Measurement-based Inter-domain Traffic Engineering

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Border 6

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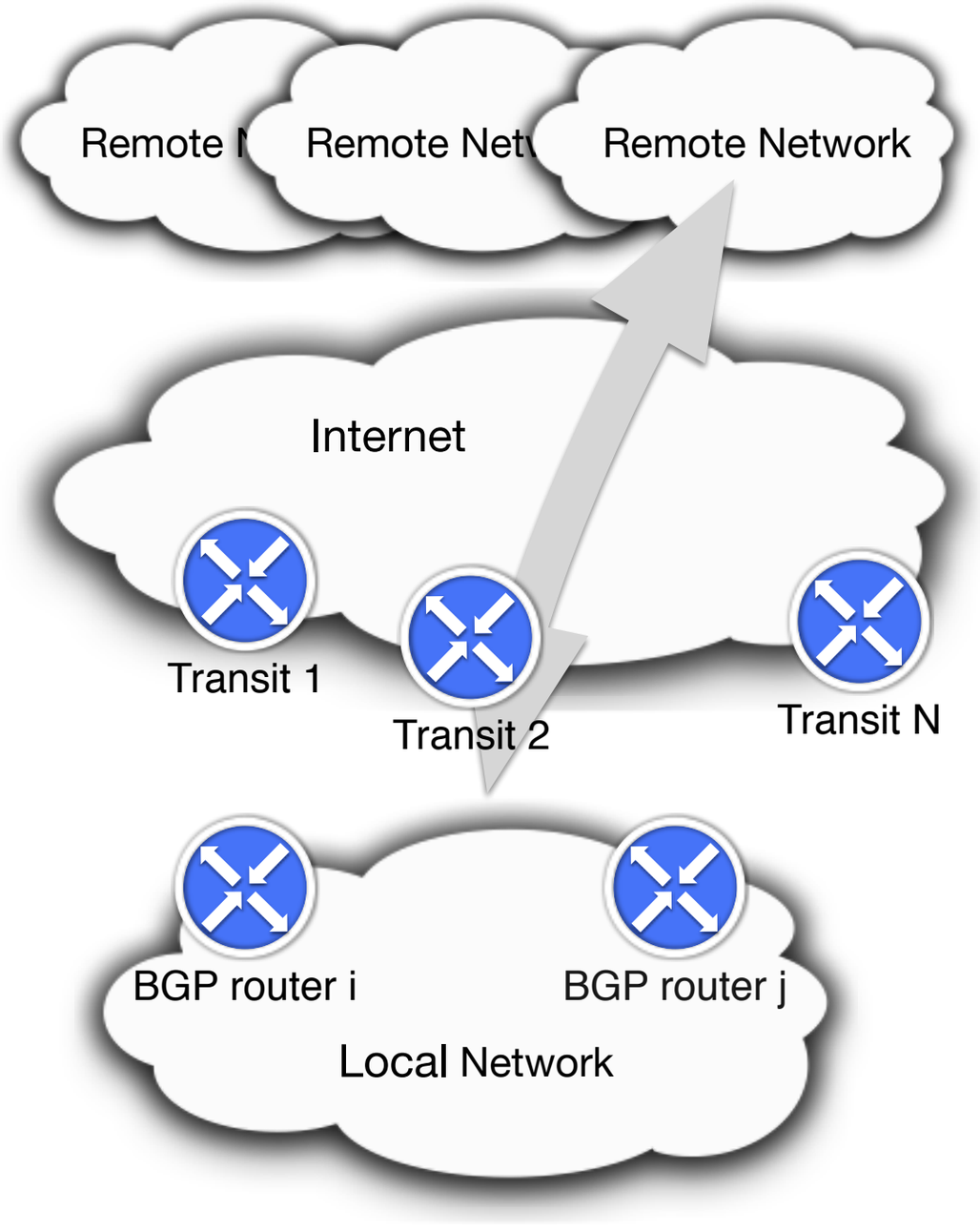
Outline

- **Measurement-based** inter-domain traffic engineering

- A problem of scalability

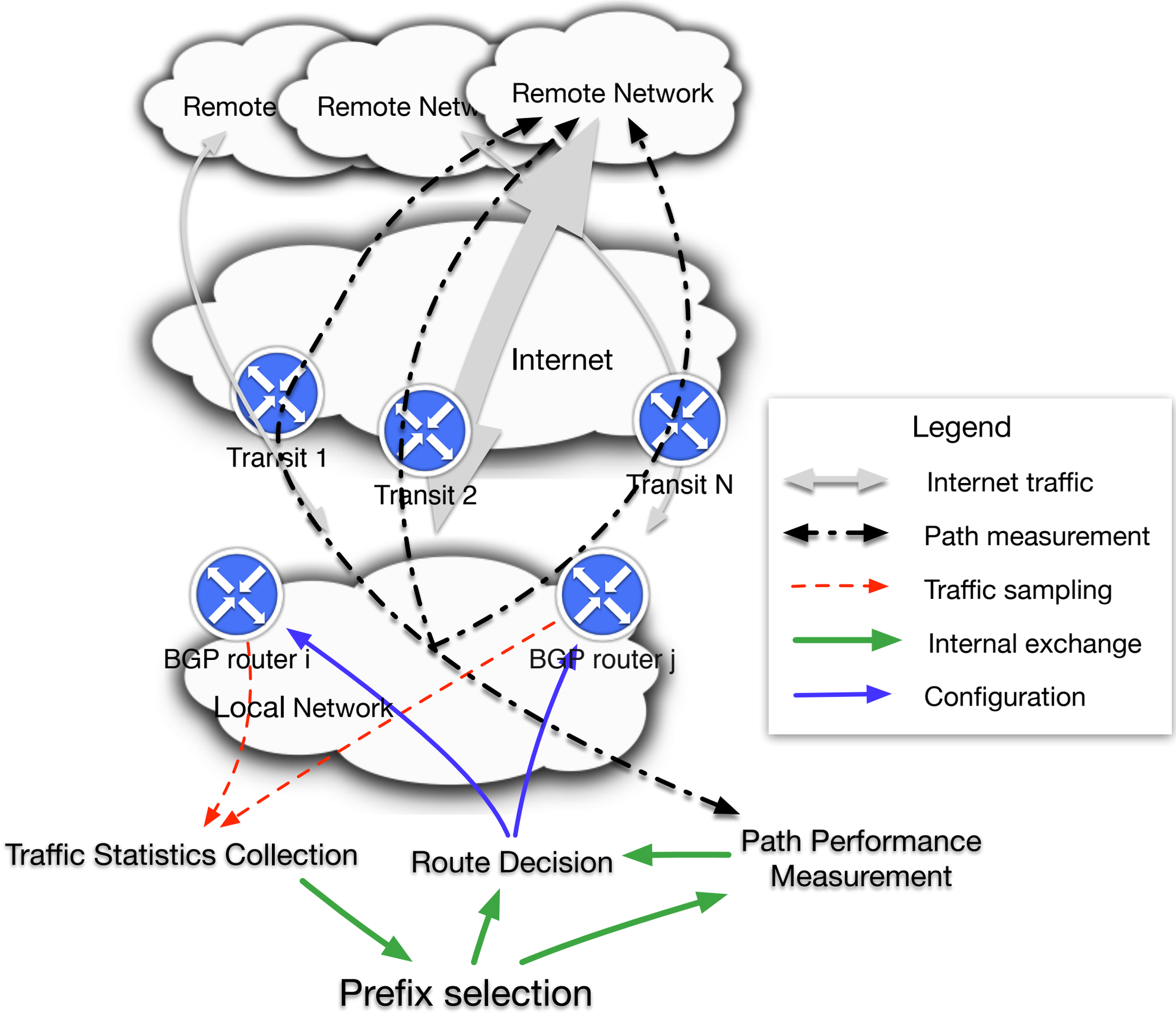
A scenario: Out-bound TE for multi-homed stub AS



