

Shortest Path Bridging IEEE 802.1aq Tutorial and Demo

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Abstract

802.1aq Shortest Path Bridging is being standardized by the IEEE as an evolution of the various spanning tree protocols. 802.1aq allows for true shortest path routing, multiple equal cost paths, much larger layer 2 topologies, faster convergence, vastly improved use of the mesh topology, single point provisioning for logical membership (E-LINE/E-LAN/E-TREE etc), abstraction of attached device MAC addresses from the transit devices, head end and/or transit multicast replication , all while supporting the full suit of 802.1 OA&M.

Applications consist of STP replacement, Data Center L2 fabric control, L2 Internet Distributed Exchange point fabric control, small to medium sized Metro Ethernet control planes. L2 wireless network backhaul....

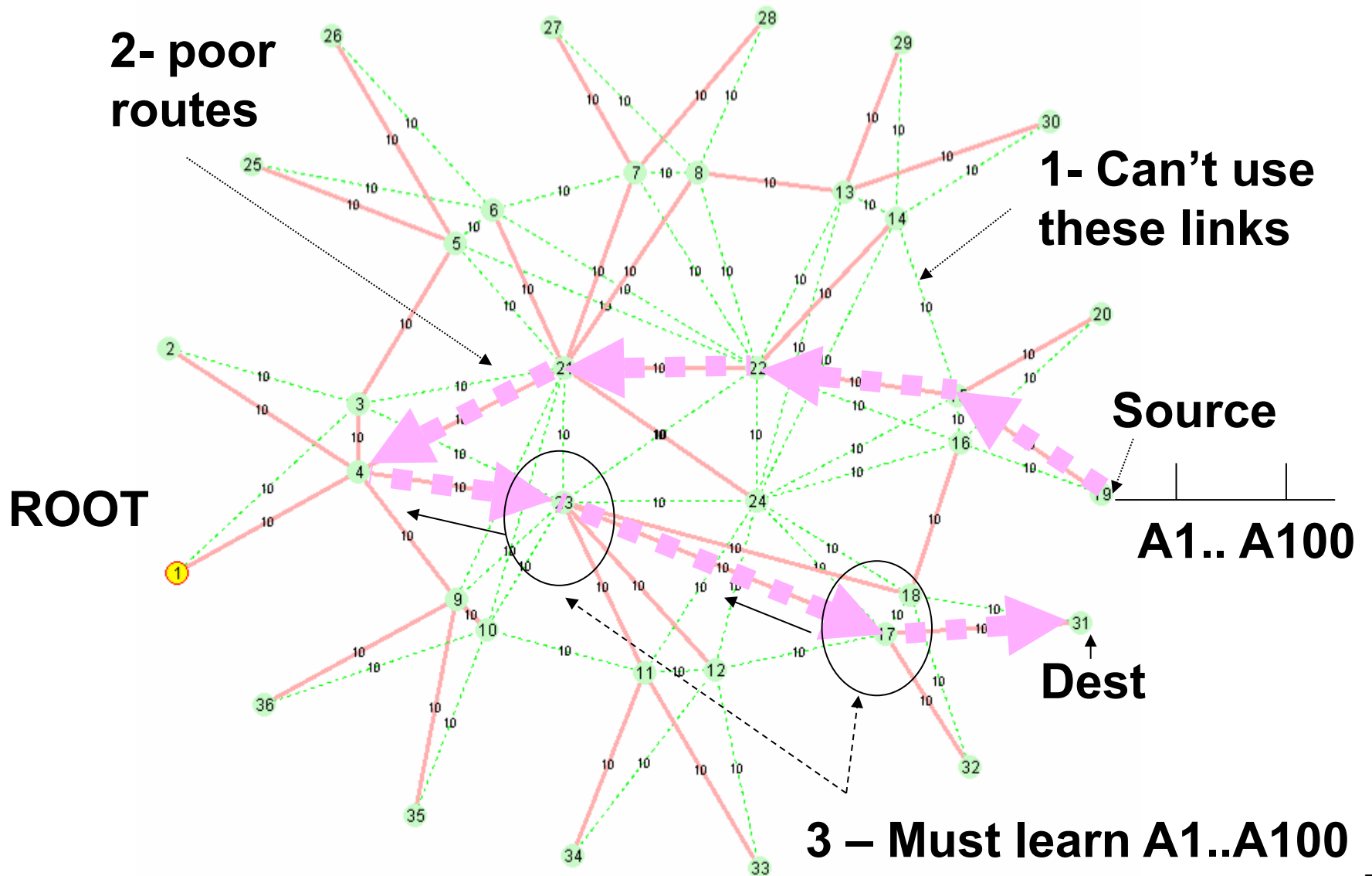
Outline

- **Challenges**
- What is 802.1aq/SPB
- Applications
- How does it work
- Example with hopefully live demo from my lab (assuming connectivity available)

