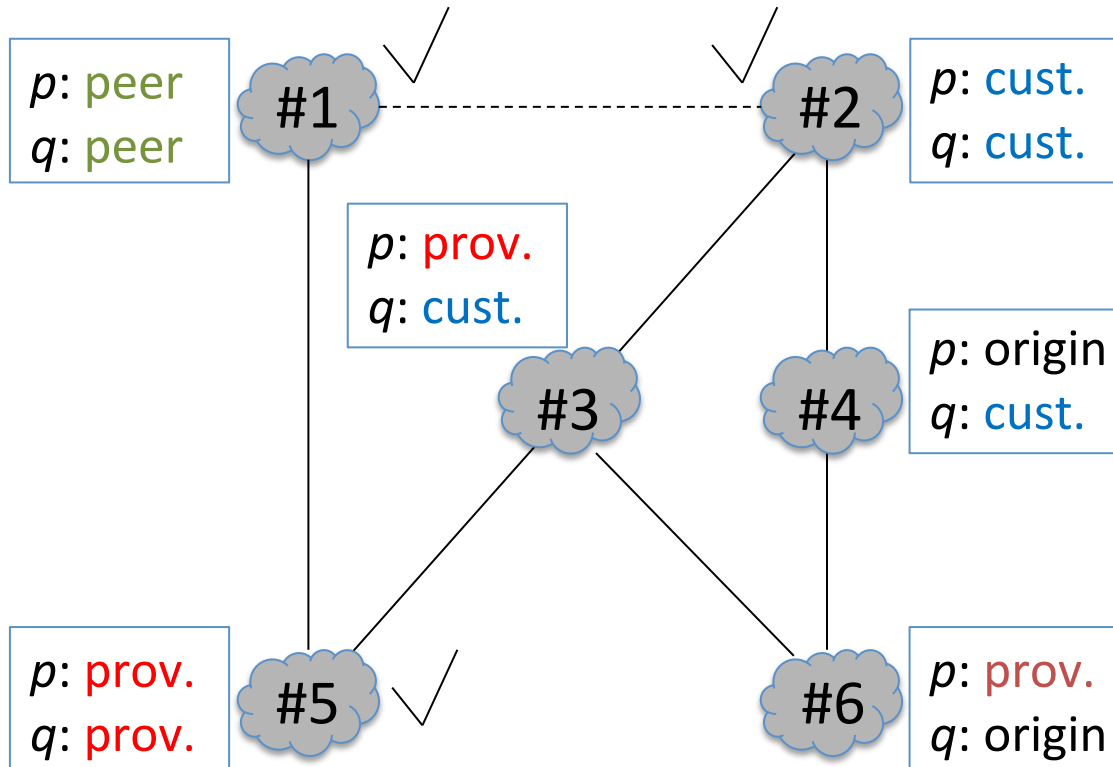
“customer”

“peer”

“provider”

forwarding: longest prefix match rule

Filtering code (FC)



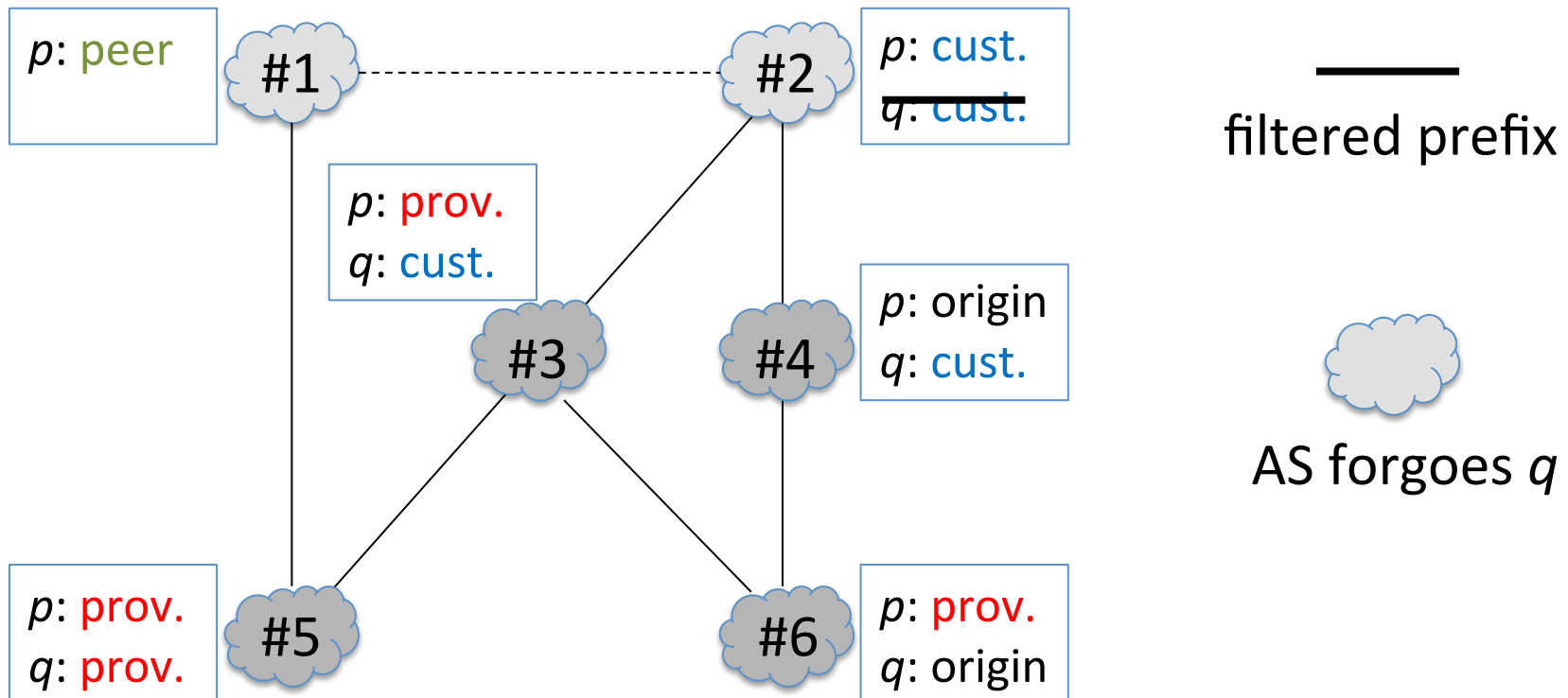
Filtering Code (FC)