1. Local_Pref
2. AS_Path
3. MED
4. eBGP > iBGP
5. tie-breaking
    1. IGP costs
    2. oldest path
    3. etc

An old scenario

- N. Feamster, D. G. Andersen, H. Balakrishnan, and M. F. Kaashoek, "**Measuring the effects of internet path faults on reactive routing**," ACM SIGMETRICS, vol. 31, no. 1, p. 126, **2003**.
- A. Akella, B. Maggs, S. Seshan, A. Shaikh, and R. Sitaraman, "**A measurement-based analysis of multihoming**," SIGCOMM, **2003**.
- D. K. Goldenberg, L. Qiu, H. Xie, Y. R. Yang, and Y. Zhang, "**Optimizing cost and performance for multihoming**," CCR, vol. 34, no. 4, p. 79, Oct. **2004**.
- A. Akella, B. Maggs, S. Seshan, and A. Shaikh, "**On the Performance Benefits of Multihoming Route Control**," IEEE/ACM Trans. Netw., vol. 16, no. 1, pp. 91–104, Feb. **2008**.


Remote
Remote Netw
Remote Network
Internet
Transit 1
Transit 2
Transit N
BGP router i
BGP router j
Local Network
Traffic Statistics Collection
Route Decision
Path Performance Measurement
Prefix selection
Legend
Internet traffic
Path measurement
Traffic sampling
Internal exchange
Configuration

Traffic statistics collection

Purpose: to select a set of ‘managed prefix’, i.e. destinations of importance;

Collector: netflow/sflow collector, PMACCT;

Storage: RAM, PostgreSQL.

Active measurements

Target: probes discovered in ‘managed’ prefixes

Method: TCP SYN -> RTT and loss;

Path: via **all available** BGP next-hops;

Steering: source-based routing, SDN, etc;

Storage: RAM, PostgreSQL.

Route decision

Objective: performance, availability & transit cost;

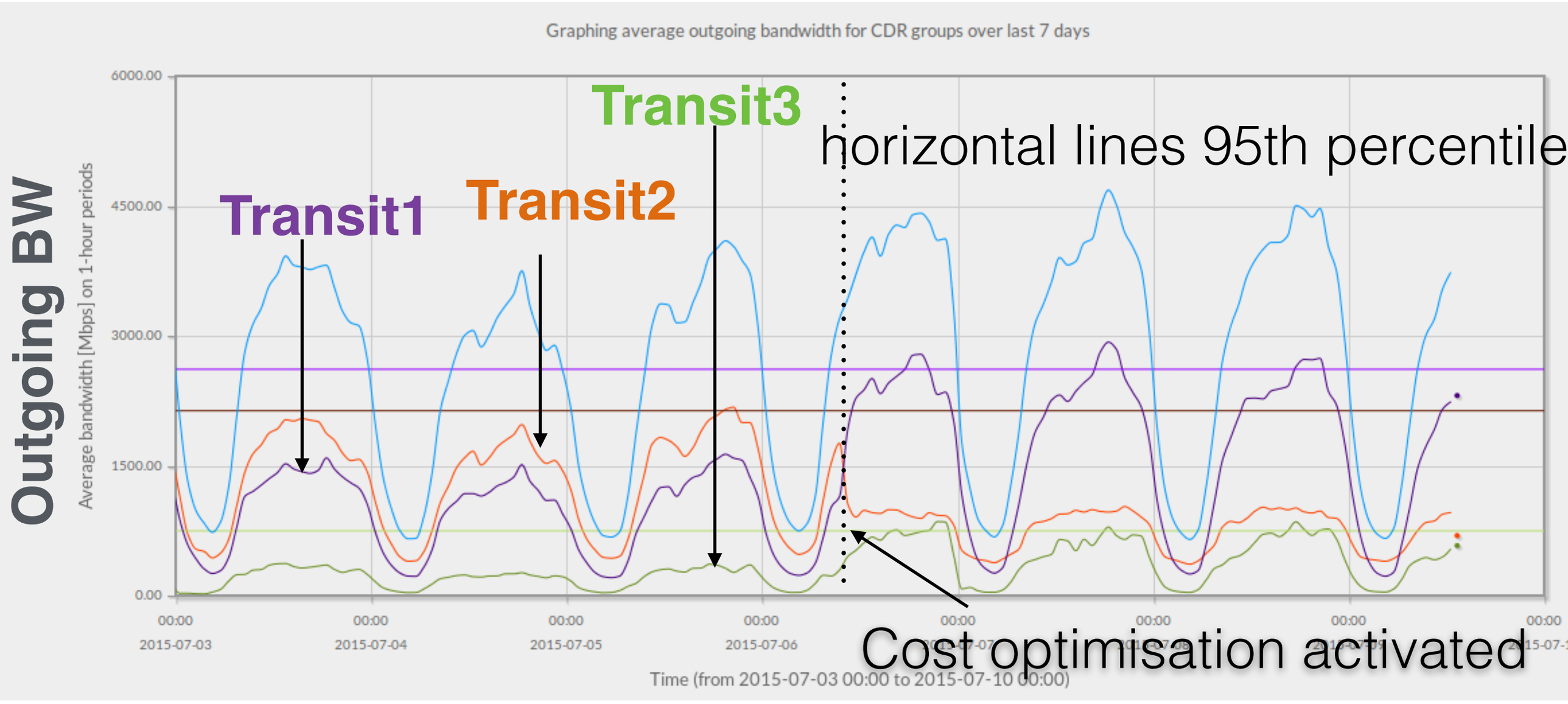
Metrics: RTT, loss and BW w.r.t. CDR;

Algorithm: depends on user's needs;

Steering: BGP as SDN southbound interface.