Challenges

- L2 networks that scale to ~1000 bridges.
- Use of arbitrary mesh topologies.
- Use of (multiple) shortest paths.
- Efficient broadcast/multicast routing and replication points.
- Avoid address learning by tandem devices.
- Get recovery times into 100's of millisecond range for larger topologies.
- Good scaling without loops.
- Allow creation of very many logical L2 topologies (subnets) of arbitrary span.
- Maintain all L2 properties within the logical L2 topologies (transparency, ordering, symmetry, congruence, shortest path etc).
- Reuse all existing Ethernet OA&M 802.1ag/Y.1731

Example STP 36 nodes



Outline

- Challenges
- **What is 802.1aq/SPB**
- Applications
- How does it work
- Short Demo (remote switches if possible)

What is 802.1aq/SPB

- **IEEE protocol builds on 802.1 standards**
- **A new control plane for Q-in-Q and M-in-M**
 - Leverage existing inexpensive ASICs
 - Q-in-Q mode called SPBV
 - M-in-M mode called SPBM
- **Backward compatible to 802.1**
 - 802.1ag, Y.1731, Data Center Bridging suite
- **Multiple loop free shortest paths routing**
 - Excellent use of mesh connectivity
 - Currently 16, path to 1000's including hashed per hop.
- **Optimum multicast**
 - head end or tandem replication

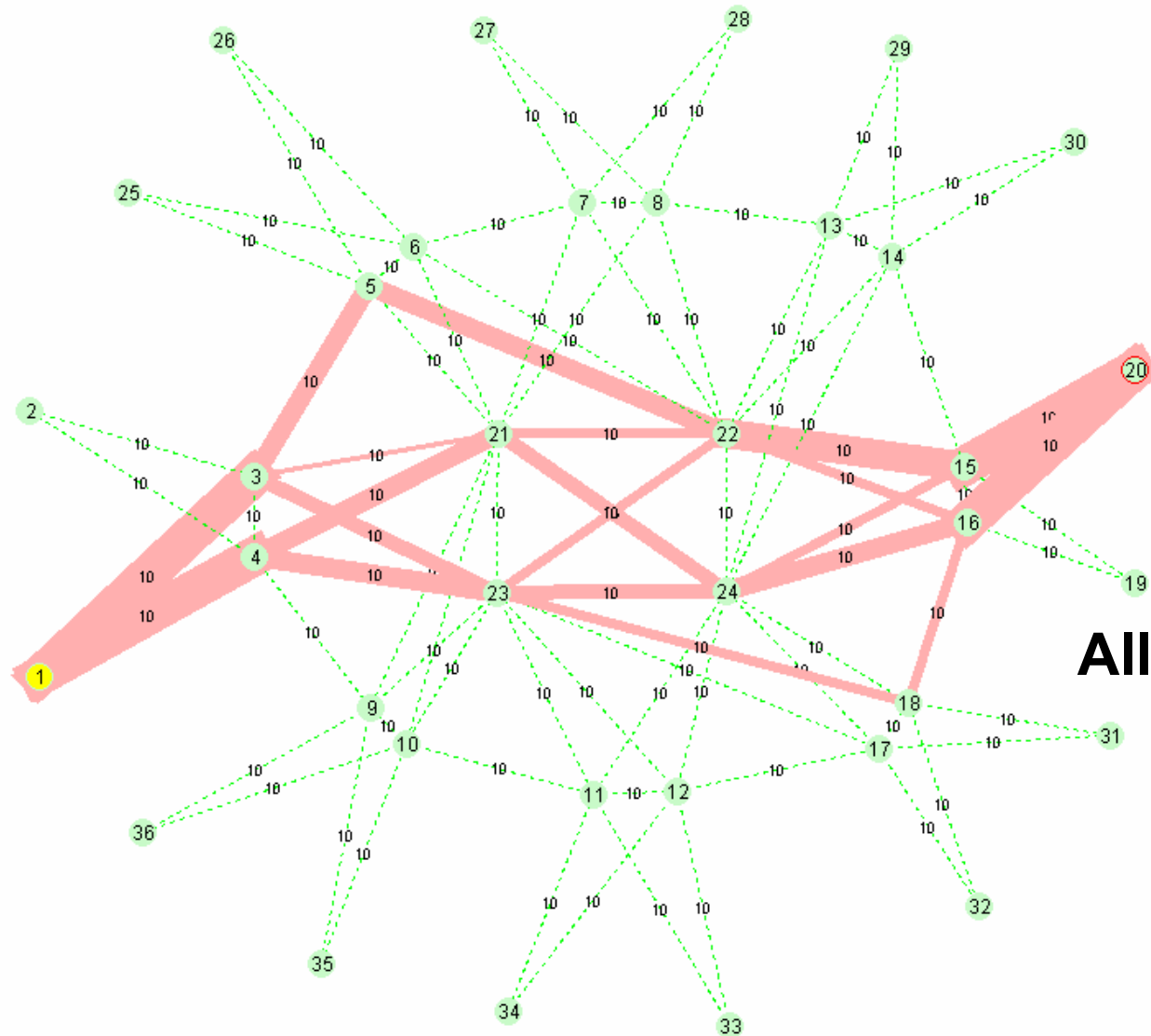
What is 802.1aq/SPB (cont'd)

