# SENSS: Software-defined Security Service

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### **Growing DDoS Attacks**

#### Average monthly size of DDoS attacks (Gbps)



### **Growing DDoS Attacks**

At March 2013, DDoS flooded Spamhaus at 300Gbps,

200 times faster than average



# **Growing Prefix Hijacking**

In 2013, prefix hijacking affects 1,500 prefixes, 150 cities Live interception attacks are on for more than 60 days Traffic from major financial companies, govs, ISPs diverted



# Motivation

- Network attacks are more frequent and powerful
  - In Q1 2014: 47% increase in total DDoS attacks.
  - Attack size more than 300Gbps.
- Network attacks are more damaging
  - 71% of data center operators report DDoS attacks
  - DDoS on Bitcoin Exchanges lowered bitcoin price from 700\$ to 540\$
- Diverse attacks
  - Data plane: Direct flooding, reflector attacks 5
  - Control plane: Prefix hijacking, interception

### **Traditional Solutions**

- Victim-based solutions are not sufficient
  - Leverage IDS boxes, or outsource to security services
  - Hard to diagnose remote root causes or trace sources
  - Have to manually call ISPs on the phone
- Research inter-ISP solutions are not adopted
  - Focus on individual attacks

Need a flexible, deployable solution for diverse attacks

# Software-defined Security Service (SENSS)

- Victim-oriented programming for diverse attacks
  - Victims have the incentives
  - Victims have knowledge of their traffic and priorities
- Victims request help from remote networks
  - To observe and control their own traffic and routes
  - Using simple and expressive interfaces at ISPs, easily implemented in today's ISPs

7

 Difficult trade-offs: all the intelligence implemented at the victim



# **SENSS** is Practical

#### ISPs

- SENSS-needed interfaces already exist in their infrastructure
- ISPs already provide manual support for victims
- ISPs can charge victims for the security services

#### Victims

- Strong incentives to fix their own problems
- Effective solutions even with partial deployment

### Challenges

- What's the right interface at ISPs?
  - Easy to implement at today's ISPs
  - Useful for a wide variety of attack defenses
- How can victims program the defenses?
  - With SENSS deployment on a few ASes
  - Without missing information (spoofing, privacy, etc)
- Security and Privacy

Simple, Flexible Interfaces at ISPs

# **SENSS Interfaces: Traffic**

- Traffic query
  - Query flows using TCP/IP header fields
  - Answer #bytes/pkts from/to each neighbor
- Traffic control

- Filter, rate limit traffic matching a traffic flow

- Similar to OpenFlow rules
  - Only allow victims to query/control traffic to/from them

```
1. src=1.2.*.*, dest=3.4.5.* \rightarrow query
```

```
2. src port=80, dest=3.4.^{*.*} \rightarrow filter
```

# **SENSS Interfaces: Routes**

#### Route query

- Query the best route to the victim prefix
- Similar to BGP route queries to neighbors
- But we extend to remote ASes

#### Route control

- Modify the route from the AS to the victim
- Demote all the routes with given AS segments
- To get around the malicious/polluted ASes

# Automated Detection/Mitigation at Victims

# DDoS w/ Signature

- The victim identifies the attack and header signature
- The victim installs filtering es at deployed AS **S6** 10 10 TrafficFilter (<header\_signatur 100 100 100 e>\_\_\_ Β С Α **S5** 8 8 . 8 8 V < 209 G Η F **S4** Ε 3 3 3 3 Κ J **S**3 201 201 201 200 **S1 S2**

# DDoS w/ Signature

- The victim identifies the attack and header signature
- The victim installs filtering s at deployed AS



# **DDoS** Without Signature

- Victims may not find a signature Spoofing; randomize packet header and contents
- Cannot simply block high traffic aggregates - May lead to high collateral damage
- SENSS: Compare traffic distribution across ASes before and after the attack
  - Track normal traffic distribution periodically
  - Compare with traffic distribution during attack
  - Filter on those AS links with big traffic growth
  - Only victim can decide which collateral damage is OK 17









