

# **OCN Experience to Handle the Internet Growth and the Future**

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**NTT Communications, OCN**

# Background

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- Internet traffic / full routes are growing more and more
- One of the most important missions of ISPs
  - to carry the traffic with stability and without any congestion
- Making the backbone robust
- We will talk about:
  - current traffic situation in Japan
  - issues at OCN when designing the backbone network
  - future visions

# Agenda

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## **1. Current situation of Internet traffic in Japan**

## **2. What is OCN?**

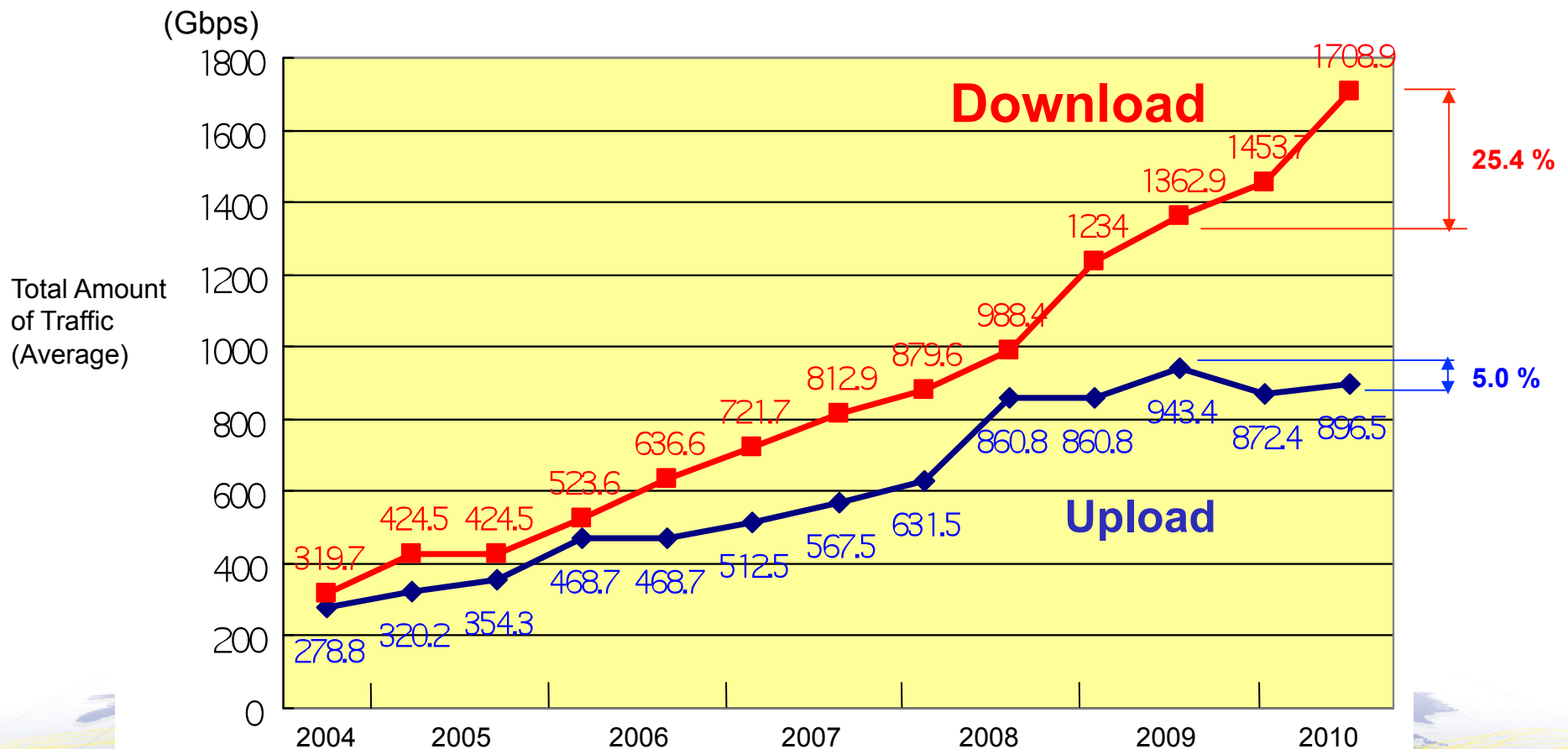
## **3. Current issues we are facing**

- Router Forwarding Table**
- Link Aggregation**

## **4. Future Plan**

# Internet Traffic Trend in Japan

- Total amount of broadband traffic is 1.7Tbps (Download)
  - 25.4% growth compared to last year
- Upload traffic decreased over the last year (896Gbps)



source: Internet Traffic Trends in Japan ( Ministry of Internal Affairs and Communications )  
[http://www.soumu.go.jp/menu\\_news/s-news/01kiban04\\_01000006.html](http://www.soumu.go.jp/menu_news/s-news/01kiban04_01000006.html) (Japanese)

# Internet Traffic Trend in Japan (cont.)

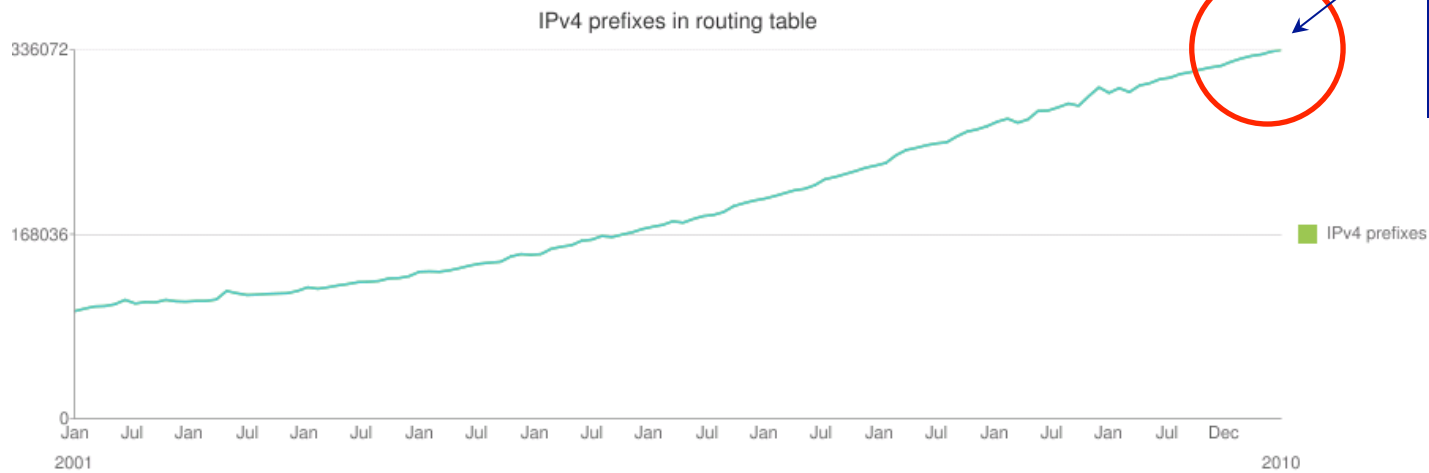
- The number of broadband subscribers and the traffic volume per subscriber are growing



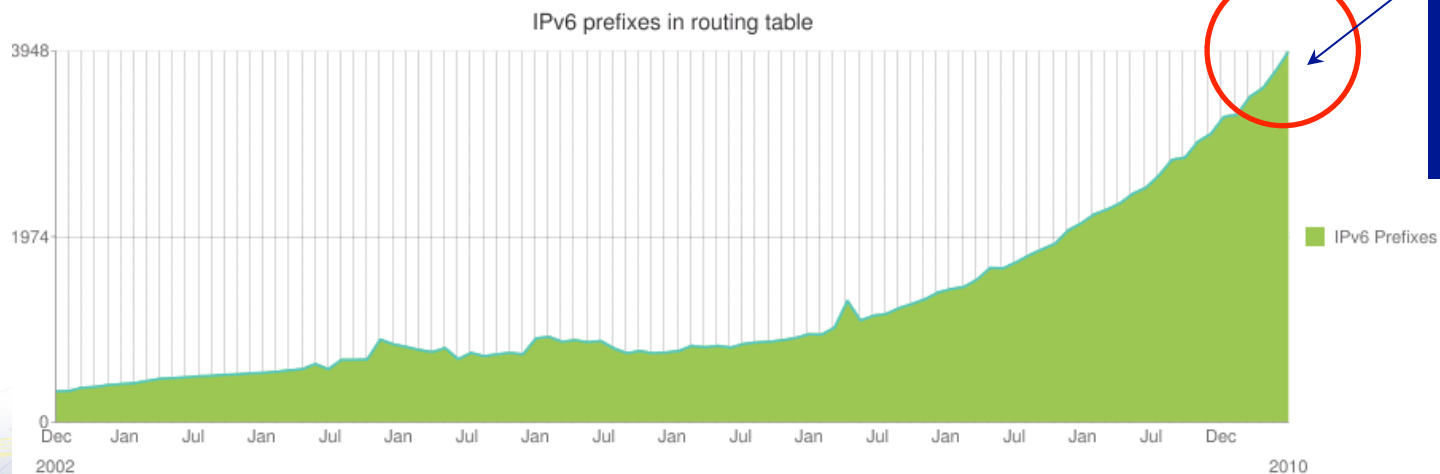
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# Internet Full Routes Trend