- **Light weight form of traffic engineering**
 - Head end assignment of traffic to 16 shortest paths.
 - Deterministic routing - offline tools predict exact routes.
- **Scales to ~1000 or so devices**
 - Uses IS-IS already proven well beyond 1000.
 - Huge improvement over the STP scales.
- **Good convergence with minimal fuss**
 - sub second (modern processor, well designed)
 - below 100ms (use of hardware multicast for updates)
 - Includes multicast flow when replication point dies.
Pre-standard seeing 300ms recovery @ ~50 nodes.
- **IS-IS**
 - Operate as independent IS-IS instance, or within IS-IS/IP, supports Multi Topology to allow multiple instances efficiently.

What is 802.1aq/SPB (cont'd)

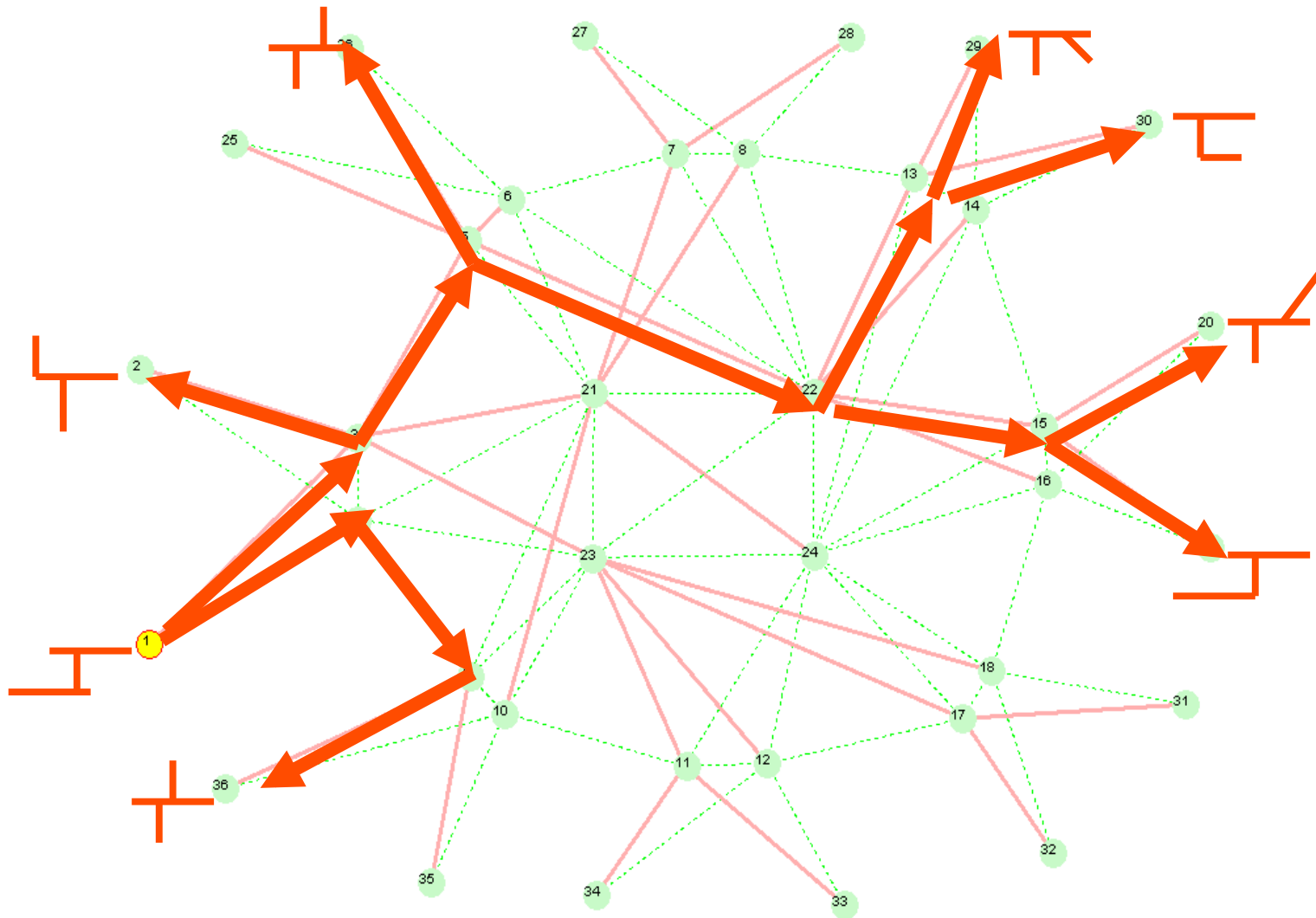
- **Membership advertised in same protocol as topology.**
 - Minimizes complexity, near plug-and-play
 - Support E-LINE/E-LAN/E-TREE
 - All just variations on membership attributes.
- **Address learning restricted to edge (M-in-M)**
 - FDB is computed and populated just like a router.
 - Unicast and Multicast handled at same time.
 - Nodal or Card/Port addressing for dual homing.
- **Computations guarantee ucast/mcast...**
 - Symmetry (same in both directions)
 - Congruence (unicast/multicast follow same route)
 - Tune-ability (currently 16 equal costs paths – opaque allows more)