<S2, 200>





<S2, 200>

### **Interception Attacks**

- Interception attacks
  - Introduce false information into the routing system
  - Claim shorter AS-PATH, hijack victim prefix
  - Traffic still reaches the victim
- Detection and mitigation
  - Data plane alone cannot reveal the root causes
     Control plane info may be inaccurate or outdated
- SENSS: Check inconsistency between control and data planes via route and traffic query



# Interception



#### Knowledge Base Control Plane S to V: <FMAV> F to V: <MAV>

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#### Knowledge Base Control Plane S to V: <FMAV> F to V: <MAV>

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Knowledge Base Control Plane S to V: ₅≤F\_MAV> F to V: <MAV>

#### Data Plane Traffic from S to V passes through B and C!!

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Knowledge BaseControl PlaneDataS to V: <FMAV>TraffF to V: <MAV>B ar

Data Plane Traffic from S to V passes through B and C!!

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### Interception



Knowledge BaseControl PlaneCS to V: <FMAV>TF to V: <MAV>E

Data Plane Traffic from S to V passes through B and C!!

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### **SENSS Use Cases**

Attacks	Query	Control
DDoS w/ signature	Traffic queries	Traffic filter
DDoS w/o signature	Traffic queries	Traffic filter
DDoS reflection	Reduces to DDoS w/	′ or w/o signature
Crossfire	Traffic queries	Bandwidth guarantees
Blackholing	Route queries	Route demotion
Interception	Route and Traffic queries	Route modification

# **Simulation Setup**

- AS-Level Internet topology from RouteViews/ RIPE
  - 41K ASes with 92K links, including 11 Tier-1 ASes
- Simulate DDoS
  - Real traffic from CDN traces and DDoS attack traces
  - Simulated traffic with different distributions
- Simulate Prefix-Hijacking
  - Select victims and attackers from different tiers in the AS hierarchy

# **DDoS Results**

- Eliminate attack traffic
  - To eliminate 95% attack traffic
  - Need only 10-30 SENSS ASes
  - Less than 36 messages are needed
  - Hold for a wide range of traffic distributions
- Small collateral damage
  - Outperforms traceback solutions with the same # of deployed ASes

# **Prefix Hijacking Results**

#### Detection

- With 30 ASes deployed, the detection accuracy can reach 90% for blackholing and 70% for interception
- The median number of queries is 3-10 for blackholing and 6-15 for interception
- Mitigation
  - Correct > 80% of polluted Ases with 18 SENSS
     ASes

# **SENSS** Implementation

- ISPs
  - Openvswitch as data plane, Quagga as control plane
  - Floodlight as controller for Openvswitch
  - Apache SENS
- Victim
  - Sends HTTPs
     requests to
     SENSS server
- Response time
   600 ms



# **Security and Privacy**

#### • Security

- Operations allowed on traffic from/to own prefixes and routes to own prefixes
- Ownership verification via RPKI, communication via SSL
- Outsource to cloud if victim has no path to SENSS server

#### Privacy

- ISPs only need to share traffic information for peer indexes, without revealing the actual peer
- Routing information is already publicly available

# Conclusion

- Software-defined security service
  - Simple, flexible interfaces at ISP
  - Victim-oriented programming for diverse attacks
- Practical security detection/mitigation services
  - Effective to mitigate large-scale attacks
  - Incentive for adoption from ISP and victims
  - Flexible for supporting new defenses for new attacks

# Adopting SENSS

- We will release SENSS for deployment

   Contact Minlan Yu (<u>minlanyu@usc.edu</u>)
- We want to hear from operators
  - What are your concerns in deploying SENSS?
  - Economics? Privacy? Effectiveness? Deployment?

http://www-bcf.usc.edu/~minlanyu/writeup/ons14senss.pdf