Other than origin of p ,
in the presence of p ,
filter q if only if:

attribute of p -route
same or preferred to
attribute of q -route

✓ ASs that filter q upon execution of FC

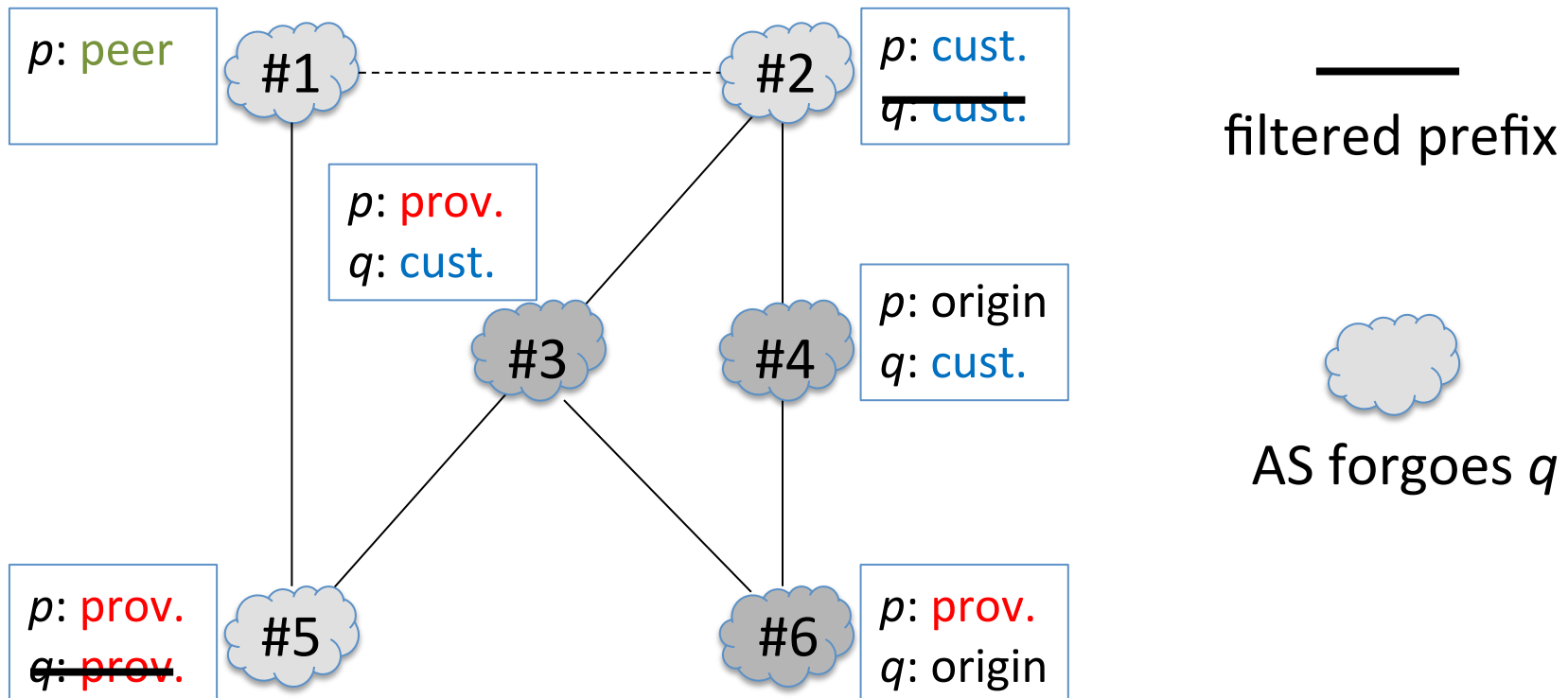
AS 2 applies FC



AS 2 filters q \rightarrow

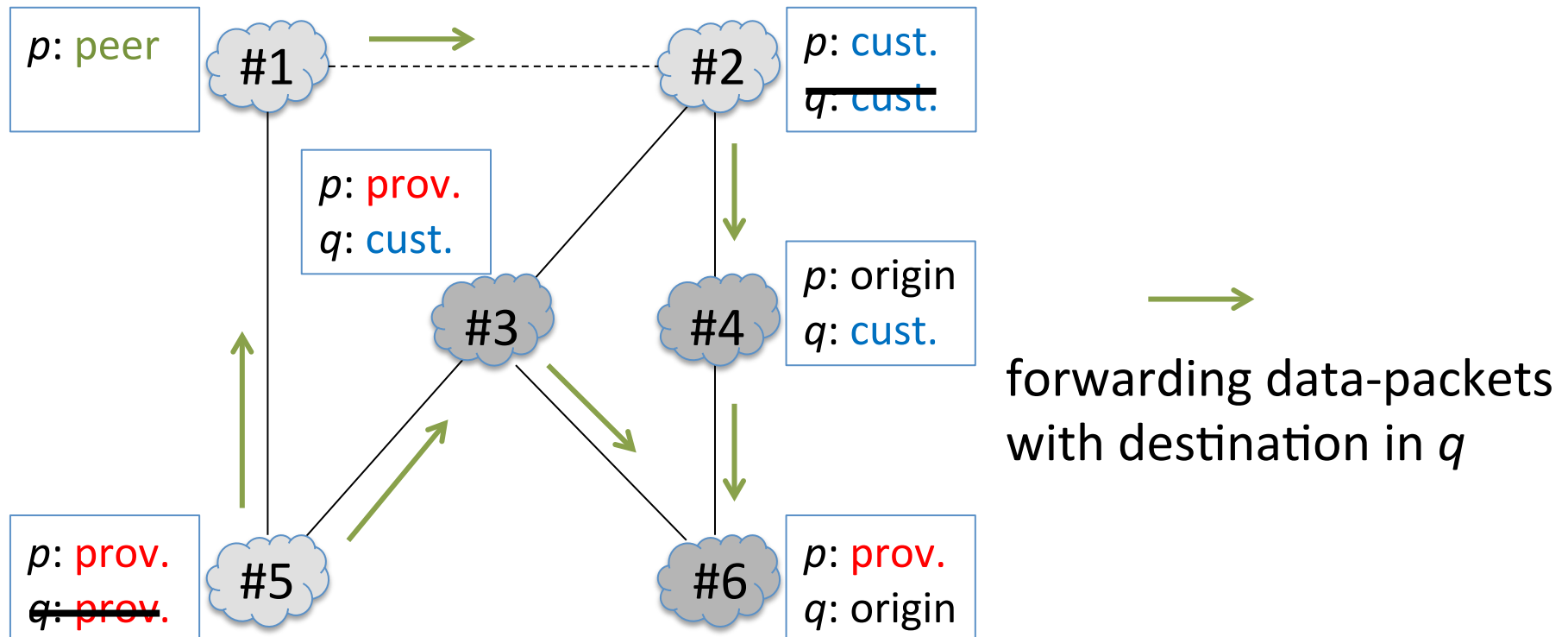
- AS 2 saves on forwarding state
- AS 1 is oblivious of q ; it saves on forwarding and routing state