End result - Visually



Multiple Shortest Path routing + Ethernet OA&M

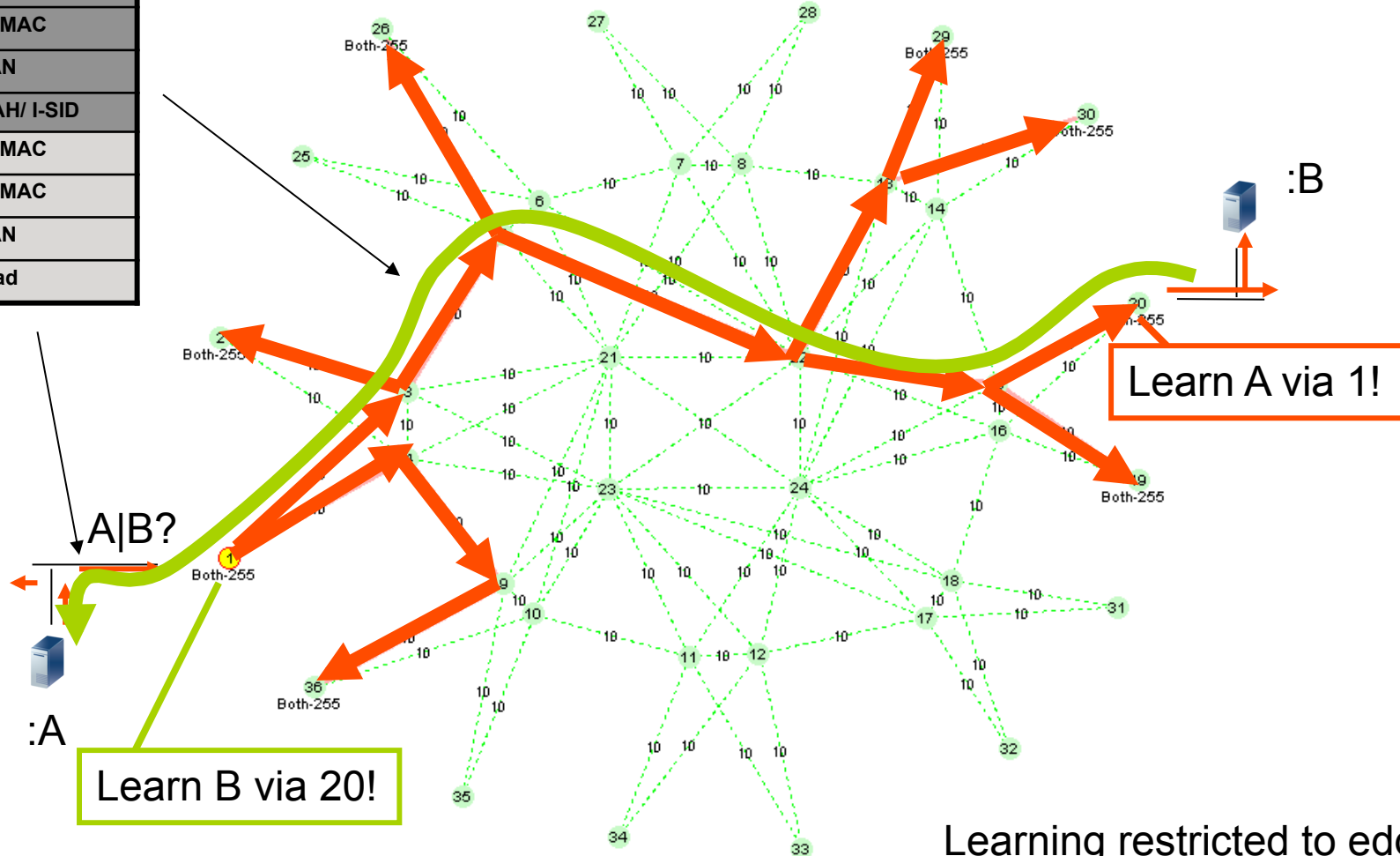
SPF trees form multicast template



Shortest Path First Tree becomes template for multicast tree and is pruned automatically to proper membership.