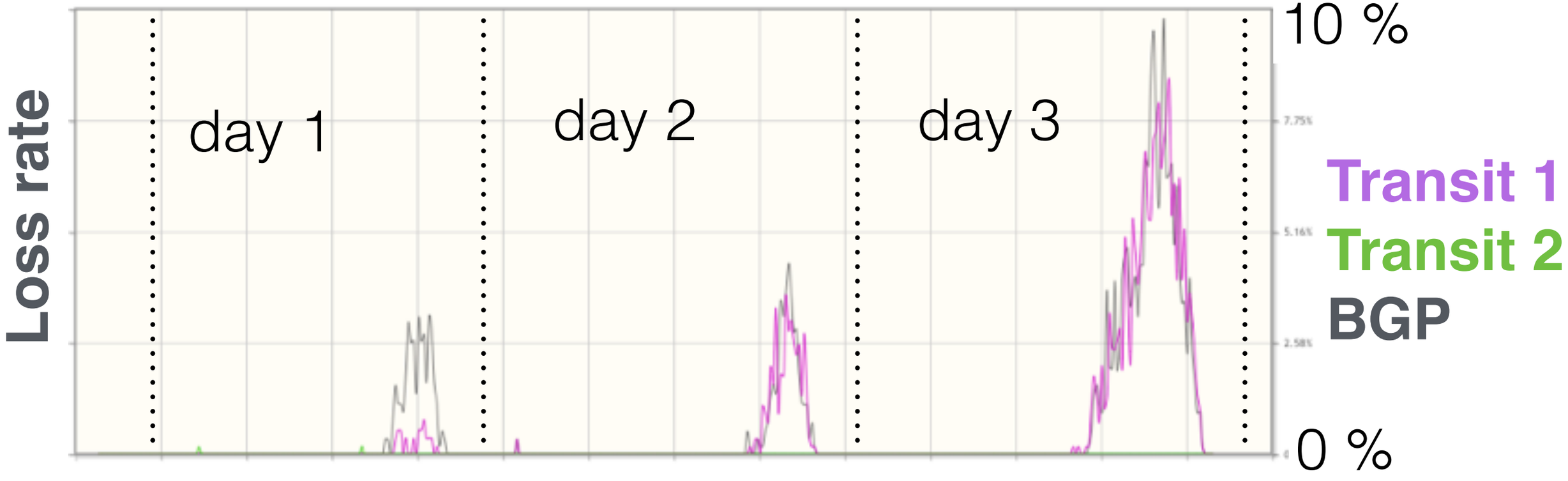
Algo example—Cost

Total **Transit1** **Transit2** **Transit3**



Transit1 has the lowest BW price.

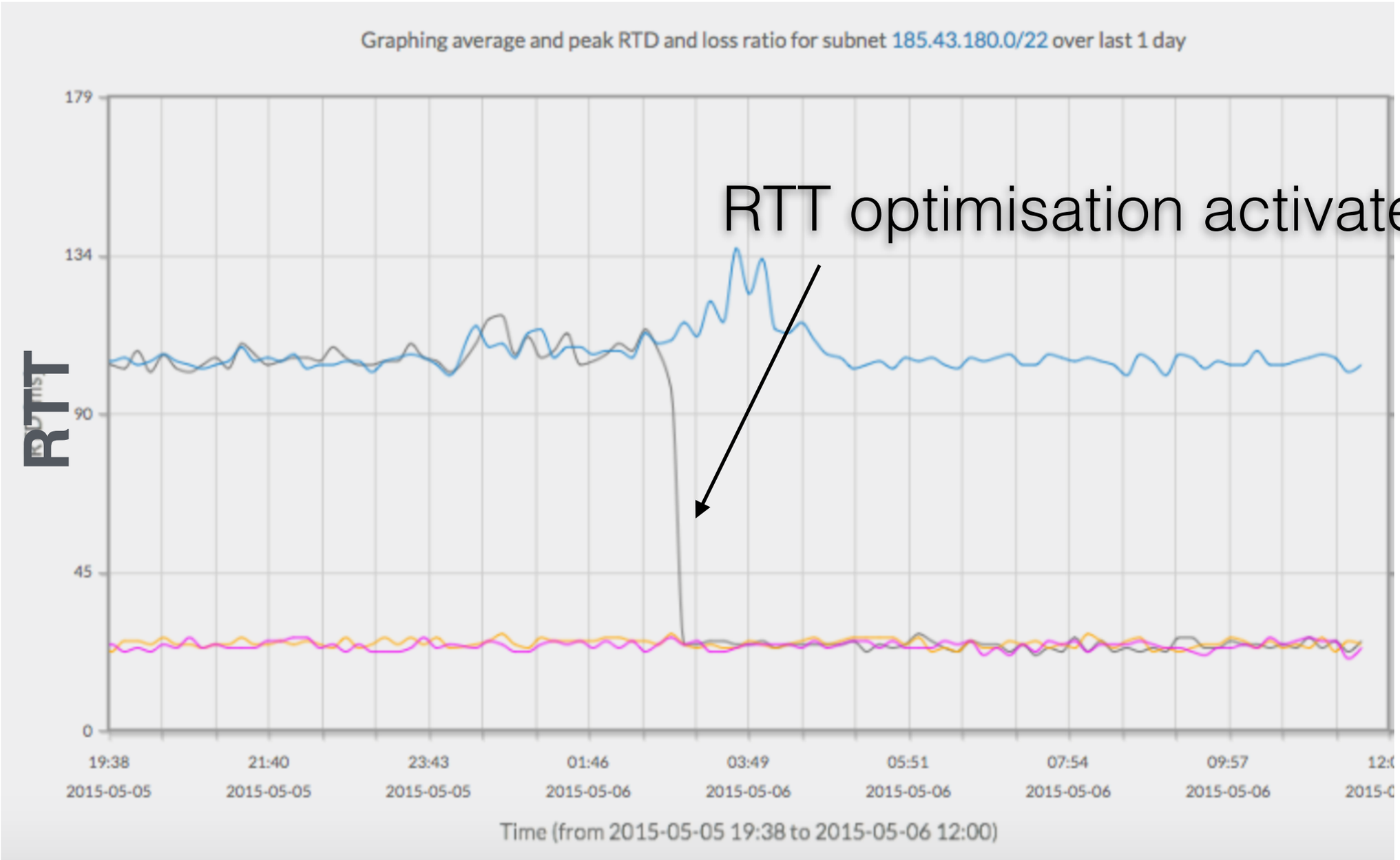
Algo example—Availability



Packet loss due to **consistent congestion**
can be avoided
by simply change a BGP next-hop.