All ASs apply FC



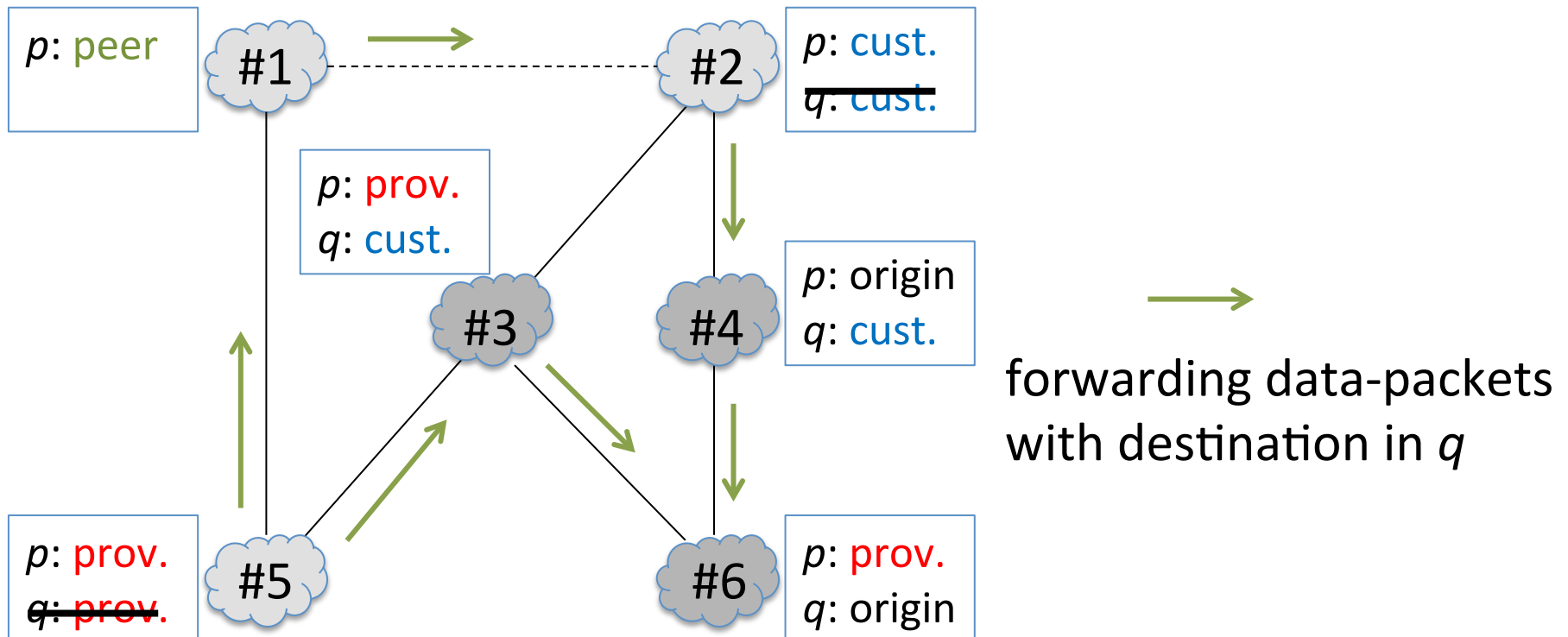
AS 1, AS 2, and AS 3 forgo q ➡ forwarding to q using less specific p

Global property: correctness



Correctness: no routing anomalies (no black holes)