Edge Learning - Visually

Dst.B-MAC
Src.B-MAC
B-VLAN
801.1AH/ I-SID
Dst.C-MAC
Src.C-MAC
C-VLAN
Payload

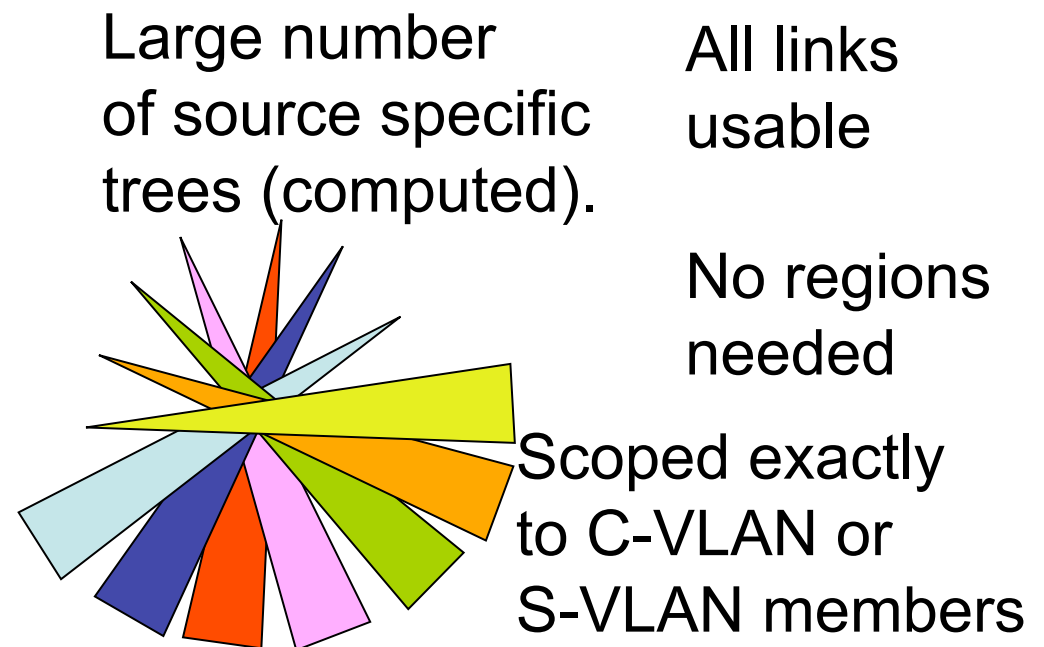
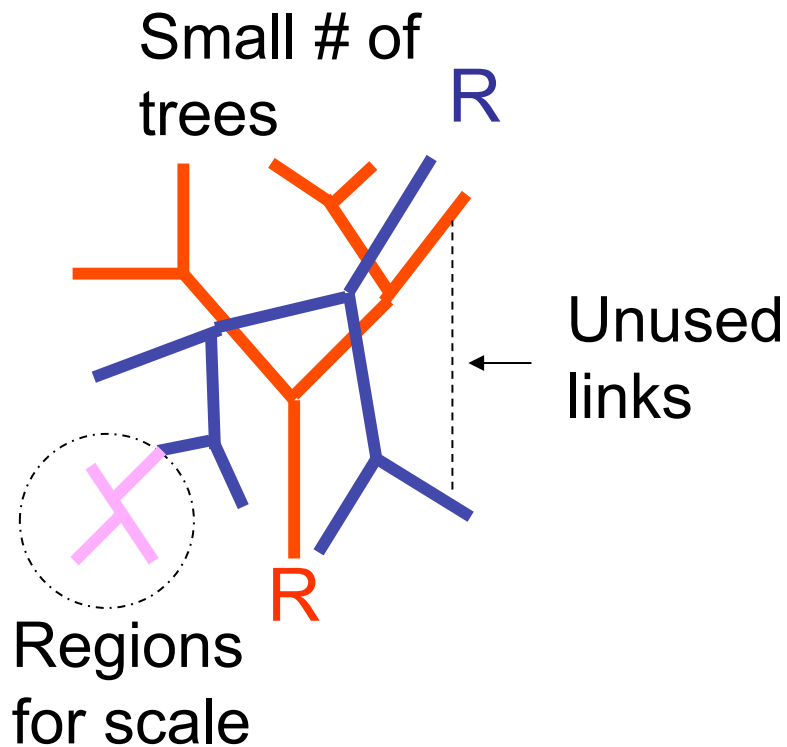


Learning restricted to edges and only where I-SID tree reaches. Mac-in-Mac encaps.