- Internet full routes growing



The number of IPv4 prefix : **over 330,000** in June 2011



The number of IPv6 prefix : **over 6,000** in June 2011

# Overview

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- Internet traffic in Japan / full routes have been growing consistently
- Traffic will keep rising in the future
  - ISPs have to ...
    - design a robust backbone network to deal with the situation
- The backbone we have been making
- The bandwidth we have

# Agenda

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**1. Current situation of Internet traffic in Japan**

**2. What is OCN?**

**3. Current issues we are facing**

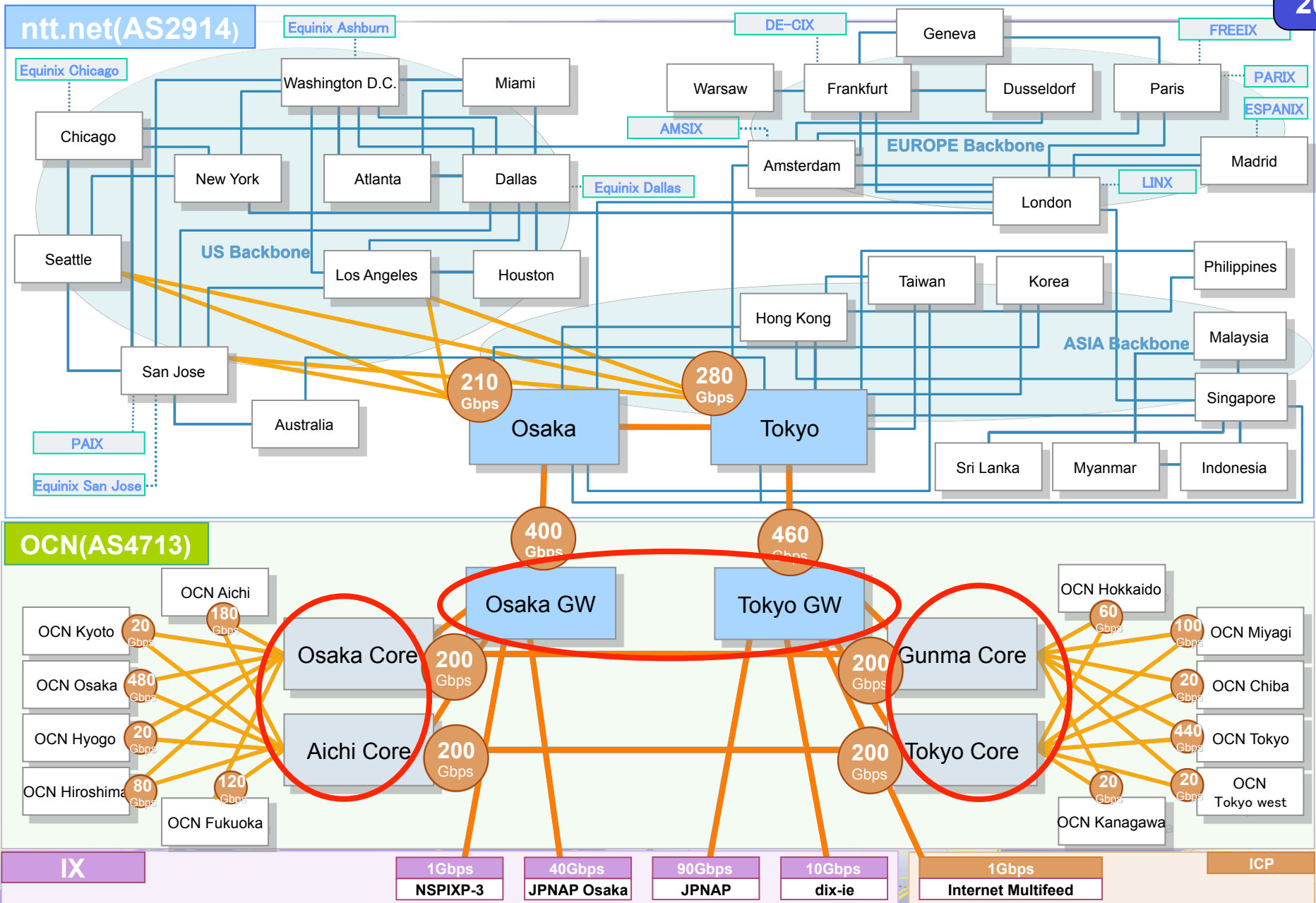
- Router Forwarding Table
- Link Aggregation

**4. Future Plan**



# NTT Communications' IP Backbone Network (ntt.net & OCN)

May 2011



# Network Design Policy of OCN

## *Full redundant network*

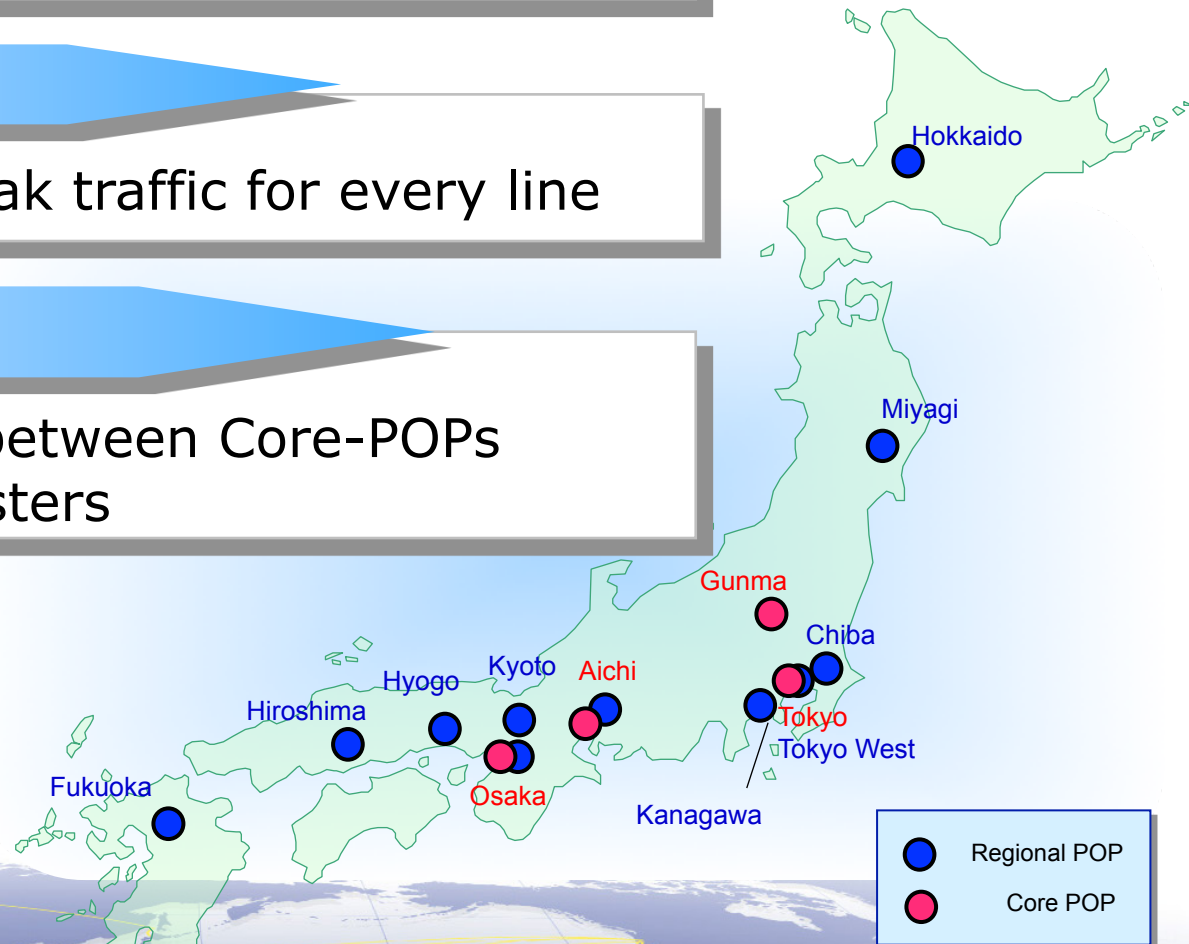
No single point of failure

## *100% Traffic Relief*

Double bandwidth of the peak traffic for every line

## *Disaster Tolerance*

More than 100km distance between Core-POPs minimize impact of the disasters



# Agenda

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1. Current situation of Internet traffic in Japan