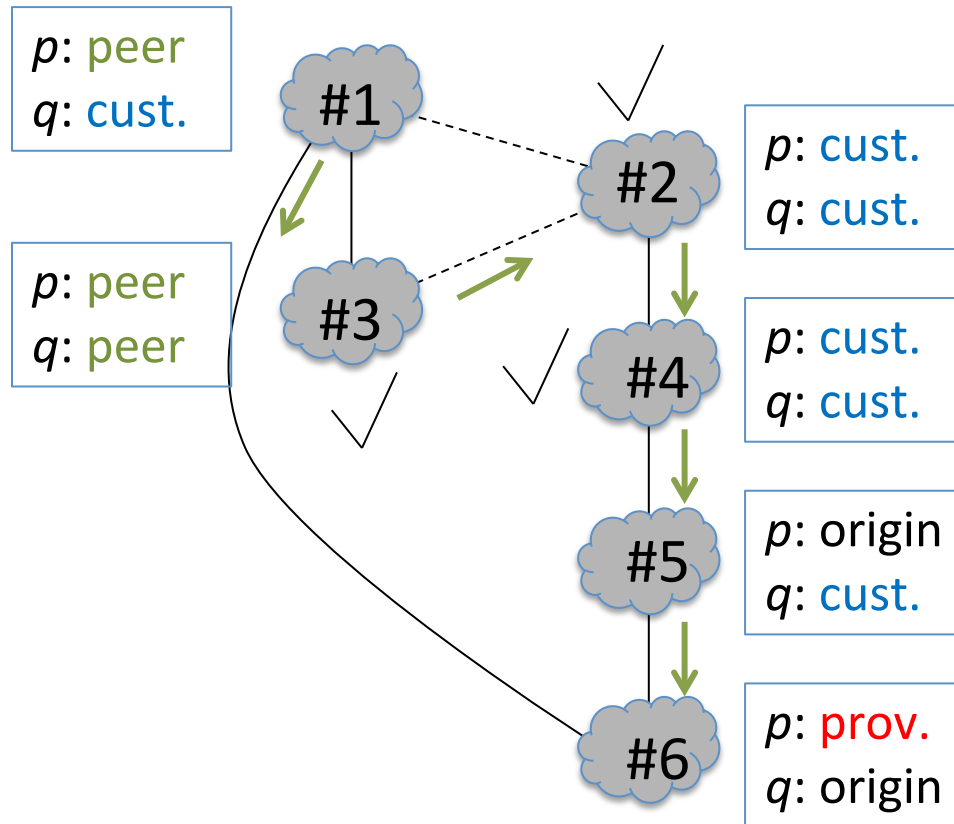
Global property: route consistency



Route consistency: attribute of route used to forward data-packets is preserved

Optimal route consistency: set of ASs that forgo q is maximal

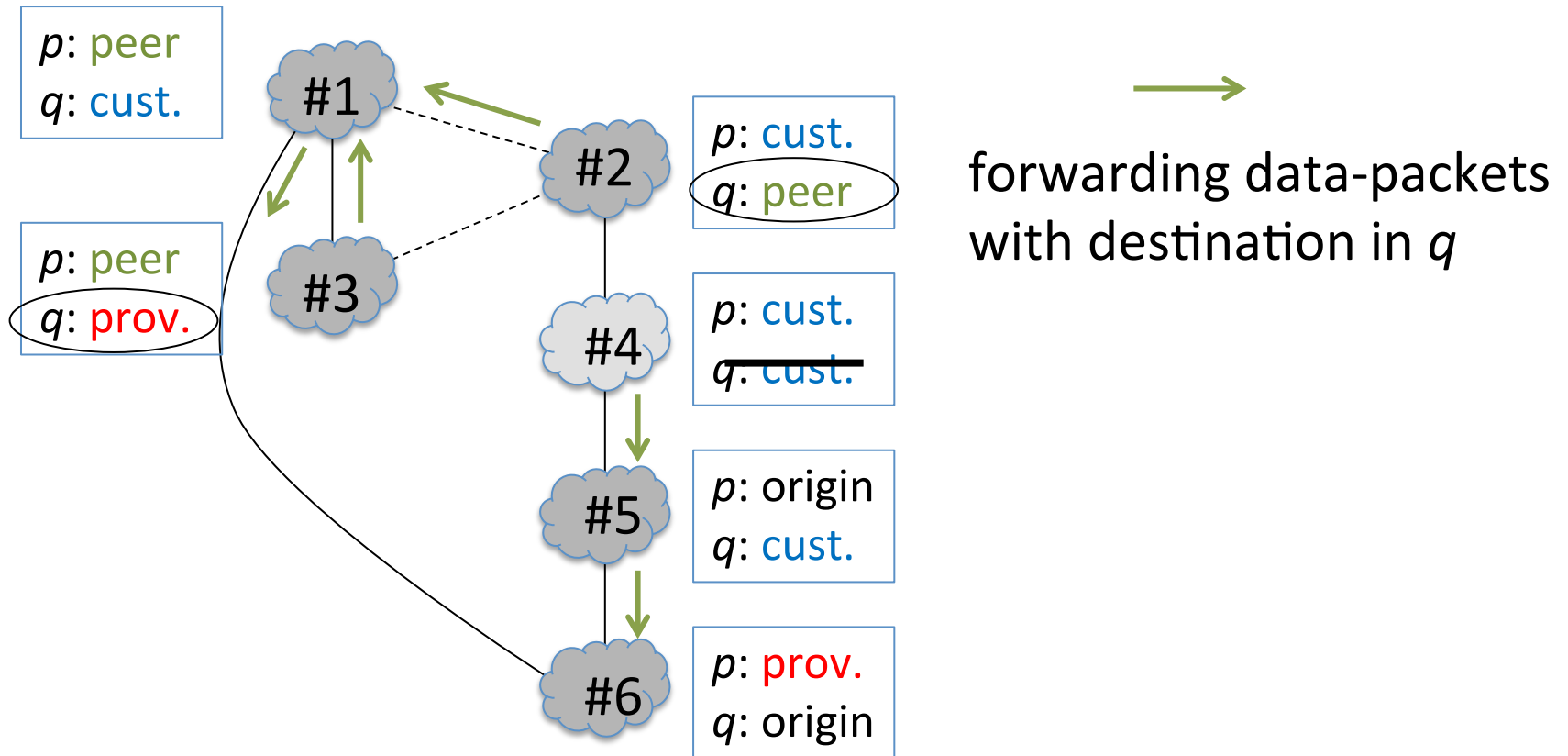
Partial deployment



→
forwarding data-packets
with destination in q

✓
ASs that filter q upon
execution of FC

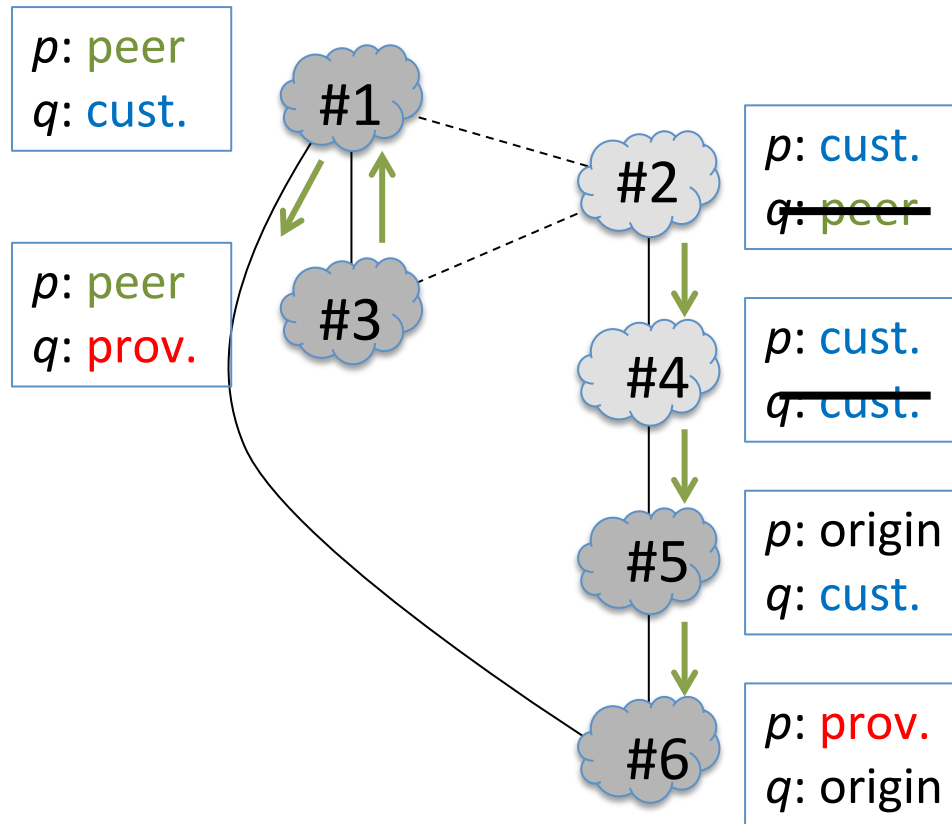