Outline

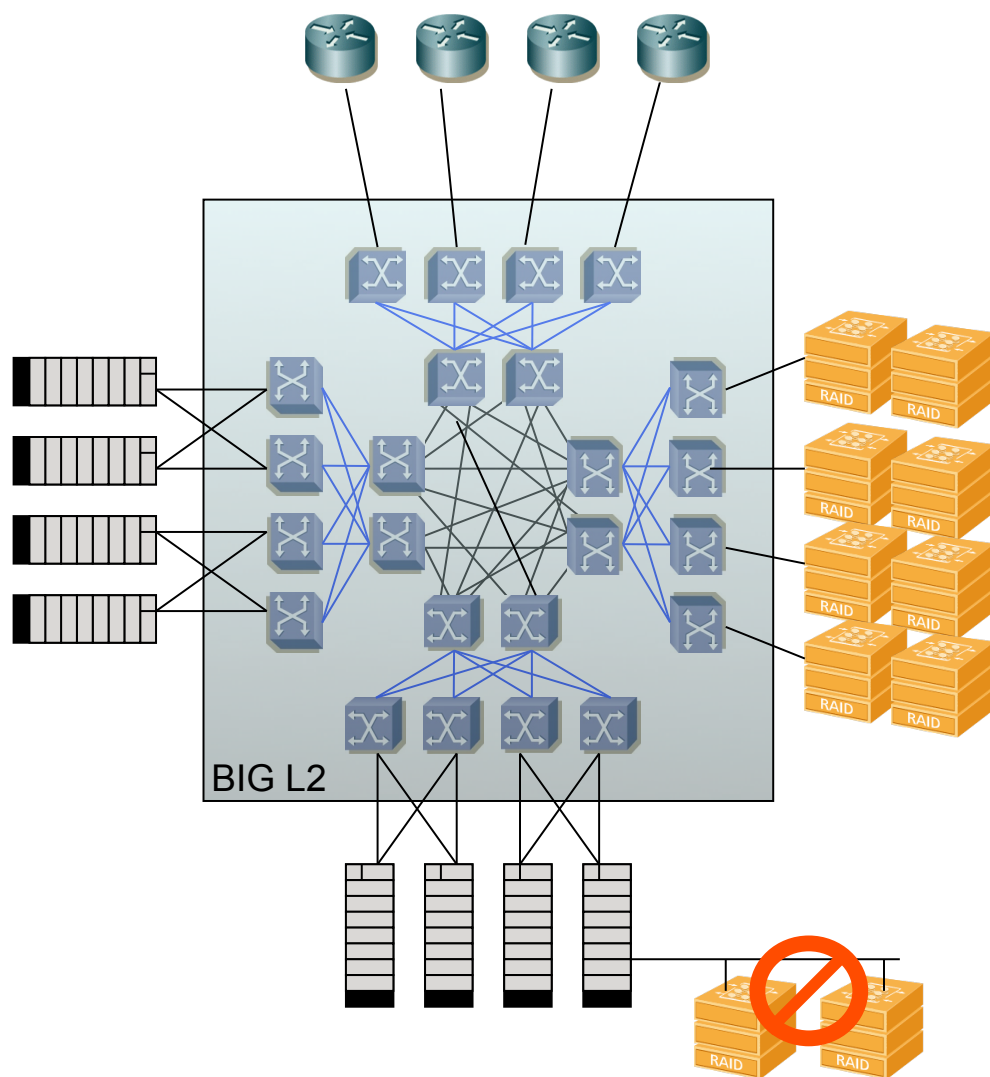
- Challenges
- What is 802.1aq/SPB
- **Applications**
- How does it work
- Short Demo (remote switches if possible)

Application (M|R)STP replacement



- Many more nodes without regions
- Low effort to get good routing