2. What is OCN?

**3. Current issues we are facing**

- Router Forwarding Table
- Link Aggregation

4. Future Plan

# The issues we are facing

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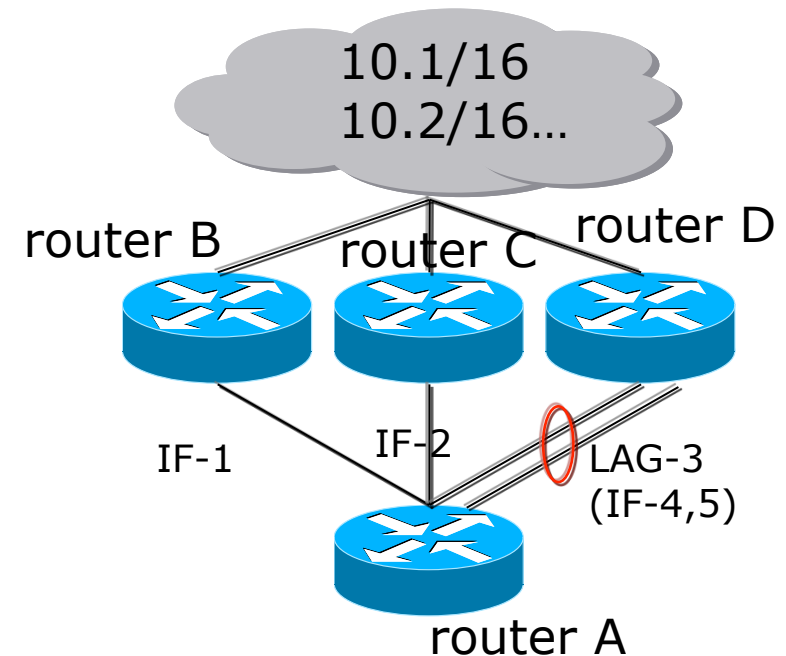
1. Routes Growth  
Scalability of Router Forwarding Tables
2. Traffic Growth  
Link Aggregation

- FIB(Forwarding Information Base) table has been growing
  - 410,000 routes in OCN (June 2011)
- Causes of growing FIB
  1. BGP full routes
  2. Prefixes with no-export (several tens of thousands in OCN)
  3. ECMP, {i, e} bgp-multipath

# Scalability of Router Forwarding Tables

- When a rerouting event occurs, potentially thousands of routes must be updated

FIB of router-A	
prefix	output interface(s)
10.1.0.0/16	IF-1
	IF-2
	LAG-3(IF- <del>4</del> , 5)
10.2.0.0/16	IF-1
	IF-2
	LAG-3(IF- <del>4</del> , 5)

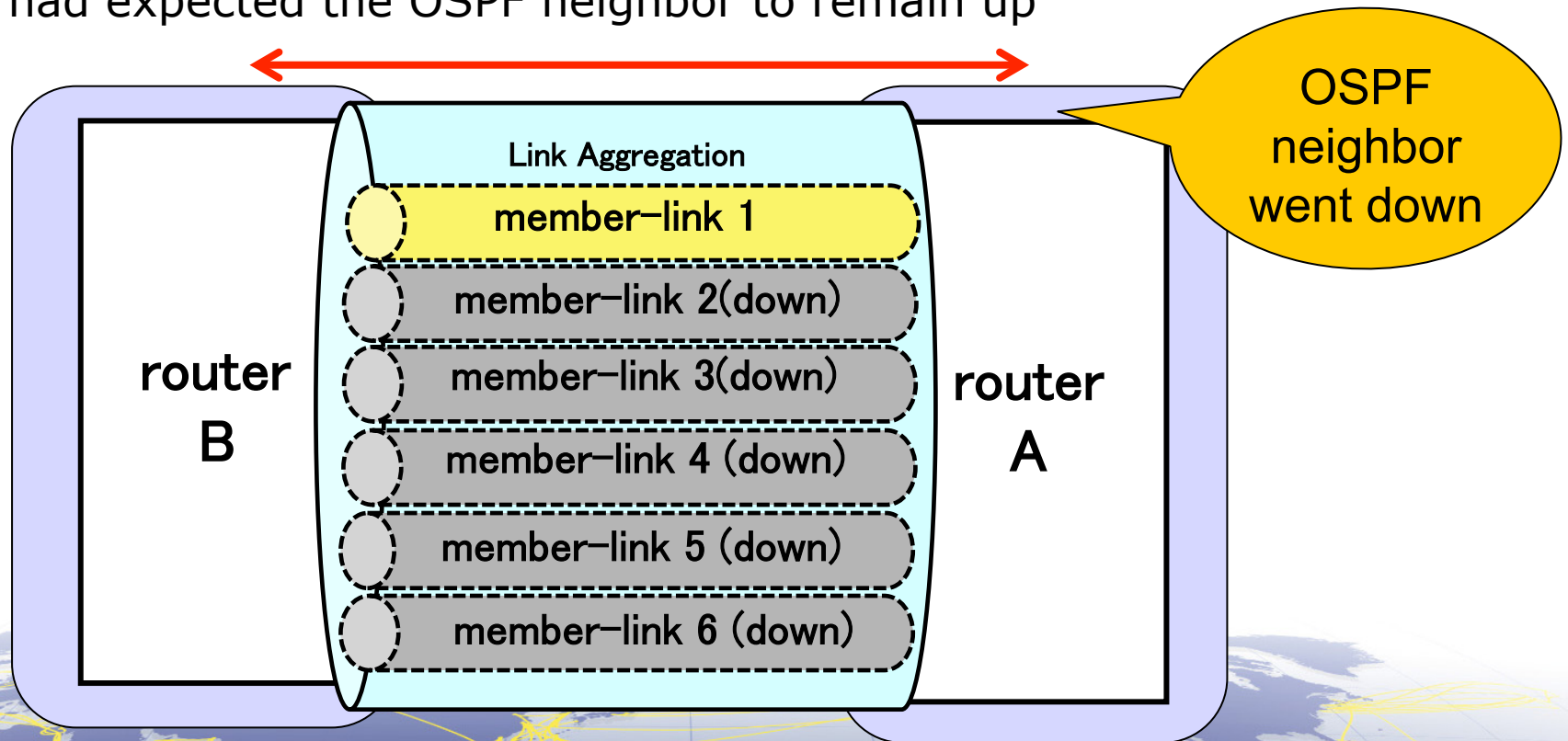


- It took a lot of time to converge the routes
  - When some member-links of a link aggregation were taken down

	FIB table(IPv4)	Convergence time (flattened FIB)
a certain router	360,000	more than 130sec
	500,000	more than 210sec