Partial deployment: incentives



AS 2 (and AS 3) has a double incentive to apply the FC:

- saves on forwarding state
- improves attribute of route used to forward data-packets

Partial deployment: incentives



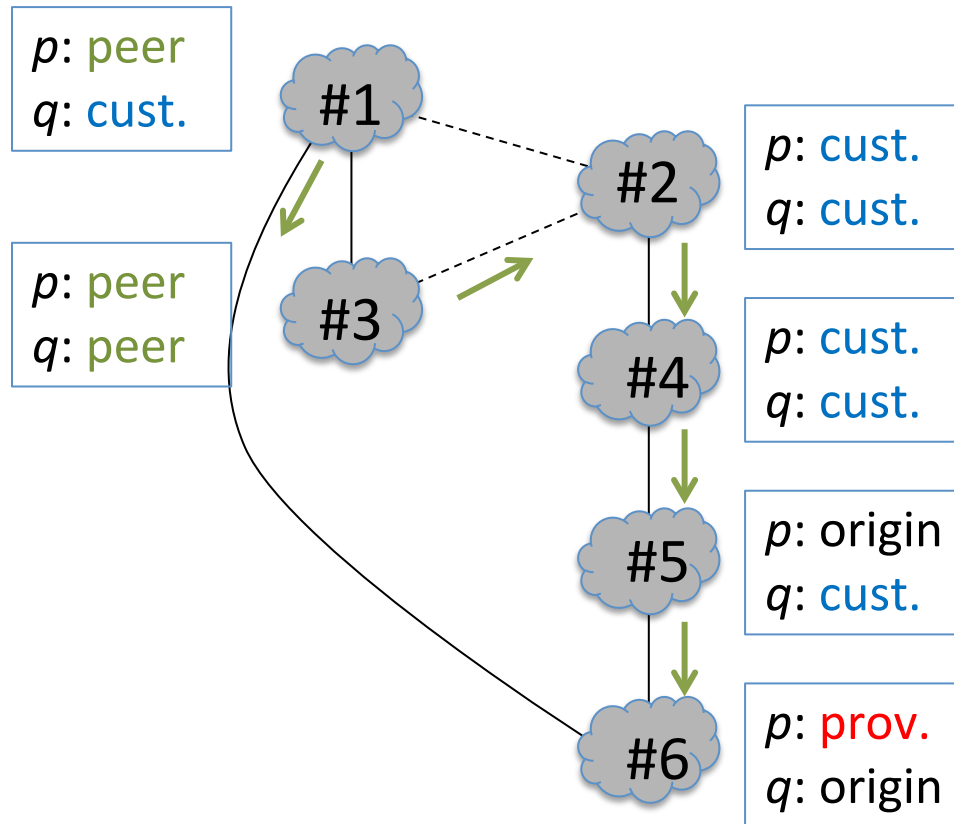
→
forwarding data-packets
with destination in q

AS 2 applies FC



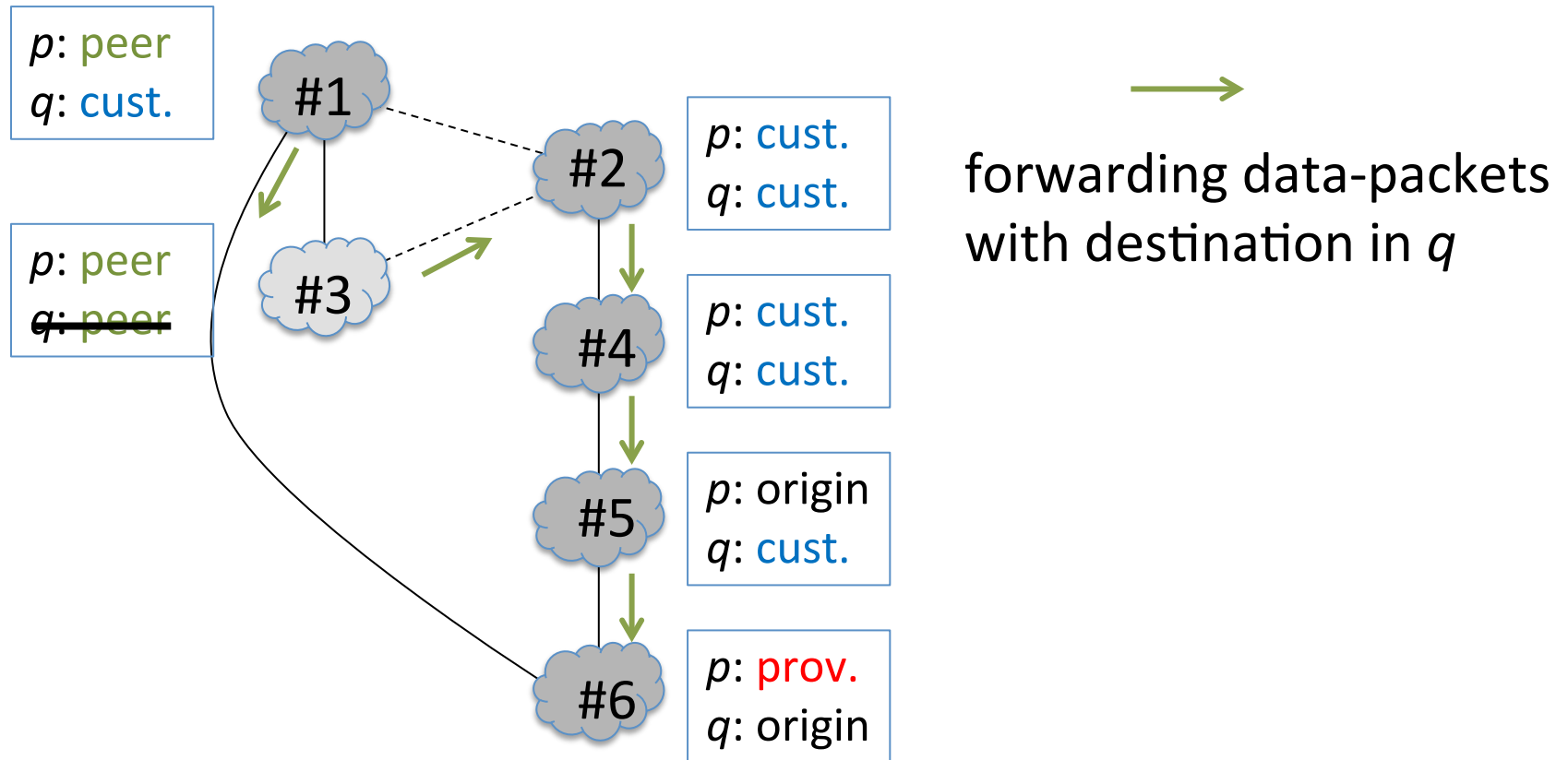
AS 2 reverts to forwarding data-packets with address in q to AS 4