Algo example—smaller RTT

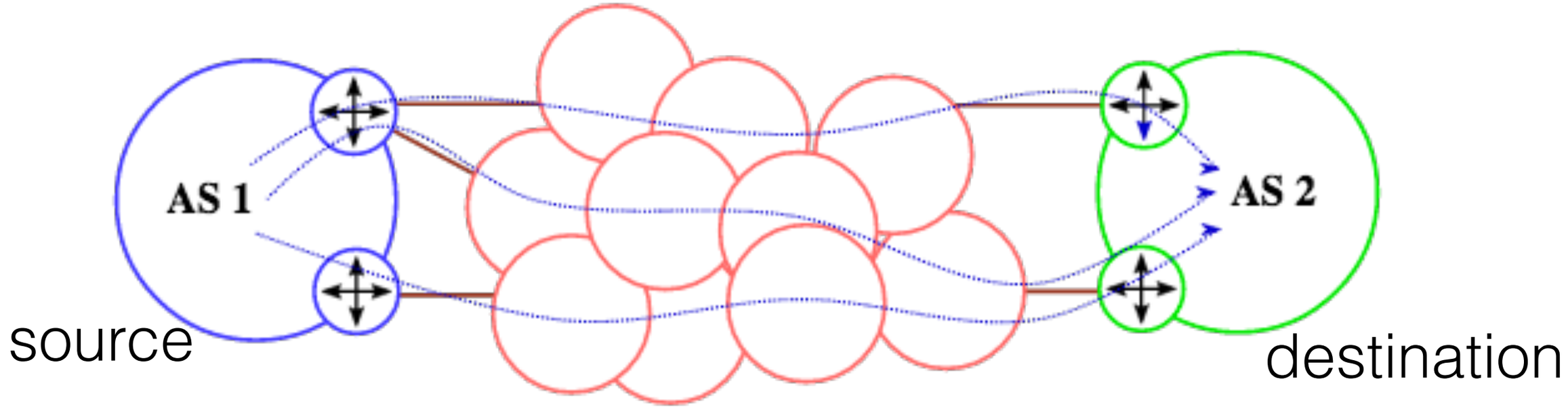
Transit 1 **Transit 2** **BGP**



For Inbound as well…

Route Preference Protocol (RPP) for **Inbound** TE

1. Tell traffic sourcing AS what is you favourite ingress point;
2. Traffic sourcing AS then does its best.



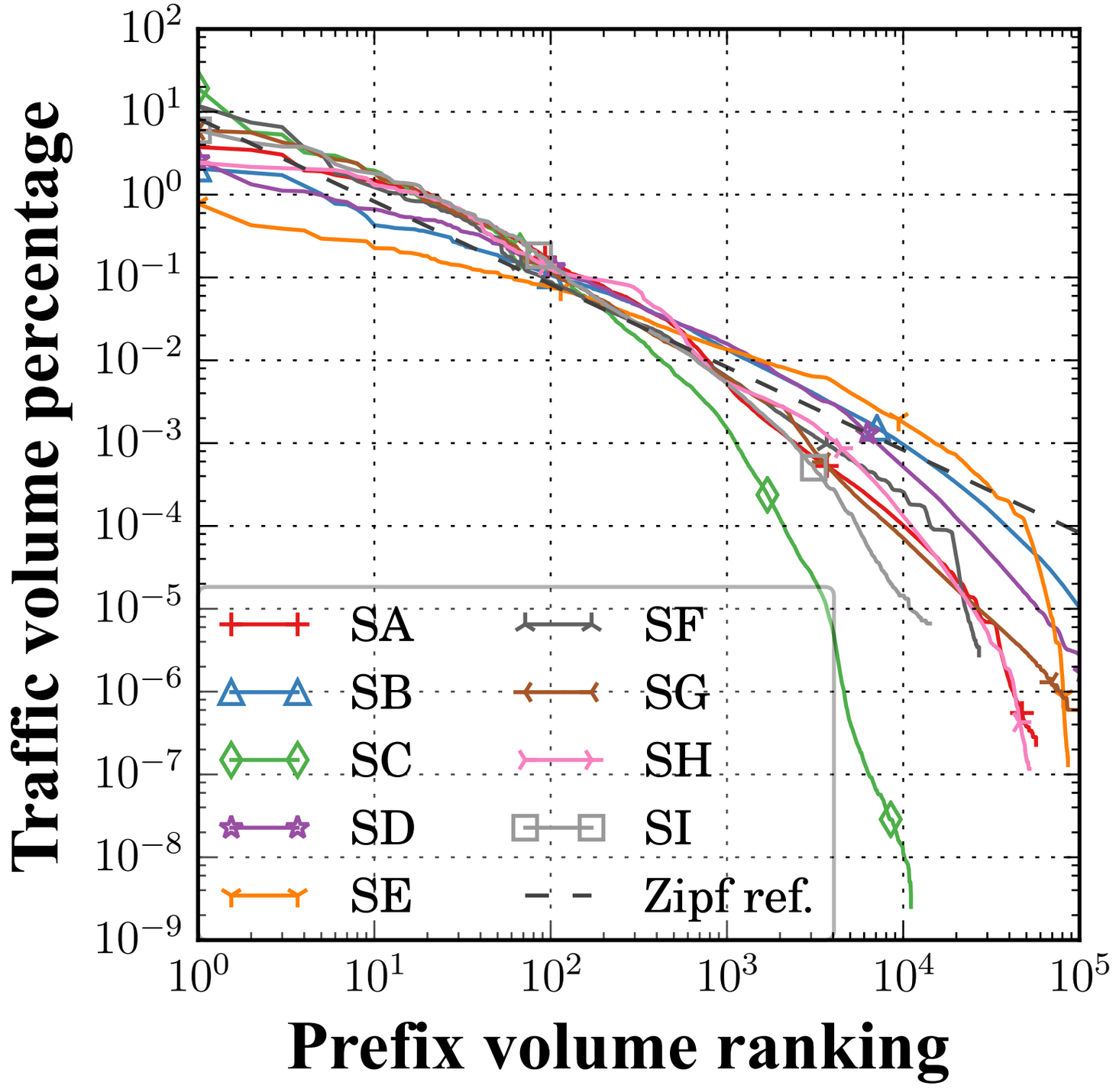
http://rpp.border6.com/
https://github.com/border6/rpp

the scalability problem
~600K BGP prefixes

How to **select prefixes** of importance?

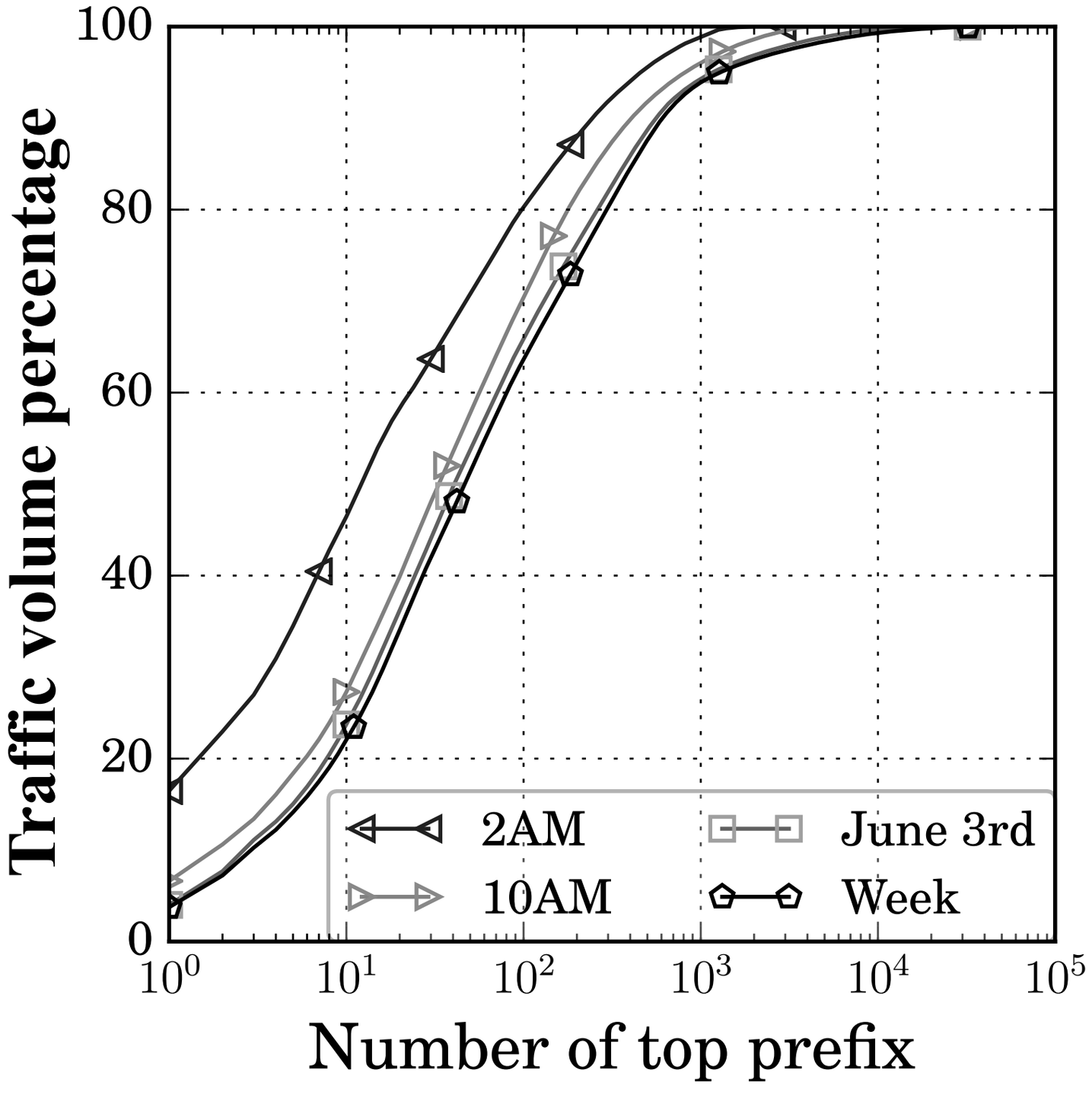
W. Shao, L. Ianonne, J.L. Rougier, F. Devienne, and M. Viste, “**Scalable BGP Prefix Selection for Effective Inter-domain Traffic Engineering**,” IEEE/IFIP NOMS, 2016.