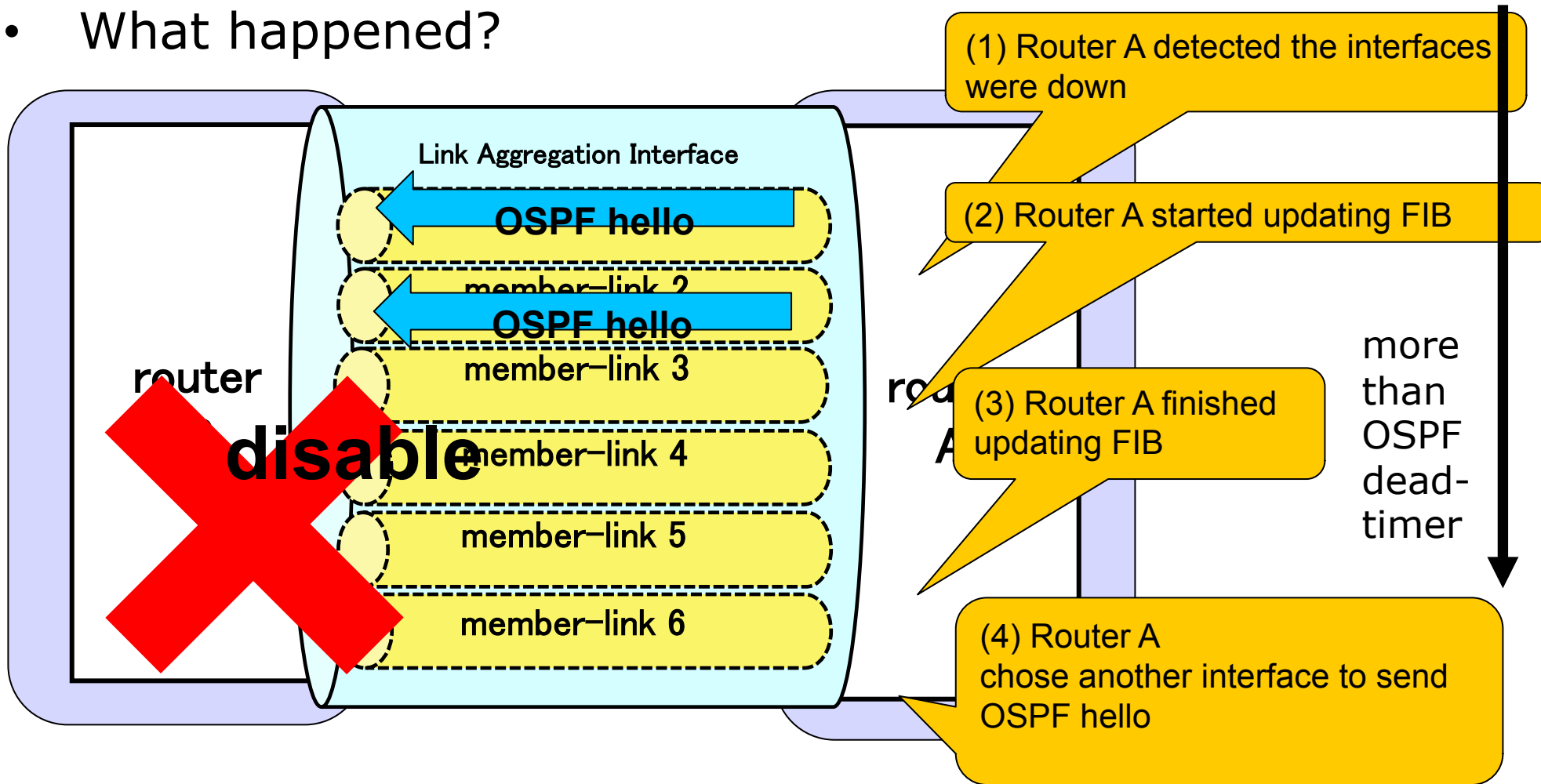
# Scalability of Router Forwarding Tables

- We were facing a problem:
  - OSPF neighbor went down due to FIB table convergence
- Between router A and B
  - Link Aggregation (LAG) had been enabled (minimum-links = 1)
  - OSPF neighbor had been connected through the LAG interface
- When all member-links but one had been disabled
  - We had expected the OSPF neighbor to remain up



# Scalability of Router Forwarding Tables

- What happened?



**Router-A could not send any OSPF hello packets during (1) - (3), then the neighbor went down**



# Scalability of Router Forwarding Tables

- Hierarchical FIB
  - Cisco: BGP Prefix Independent Convergence(PIC)
  - Juniper: indirect-nexthop

For more information: BGP Convergence in much less than a second

<http://www.nanog.org/meetings/nanog40/presentations/ClarenceFilsfils-BGP.pdf>

- Fewer routes to be updated
- Improving the route convergence time

a certain router	FIB table(IPv4)	Convergence time (flattened FIB)	Convergence time (hierarchical FIB)
	500,000	more than 210sec	around 25sec

# Agenda

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1. Current situation of Internet traffic in Japan

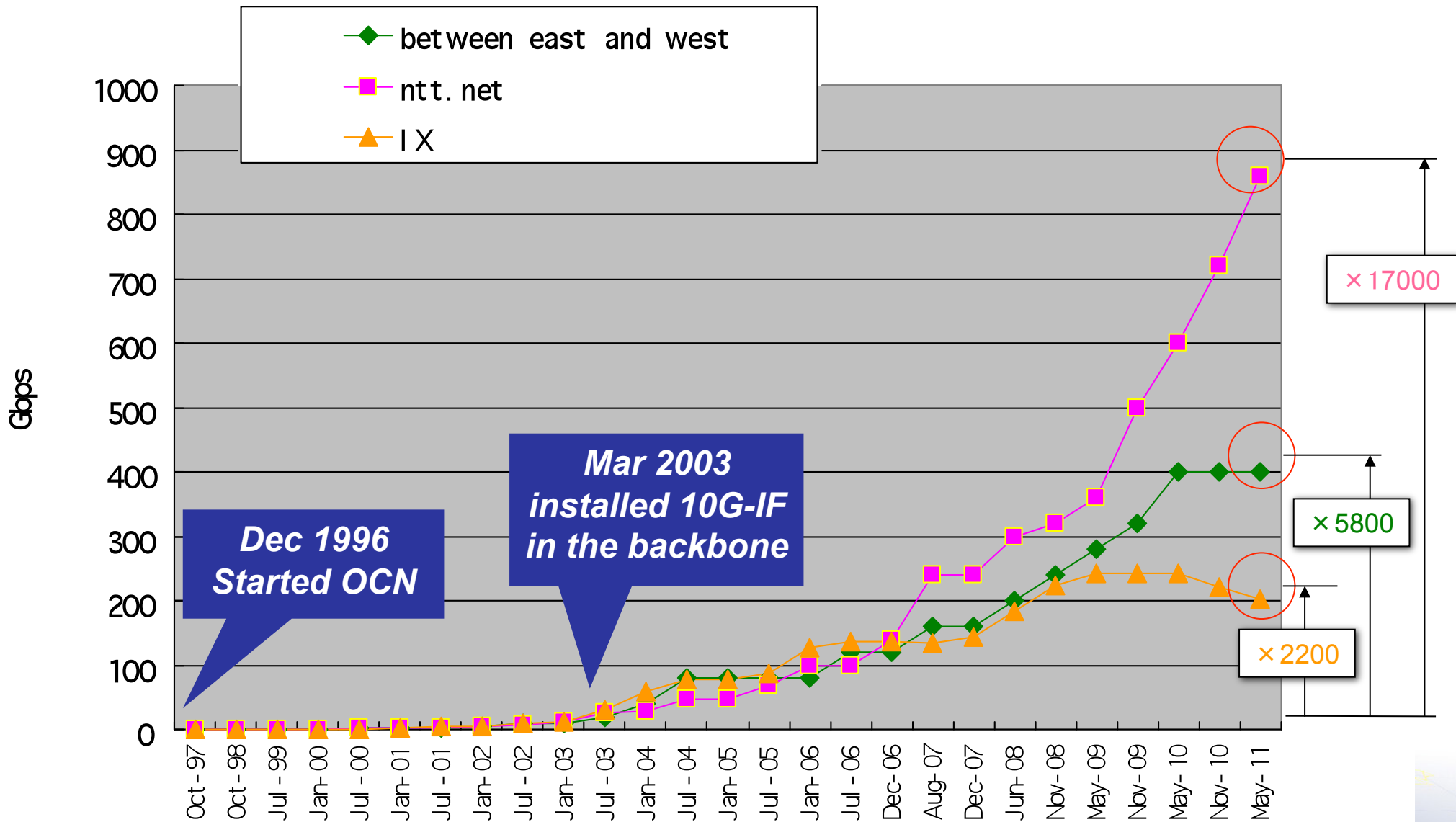
2. What is OCN?