Partial deployment: route consistency



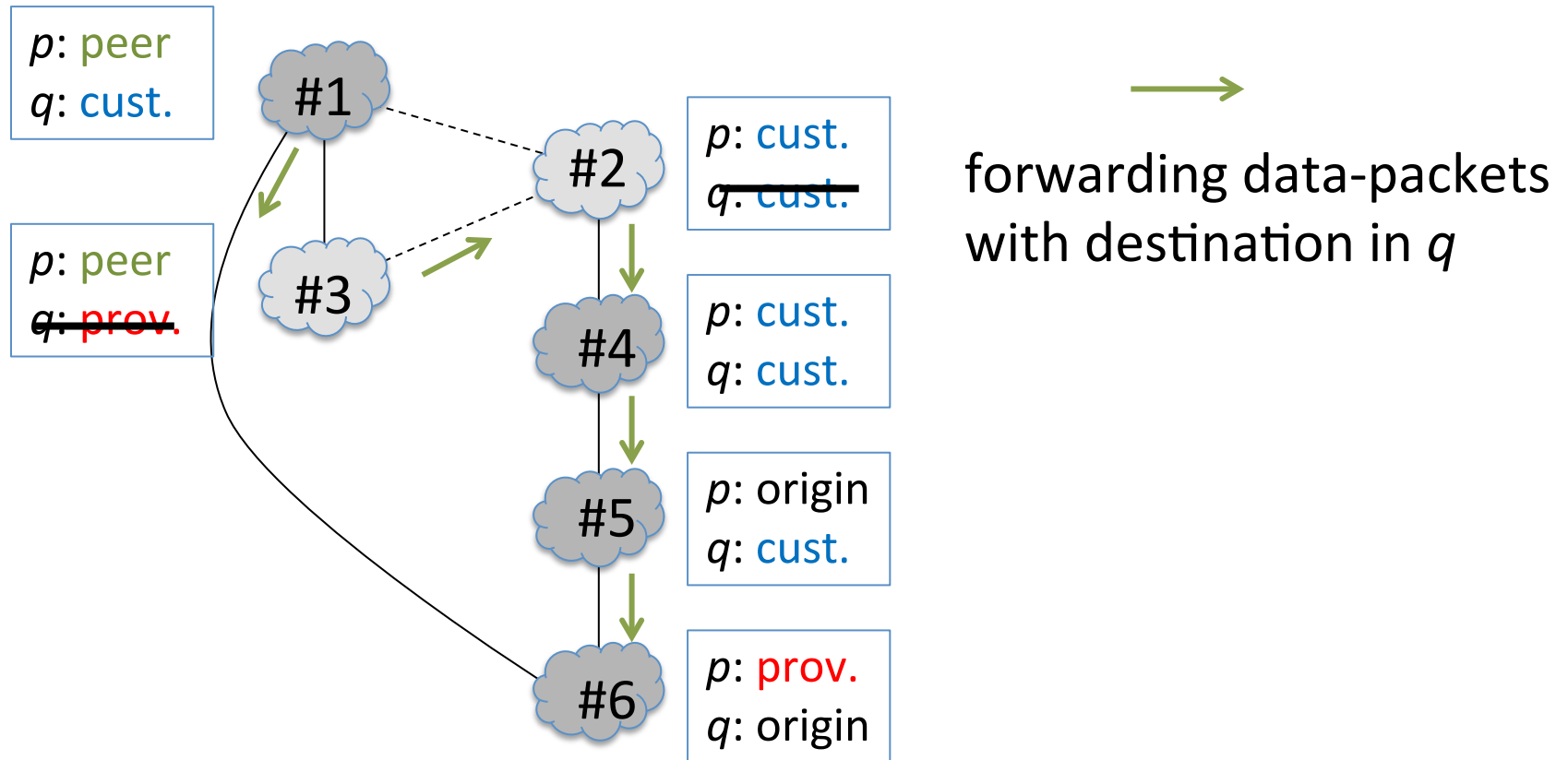
→
forwarding data-packets
with destination in q