A large portion traffic is concentrated on a small number of prefixes (log-log graph not CDF)



However…

traffic value associated to each prefix evolves over time.



However…

we have **some ten thousands** of time-series/prefixes to predict.

Predictively select BGP destination prefixes that stand for **a large portion** of traffic, with a **simple and efficient** method.

Wisely save sources for measurement and optimization.

Save resources of **data plane** as well.

full BGP routing table → one optimized route for each selected prefix + one default route for the rest

expensive edge router → cheaper DC switch/ white-box SDN switch

FIB caching/prediction

- W. Zhang, J. Bi, J. Wu, and B. Zhang, "**Catching popular prefixes at AS border routers with a prediction based method**," Comput. Networks, vol. 56, no. 4, pp. 1486–1502, Mar. 2012.
  - **GM(1,1)** outperforms LFU, LRU

David Barroso, Spotify
*Building an extensible **SDN** Internet Router with commodity hardware*
https://youtu.be/o1njanXhQqM

Difference from FIB caching/prediction

Basically, a matter of time scales

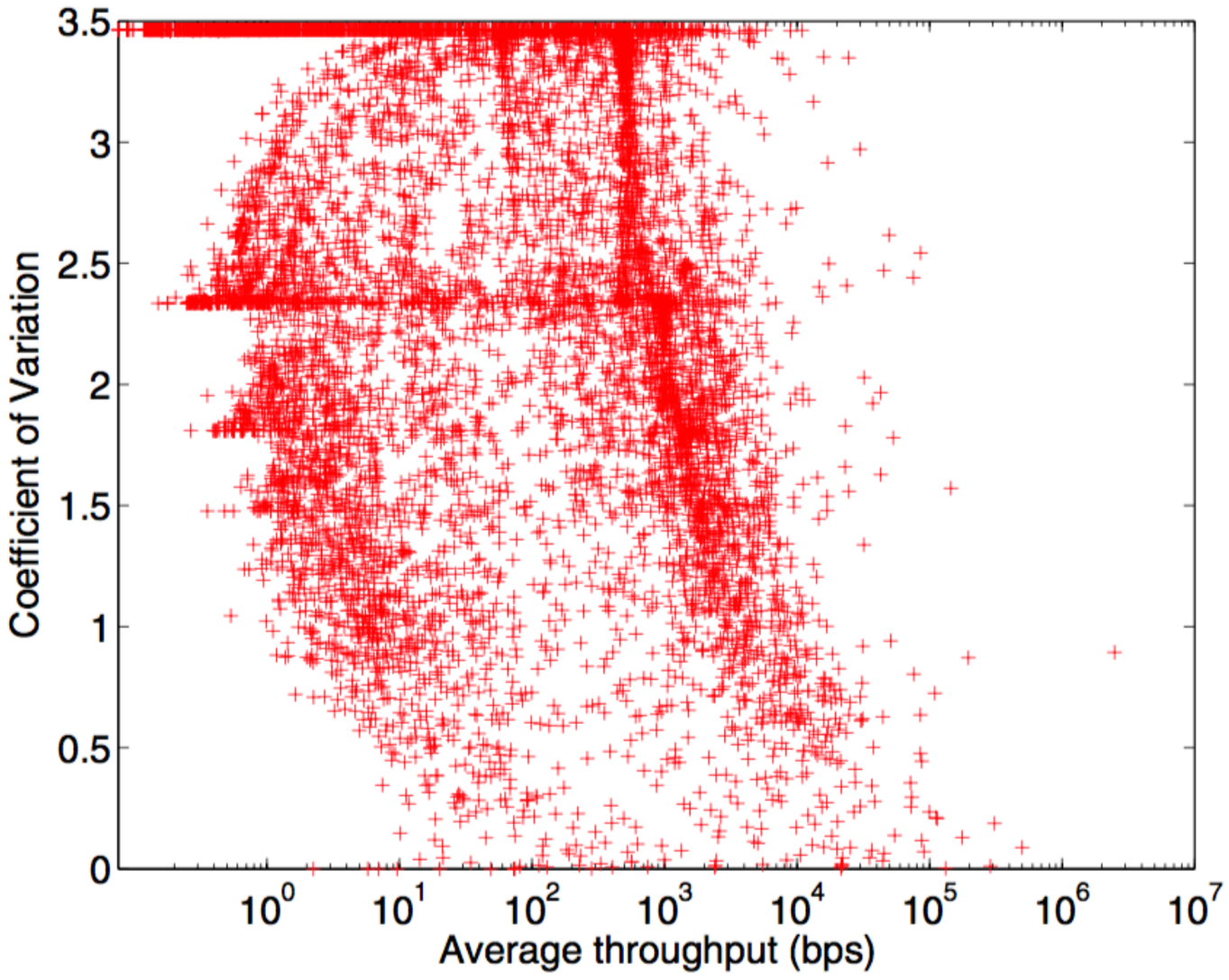
FIB caching, 5min interval or less, memory of hours;
Prefix selection, 1 hour interval, memory of days till weeks.

Why?

Not only data-plane is involved.
Probes discovery, probe selection, long term pattern and bursty-traffic, prefix churn, etc.

And **more** than that….

Traffic dynamism



K. Papagiannaki, N. Taft, and C. Diot, "**Impact of Flow Dynamics on Traffic Engineering Design Principles**," INFOCOM, 2004.