**3. Current issues we are facing**

- Router Forwarding Table
- **Link Aggregation**

4. Future Plan

# Bandwidth History of OCN





# Link Aggregation Issues

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## A) Traffic load-balancing issues (Traffic Polarization)

- Background
  1. Traffic-unbalance by variation of flow - may skip -
  2. Limited number of hash elements
  3. Combination of ECMPs and LAGs
    - Case 1: ECMP and LAG at the same Node
    - Case 2: ECMP and LAG at different Node
    - Case 3: ECMP and ECMP at different Node

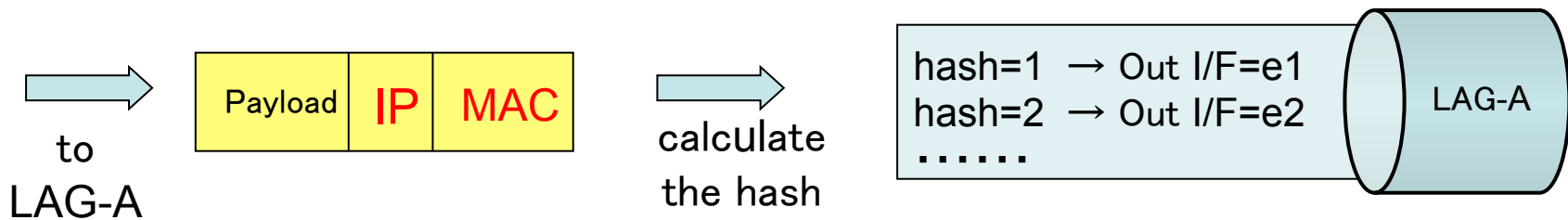
## B) Operational Considerations

1. LACP - might skip -
2. minimum-links - may skip -
3. Ping to each physical interface

## C) Other issues

# A) Traffic load-balancing issues: Background

- Condition of traffic load-balancing method in the LAG
  - **Can't use per-packet round-robin**
    - Simple round-robin bring about packet reordering in a flow
    - Should use flow-based traffic load-balancing method
- **Hash value is used for flow-based traffic load-balance**
- Hashing algorithm: calculate the hash value based on the packet information (IP address, MAC address, and etc.) to decide Output I/F



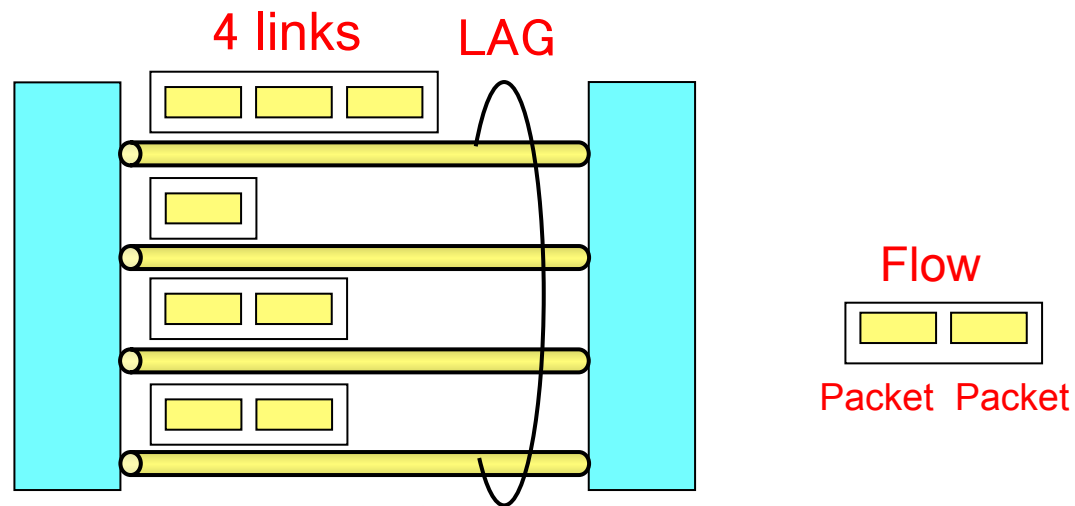
# A) Traffic load-balancing issues

- Issue 1: Traffic-unbalance by variation of flow

- Each flow has each size
- Small issue

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due to limited time**

- Each 10Gbps physical link has a huge number of flows



# A) Traffic load-balancing issues

- Issue 2: Limited number of hash elements
  - Due to this, traffic cannot be evenly distributed
    - Less effective use of bandwidth
  - The less # of hash elements, the worse traffic balance

e.g.: Traffic balance in a LAG when # of hash elements is 8

5 link 10GE LAG	4 link 10GE LAG	3 link LAG
IF#1 H1、H6	IF#1 H1、H5	IF#1 H1、H4、H7
IF#2 H2、H7	IF#2 H2、H6	IF#2 H2、H5、H8
IF#3 H3、H8	IF#3 H3、H7	IF#3 H3、H6
IF#4 H4	IF#4 H4、H8	
IF#5 H5		
2:2:2:1:1	2:2:2:2	3:3:2
10+10+10+10*1/2+10 *1/2 = <u>40</u>	10+10+10+10= <u>40</u>	10+10+10*2/3 = <u>26.7</u>

<- Traffic balance ratio  
<- Effective bandwidth in the LAG

**Use only 40G / 50G**

**Use only 27G / 30G**



# A) Traffic load-balancing issues

- cf. Difference in traffic load-balance by # of hash elements

**e.g.1: Traffic balance in a LAG when # of hash elements is 8**

5 links LAG	4 links LAG	3 links LAG
IF#1 H1, H6	IF#1 H1, H5	IF#1 H1, H4, H7
IF#2 H2, H7	IF#2 H2, H6	IF#2 H2, H5, H8
IF#3 H3, H8	IF#3 H3, H7	IF#3 H3, H6
IF#4 H4	IF#4 H4, H8	
IF#5 H5		
<b><u>40</u></b>	<b><u>40</u></b>	<b><u>26.7</u></b>

The more # of hash elements, the better traffic balance

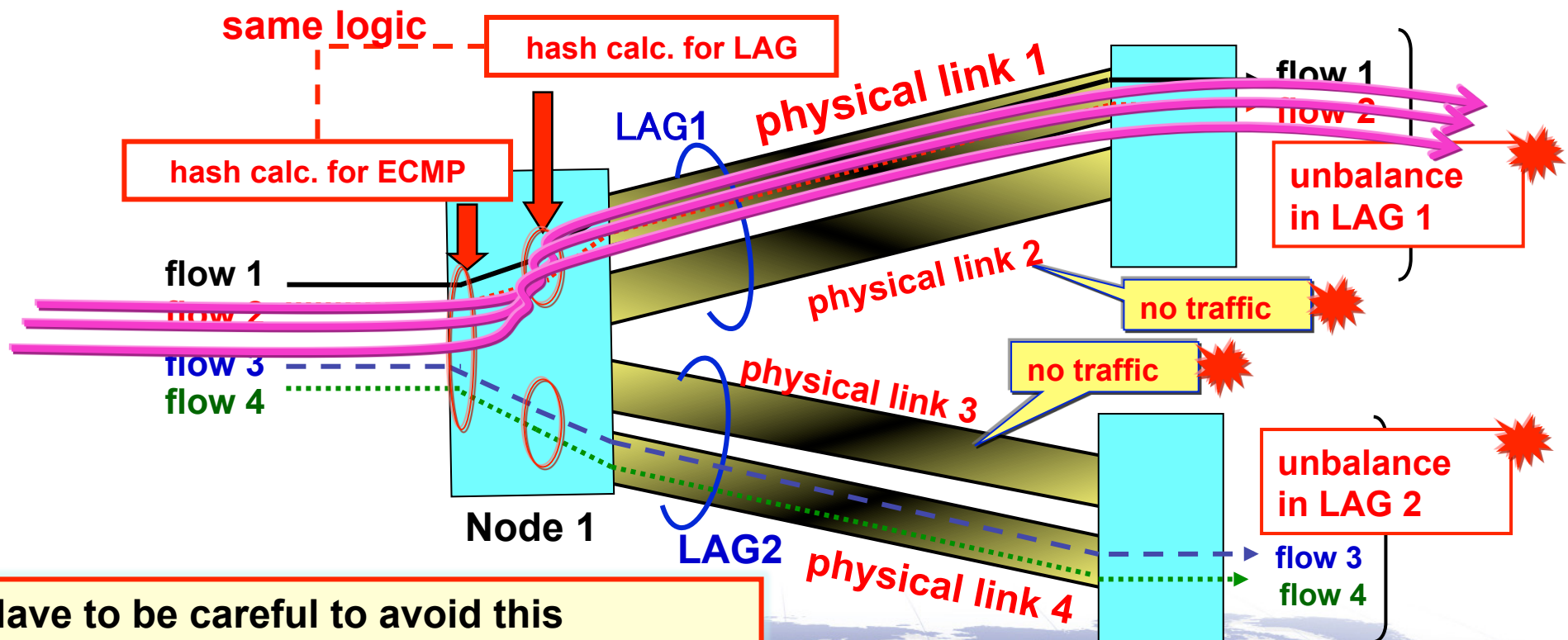
Should avoid odd number of member-links in a LAG