Partial deployment: route consistency



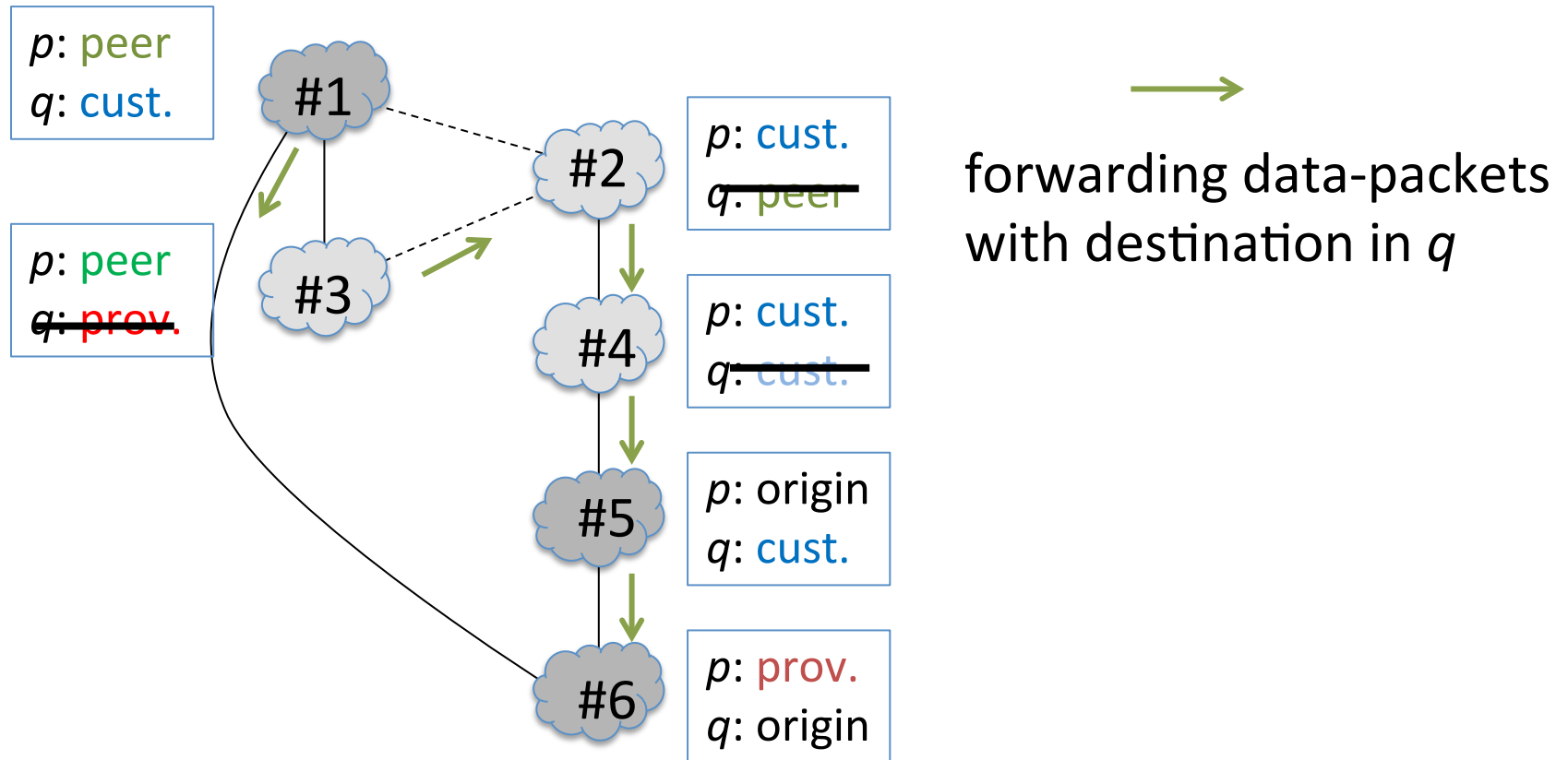
First to apply FC are ASs that elect a peer or provider q -route

Partial deployment: route consistency



Next to apply FC are ASs for which providers have already applied FC

Partial deployment: route consistency



Next to apply FC are ASs for which providers have already applied FC

Filtering strategy: general case

- Trees of prefixes learned from BGP
 - FC for a prefix in relation to the parent prefix
 - Correctness
 - for the routing policies for which BGP is correct
 - Route consistency (optimal and through partial deployment)
 - for *isotone* routing policies (includes Gao-Rexford)
- Optimal route consistency is not synonymous with *efficiency* (think shortest paths)**

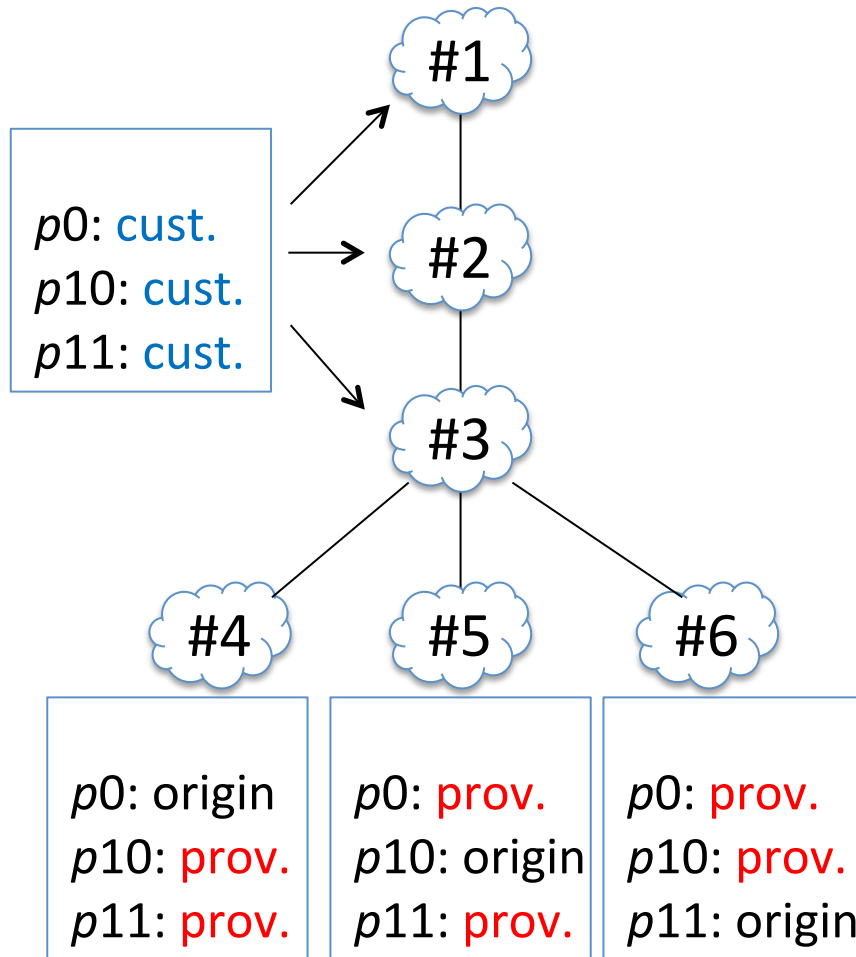
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Aggregation strategy