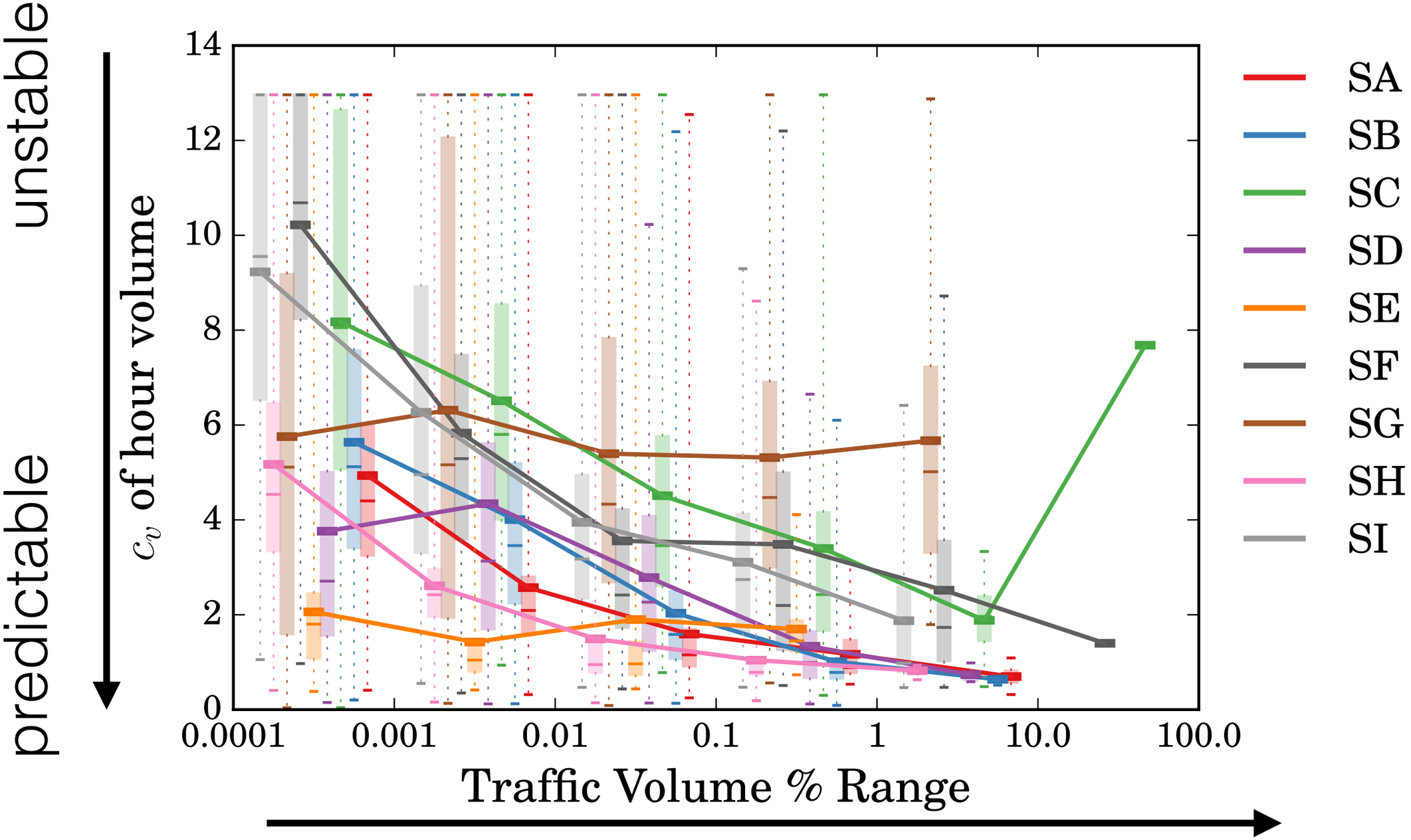
> "… shows that there is no clear correlation between the mean and the coefficient of variation of the bandwidth of a network prefix flow."

What is case for traffic aggregated by BGP prefix over longer interval?