**e.g.2: Traffic balance in a LAG when # of hash elements is 32**

5 links LAG	4 links LAG	3 links LAG
IF#1 H1, H6, ...H26, H31	IF#1 H1, H5, ...H29	IF#1 H1, H4, ...H28, H31
IF#2 H2, H7, ...H27, H32	IF#2 H2, H6, ...H30	IF#2 H2, H5, ...H29, H32
IF#3 H3, H8, ...H28	IF#3 H3, H7, ...H31	IF#3 H3, H6, ...H30
IF#4 H4, H9, ...H29	IF#4 H4, H8, ...H32	
IF#5 H5, H10, ...H30		
7:7:6:6:6	8:8:8:8	11:11:10
$10+10+10*6/7+10*6/7+10*6/7 = \mathbf{45.7}$	$10+10+10+10 = \mathbf{40}$	$10+10+10*10/11 = \mathbf{29.1}$

# A) Traffic load-balancing issues

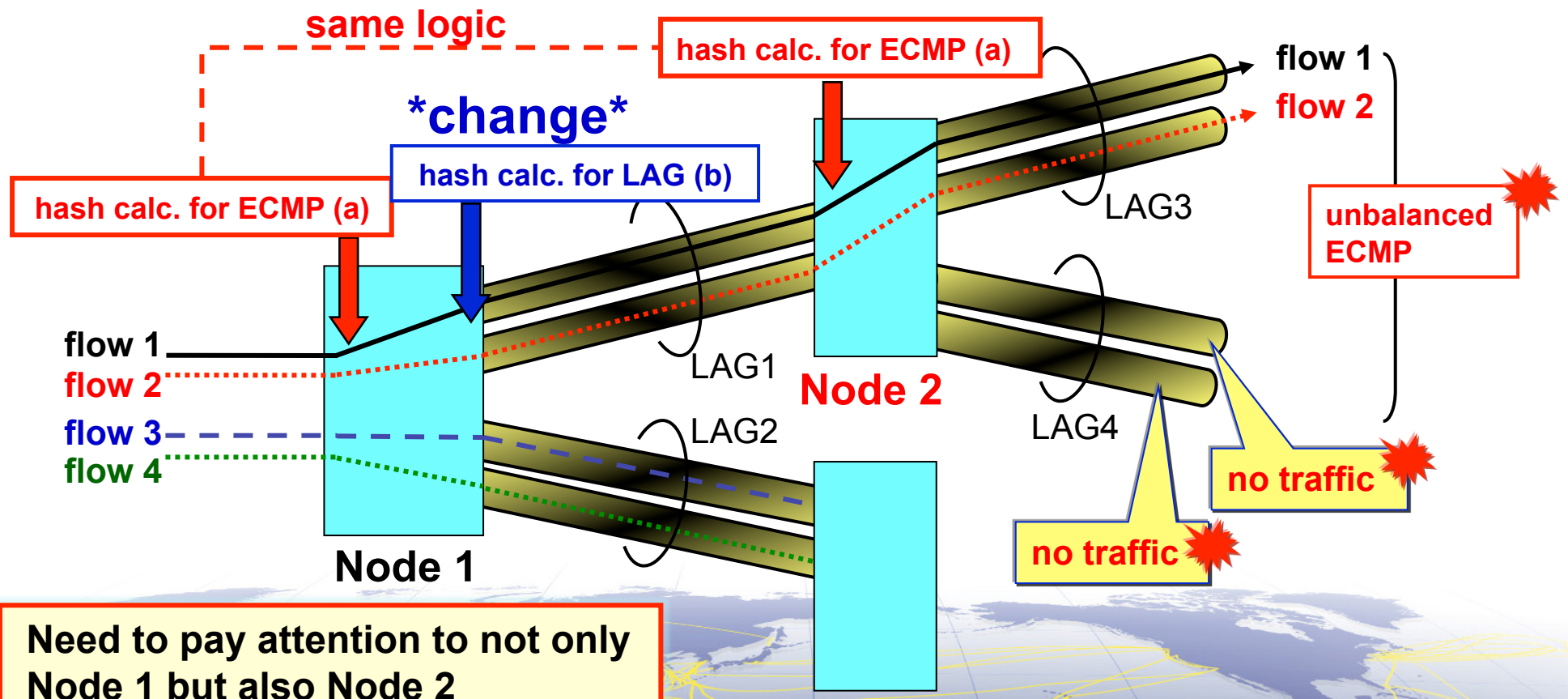
- Issue 3: Combination of ECMPs (Equal Cost Multi Path) and LAGs
- Case 1:
  - When hash calculation logic of LAG is the same as ECMP's, it will bring about unbalanced traffic in physical links



- Have to be careful to avoid this
- Some routers have the same calculation logics for ECMP and LAG as a default

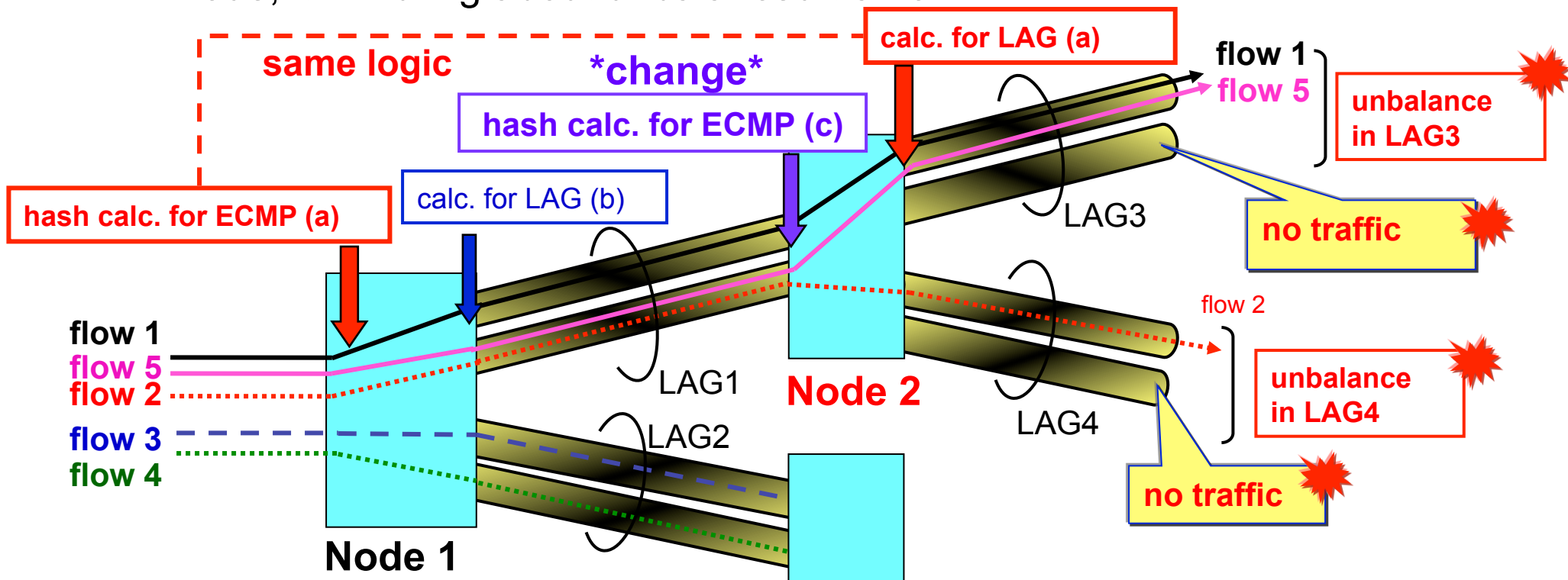
# A) Traffic load-balancing issues

- Issue 3: Combination of ECMPs and LAGs
- Case 2:
  - When calculation logic of ECMP is the same as that of next node, it will bring about unbalanced traffic



# A) Traffic load-balancing issues

- Issue 3: Combination of ECMPs and LAGs
- Case 3:
  - When calculation logic of ECMP is the same as that of LAG at the next node, it will bring about unbalanced traffic