- Locally originate aggregation prefixes when beneficial
 - new address space is *not* created
 - allow filtering of provider-independent prefixes
 - self-organization when more than one AS originates the same aggregation prefix
- *Again, exchange routing information with standard BGP*

Aggregation prefix

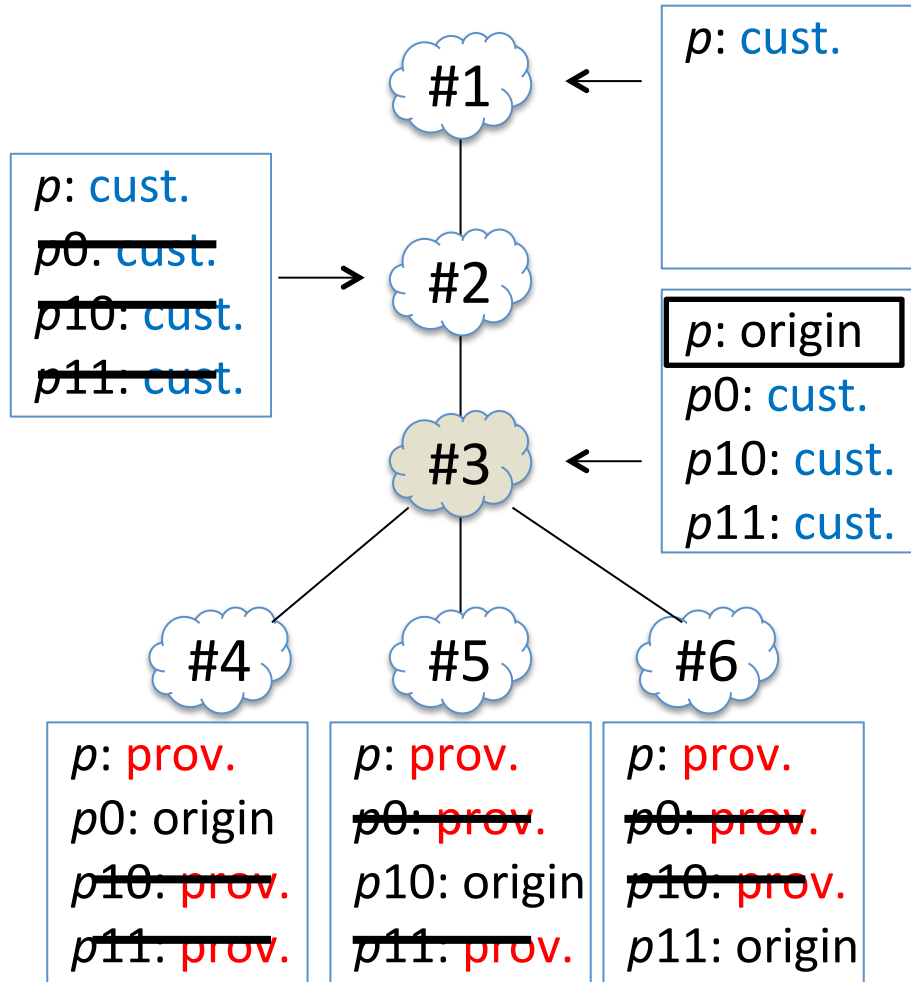


Aggregation prefix

1. no routable address space is created
2. at least two covered prefixes
3. customer route is elected for each of the covered prefixes

$p0 + p10 + p11 = p$; p is an aggregation prefix at AS 3

AS 3 originates p



AS 1 is oblivious of $p0$, $p10$, and $p11$

AS 2 filters $p0$, $p10$, and $p11$

AS 4 filters $p10$ and $p11$

AS 5 filters $p0$ and $p11$

AS 6 filters $p0$ and $p10$

Aggregation strategy: general case

- Trees of prefixes learned from BGP
 - aggregation prefixes cover parentless prefixes
- Self-organization
 - for the routing policies for which BGP is correct
- Optimal origins
 - for *isotone* routing policies (includes Gao-Rexford)

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Data-sets

- Annotated topology (CAIDA, Feb. 2015)
 - ~50K ASs; ~42K stub ASs
 - ~94K provider links; ~94K customer links; 180K peer links
- IPv4-prefixes-to-ASs mapping (CAIDA, Feb. 2015)
 - ~530K prefixes
 - ~270K parentless prefixes
 - ~210K prefixes have same origin AS as parent

FIB filtering efficiency: definition

Normalized amount of reduction brought
by DRAGON to the forwarding tables of an
AS

$$\text{FilterEff} = \frac{\# (\text{FIB entries BGP}) - \# (\text{FIB entries DRAGON})}{\# (\text{FIB entries BGP})}$$

FIB filtering efficiency: results

	Basic DRAGON filtering	Full DRAGON filtering & aggregation
Min. FilterEff	47%	
% of ASs with at least Min. FilterEff	100%	
Max. FilterEff	49%	
% of ASs attaining Max. FilterEff	87%	

FIB filtering efficiency: results

	Basic DRAGON filtering	Full DRAGON filtering & aggregation
Min. FilterEff	47%	69%
% of ASs with at least Min. FilterEff	100%	100%
Max. FilterEff	49%	79%
% of ASs attaining Max. FilterEff	87%	87%

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- **Conclusions**

Conclusions

- DRAGON is a BGP add-on to scale the Internet routing system
- DRAGON can be deployed incrementally
- DRAGON reduces the amount of forwarding state by approximately 80%
- DRAGON is – more fundamentally – a solid framework to reason about route aggregation

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Thank you!