Volume importance vs predictability

Prefix volume share — Cv

Cv = std/mean



prefix of **small** volume share

prefix of **big** volume share

A solution

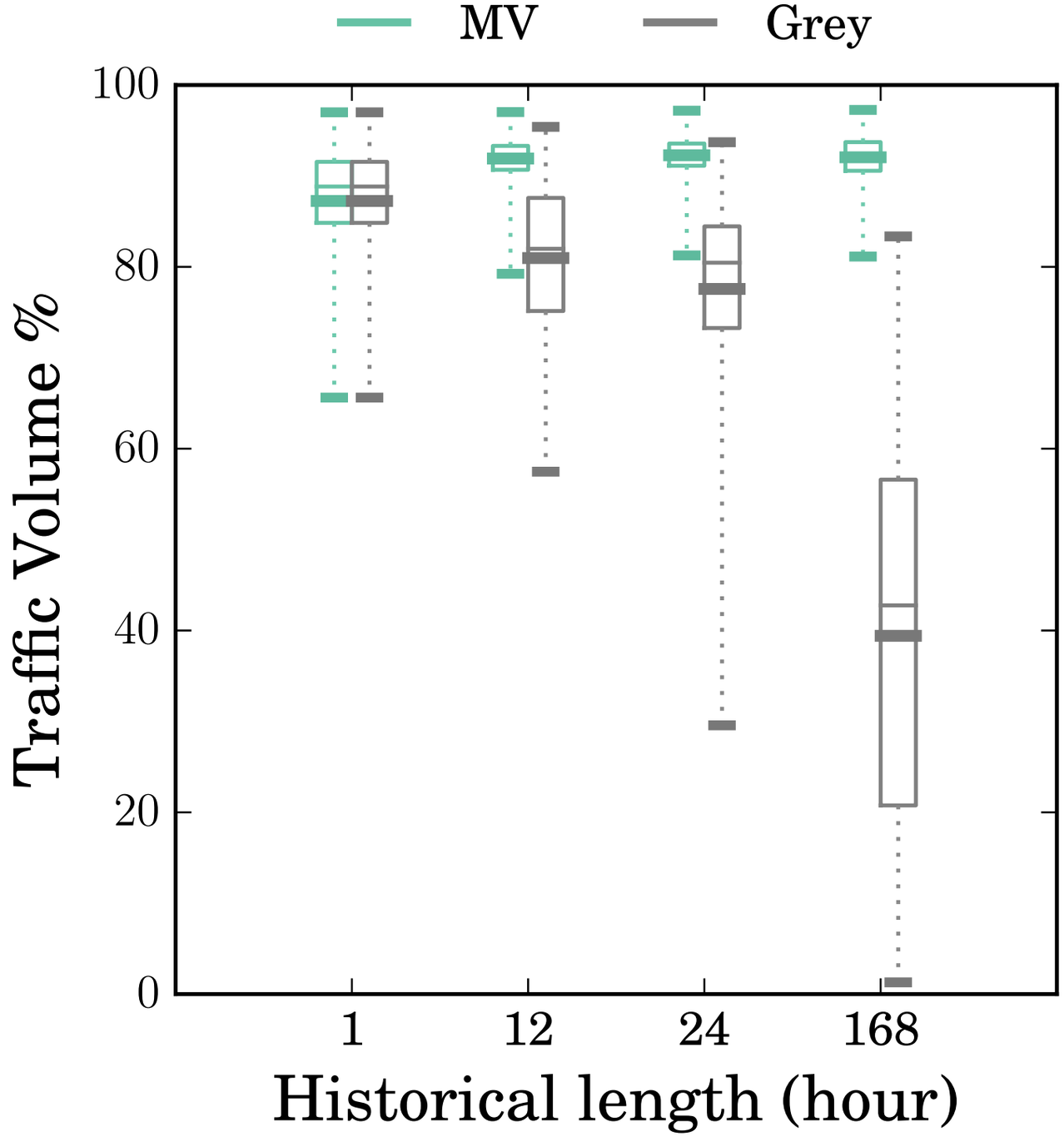
as simple as

Moving Average

Volume coverage of **MV** compared to Grey Model **GM(1,1)**

good coverage

Poor coverage



short memory

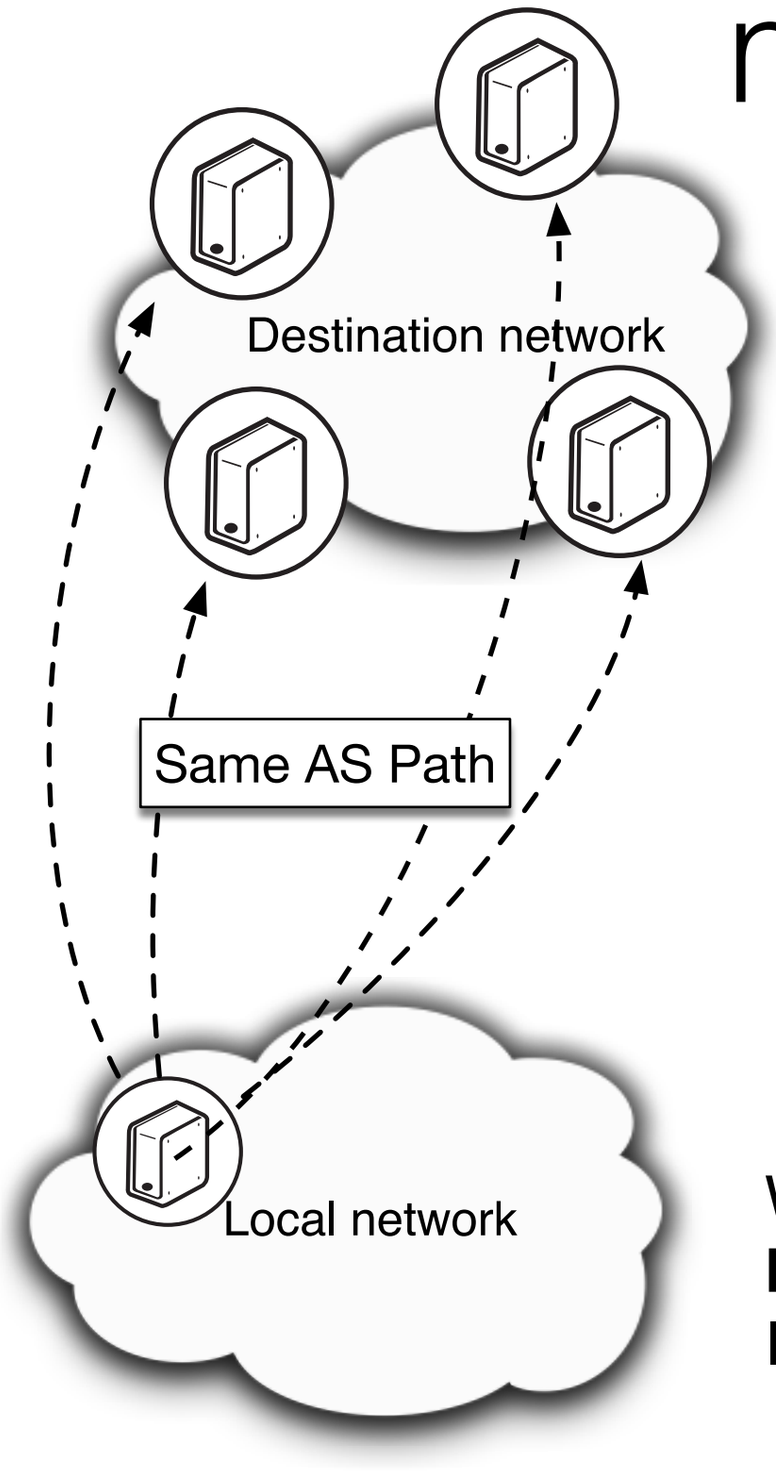
long memory

Summary

- An old scenario with many remaining challenges;
- A possible approach realizing it.

Many other challenges…

Measure a same AS path with multiple probes



Are these measurements different? How?

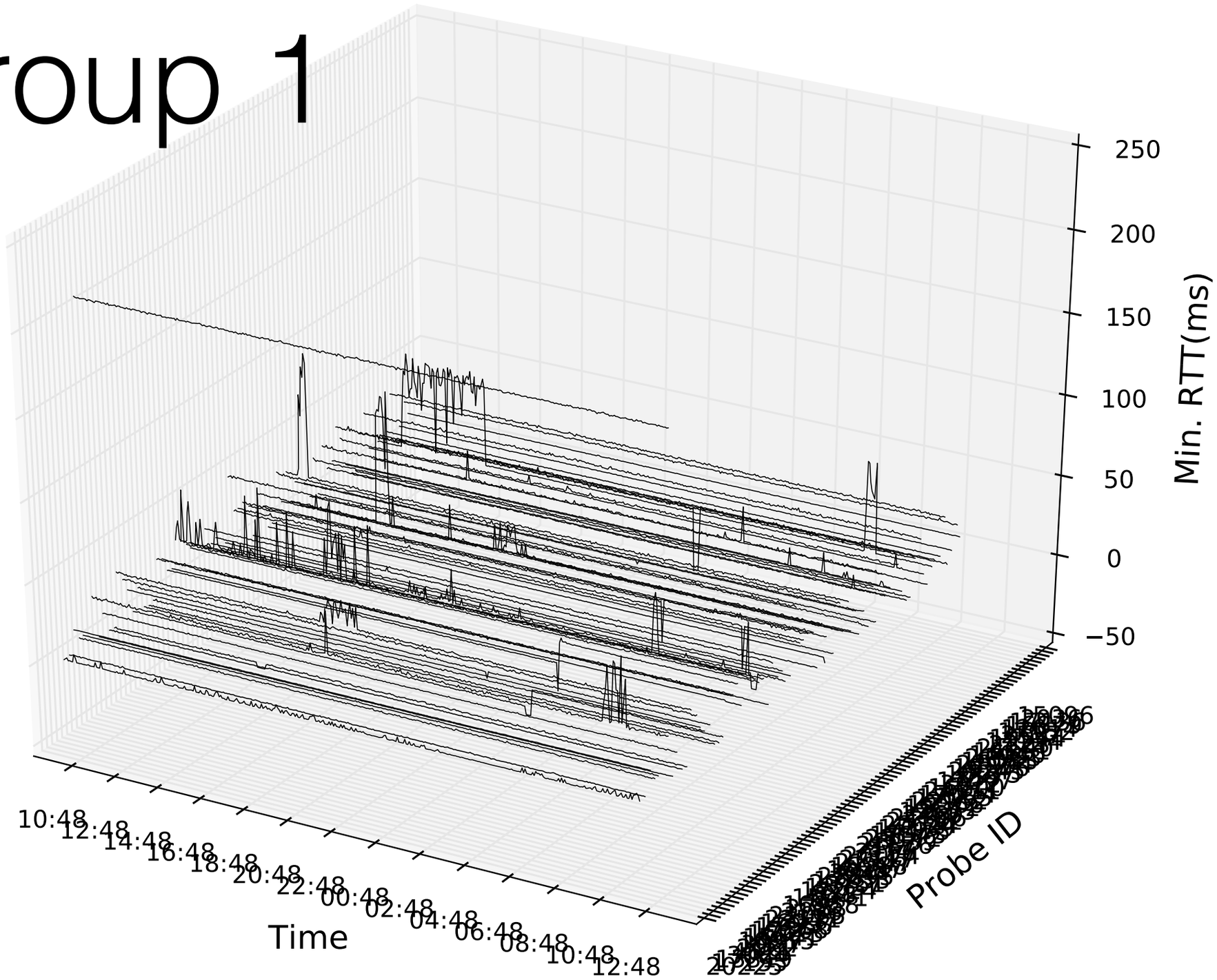
Where does the difference take place?