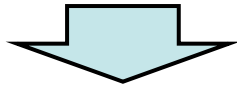


\* Some latest routers can include a router-ID in the seed of hash to avoid case 2,3

**Need to consider balance logics, network topology, configurations**

## B) Operational Considerations

- Consideration 1:
  - In the case of silent-failure, traffic through the fault link will drop

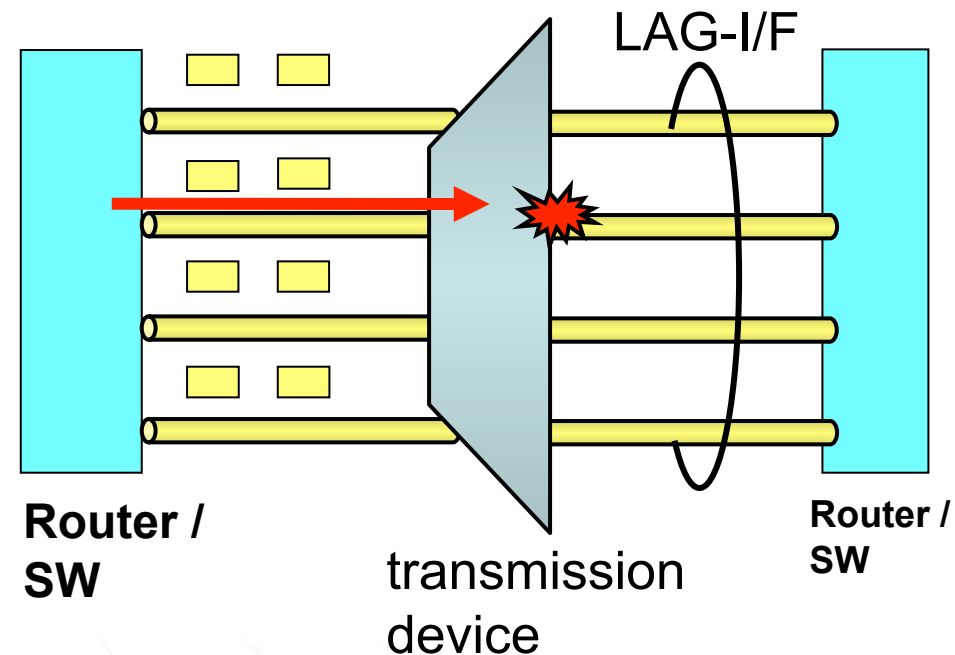


### LACP (Link Aggregation Control Protocol)

- Send and receive control frames in physical links
- Attention to detail Interoperability
  - Basically good
  - Different default mode (fast / slow)
  - Different reaction to null ID (bug) LACP (keep down / once down then go up)

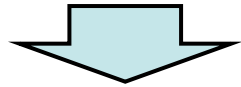
BFD Per Member Link  
(Bidirectional Forwarding Detection)

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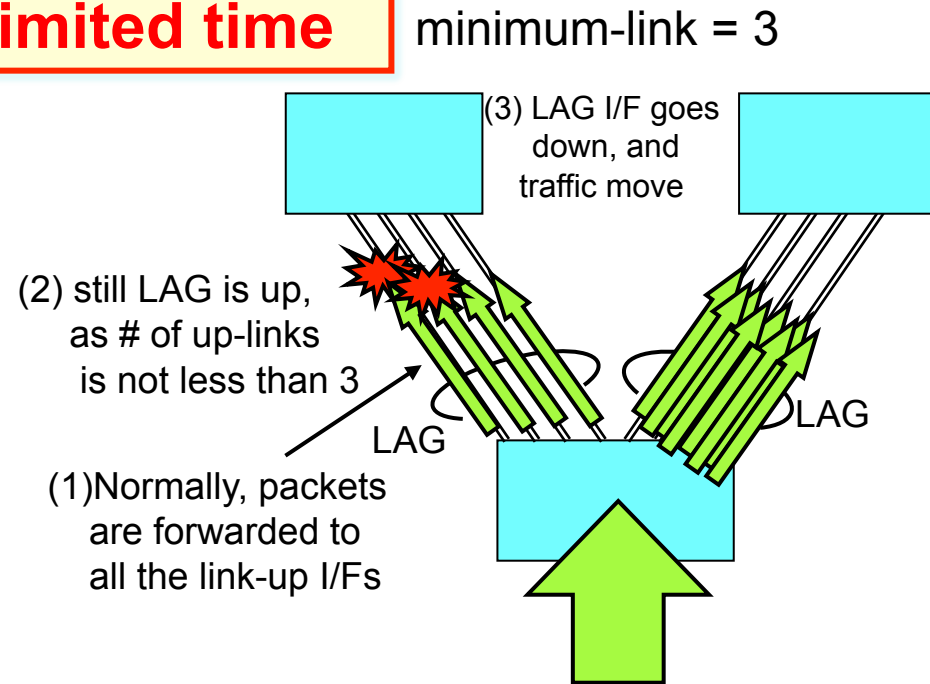
## B) Operational Considerations

- Consideration 2:
  - Switching policy of LAG-I/F
    - minimum-link (trunk-threshold)
    - threshold whether LAG-I/F is up or down



- This switching policy is important for effective use of LAG
- **should consider the entire network topology to use minimum-links**

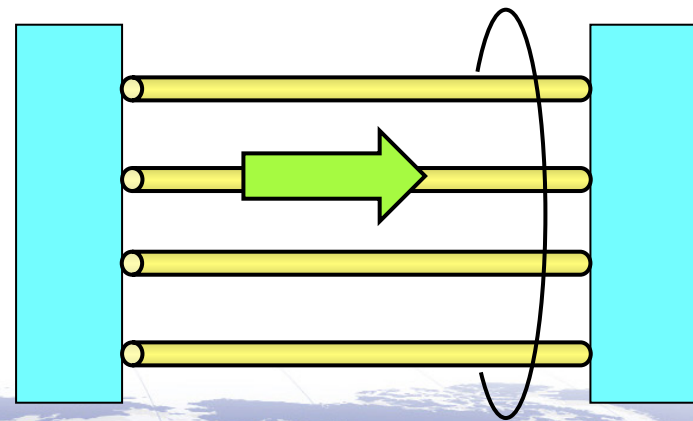
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# of links in LAG	3	4	5	6	7	8	9	10
minimum-link	3	3	4	5	5	6	7	7

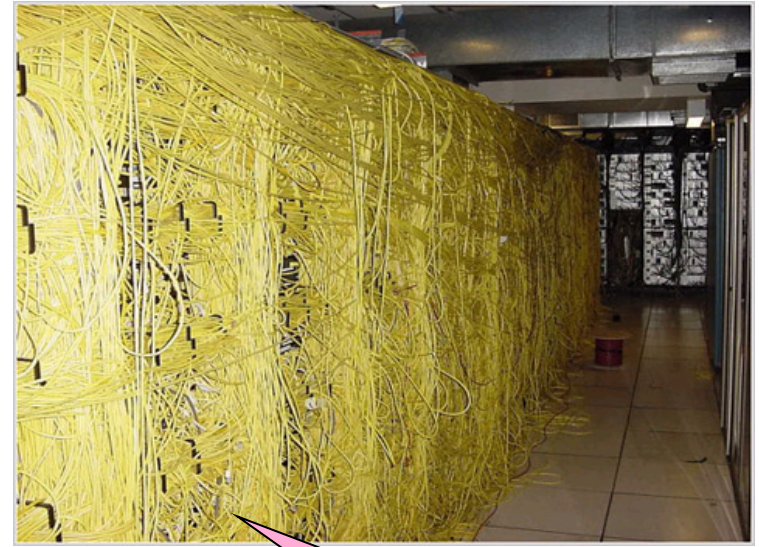
## B) Operational Considerations

- Consideration 3:
  - Ping for test
    - Packet goes through only one physical interface
    - Need to test each interface with letting the rest go down
    - Some recent routers and switches support Ethernet OAM to avoid this troublesome job



## C) Other Issues

- Limitations on # of links in a LAG
- Issues of physical wiring
  - Increased # of physical links
  - > Complicated maintenance
- Need a well-thought-out plan for LAG
  - How to assign physical links to Line Cards
    - Redundant policy
    - MTBF for each part
    - Cost, etc.
  - e.g. Policy 1: keep LAG-I/F up as much as possible
    - assign each physical link to each LC, minimum-link = 1
  - e.g. Policy 2: Switching traffic to the other LAG immediately
    - assign all physical links to one LC, minimum-link = # of links
  - e.g. Policy 3: Between policy 1 and policy 2