Which probe should we use?

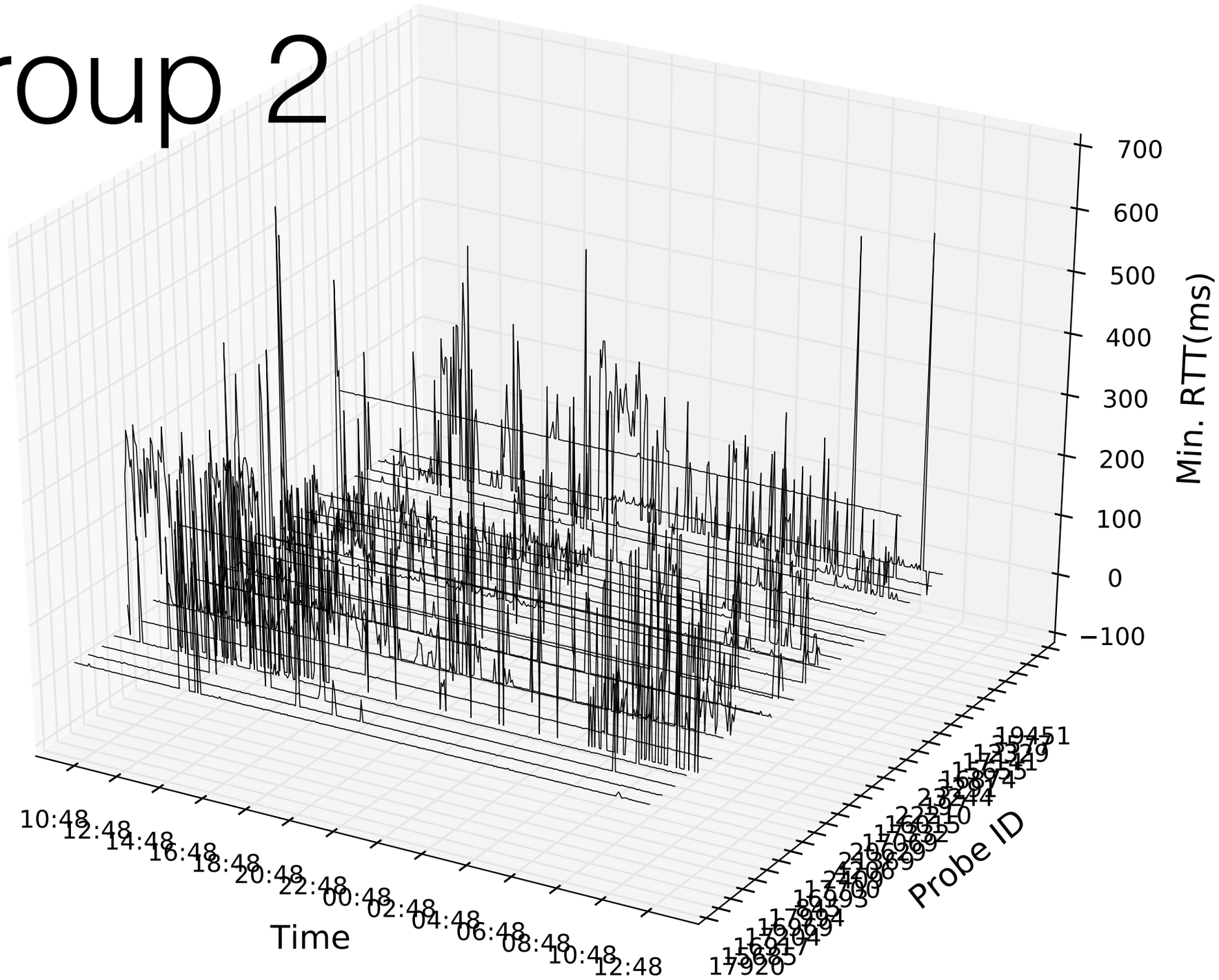
W. Shao, J.L. Rougier, F. Devienne, and M. Viste, “**Improve RTT Measurement Quality via Clustering in Inter-Domain TE** ,” IEEE/IFIP AnNet, 2016.

end2end RTT

Group 1



end2end RTT
Group 2



To measure is to see.

To see is to understand.

Understanding allows automation.

Q&A

Appendix

Not all references are given.
The listed ones could be a good starting point.

https://github.com/WenqinSHAO/NANOG67_prez.git

Appendix-I

Measurement-based Inter-domain TE

- N. Feamster, D. G. Andersen, H. Balakrishnan, and M. F. Kaashoek, “Measuring the effects of internet path faults on reactive routing,” ACM SIGMETRICS Perform. Eval. Rev., vol. 31, no. 1, p. 126, 2003.
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Appendix-II

Traffic volume forecasting

- ARMA family, e.g. ARIMA, SARIMA, FARIMA, O(L^2), pre-process needed for each individual trace, L being historical record length.
- Artificial Neural Network family, e.g. TLFN, O(L*M), M for number of hidden nodes, usually bigger than L.
- Wavelet,O(L), pre-process needed, less accurate than FARIMA and ANN. H. Feng and Y. Shu, “Study on network traffic prediction techniques,” Proceedings. Int. Conf. Wirel. Commun. Netw. Mob. Comput., vol. 2, no. 3, pp. 995-998, 2005.
- Grey model GM(1,1) predicts the accumulated value of a time series, O(L). D. Julong, “Introduction to Grey System Theory,'' J. Grey Syst., vol. 1, pp. 1-24, 1989.

Appendix-III

Traffic dynamism

- K. Papagiannaki, N. Taft, and C. Diot, “Impact of Flow Dynamics on Traffic Engineering Design Principles,” INFOCOM, 2004.
  - no clear correlation between throughput and its stability.

- J. J. Wallerich and A. Feldmann, “Capturing the variability of internet flows across time,” INFOCOM, 2006.
  - throughput ranking of flows can change drastically over time.

- W. Zhang, J. Bi, J. Wu, and B. Zhang, “Catching popular prefixes at AS border routers with a prediction based method,” Comput. Networks, vol. 56, no. 4, pp. 1486–1502, Mar. 2012.
  - assumed positive correlation between popularity and stability.

Appendix-IV

FIB caching

- L. Iannone and O. Bonaventure, “On the cost of caching locator/ID mappings,” CoNEXT, 2007.
- C. Kim, M. Caesar, A. Gerber, and J. Rexford, “Revisiting route caching: The world should be flat,” PAM, 2009.
- H. Ballani, P. Francis, T. Cao, and J. Wang, “Making routers last longer with ViAggre,” NSDI, pp. 453–466, 2009.
- N. Sarrar, S. Uhlig, A. Feldmann, R. Sherwood, and X. Huang, “Leveraging Zipf’s law for traffic offloading,” CCR, vol. 42, no. 1, p. 16, Jan. 2012.
- W. Zhang, J. Bi, J. Wu, and B. Zhang, “Catching popular prefixes at AS border routers with a prediction based method,” Comput. Networks, vol. 56, no. 4, pp. 1486–1502, Mar. 2012.
  - GM(1,1) outperforms LFU, LRU