NOTE: this is NOT NTT Communications' equipment

- **LAG is troublesome**
  - many LAGs, many member-links



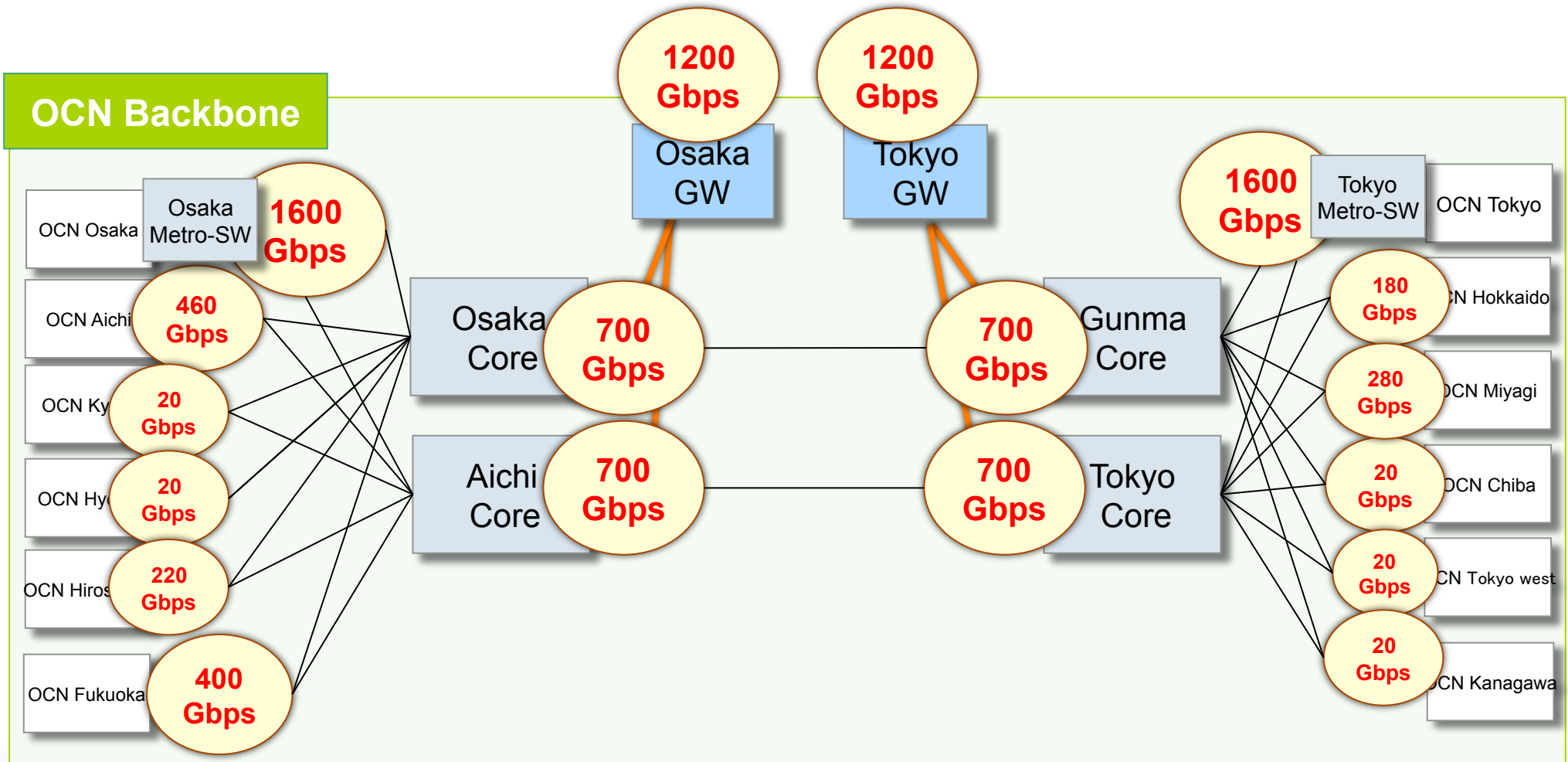
# Agenda

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# OCN future plan

- More bandwidth

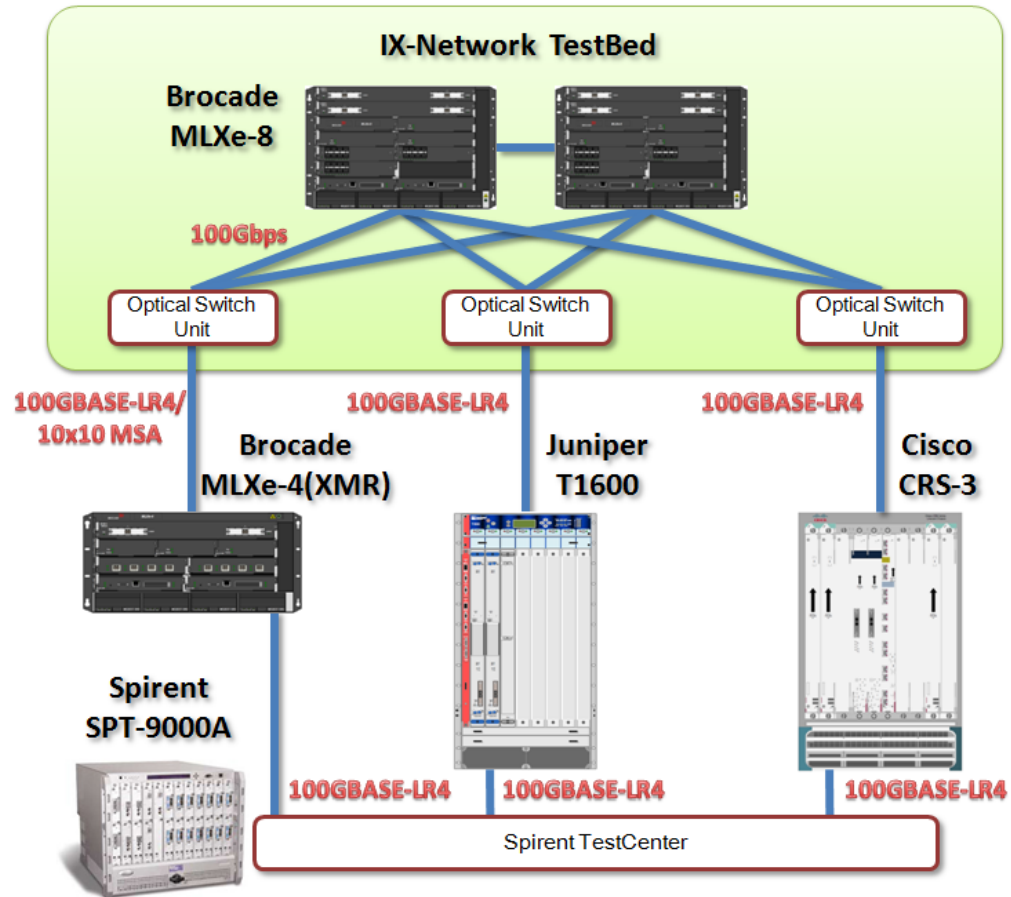


# Expectation for 100GE

- Need 100GE I/Fs
  - Bandwidth over 1Tbps
  - LAG is troublesome
- Request
  - Lower price
    - CFP is expensive
    - 10 x10 MSA (LR10)
  - Long-distance transmission (ER4)
  - Higher Capacity
    - Capacity per chassis will be decreased when migrating from 10GEs to 100GEs in some current routers
  - LAG of 10GE and 100GE simultaneously
  - good Interoperability, easy-operation 100GE LAG, convenient Ether OAM
  - Next step: 400GE, 1T Ether

# 100GE Joint Interoperability Test at JPNAP

- Brocade, Cisco, Juniper and Toyo Corporation (Spirent)
- JPNAP, IIJ, and NTT Communications
- Success of 100 Gigabit Ethernet joint interoperability test at IX
- Confirmed the interoperability between different vendors' products especially at an IX environment
  - Good interoperability
  - Some small issues with each product
    - feedback to vendors with some requests
- Further information is available at:  
<http://www.mfeed.co.jp/english/press/2011/20110601-e.html>



# Summary

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- The traffic in Japan and BGP table has been consistently growing
- We need to consider growth of both routes and traffic to keep our backbone stable
- LAG is troublesome
- We need 100GE to deal with the traffic growth

# Thank you!