- Fast convergence – link state v.s. distance vec
- Address isolation m-in-m.

Data Center - trends

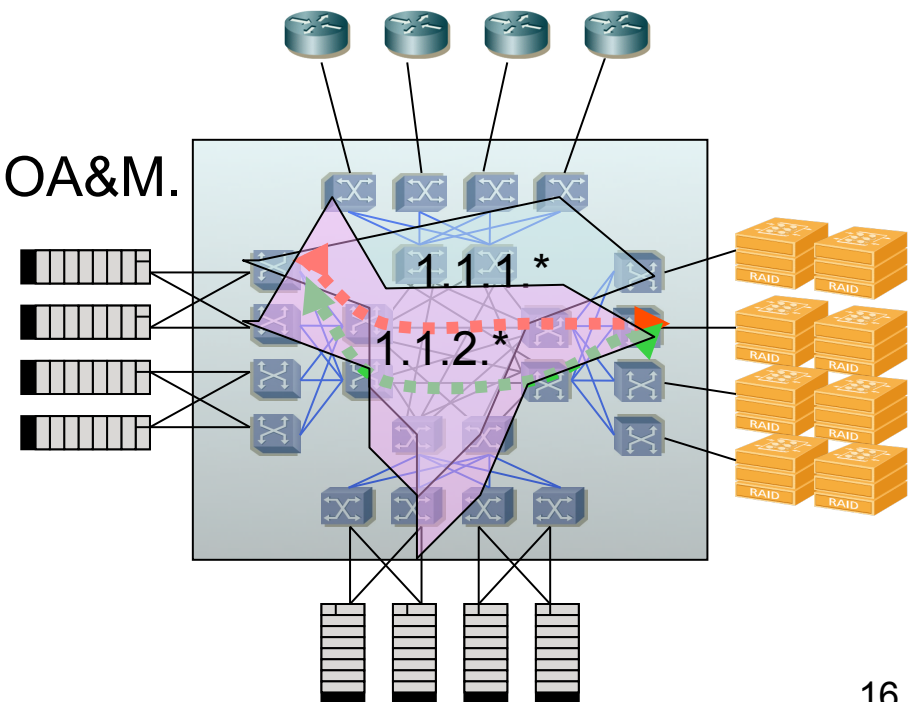


Treat DC network as one big L2 switch by combining 100's of smaller switches in 'non blocking' topology – why?

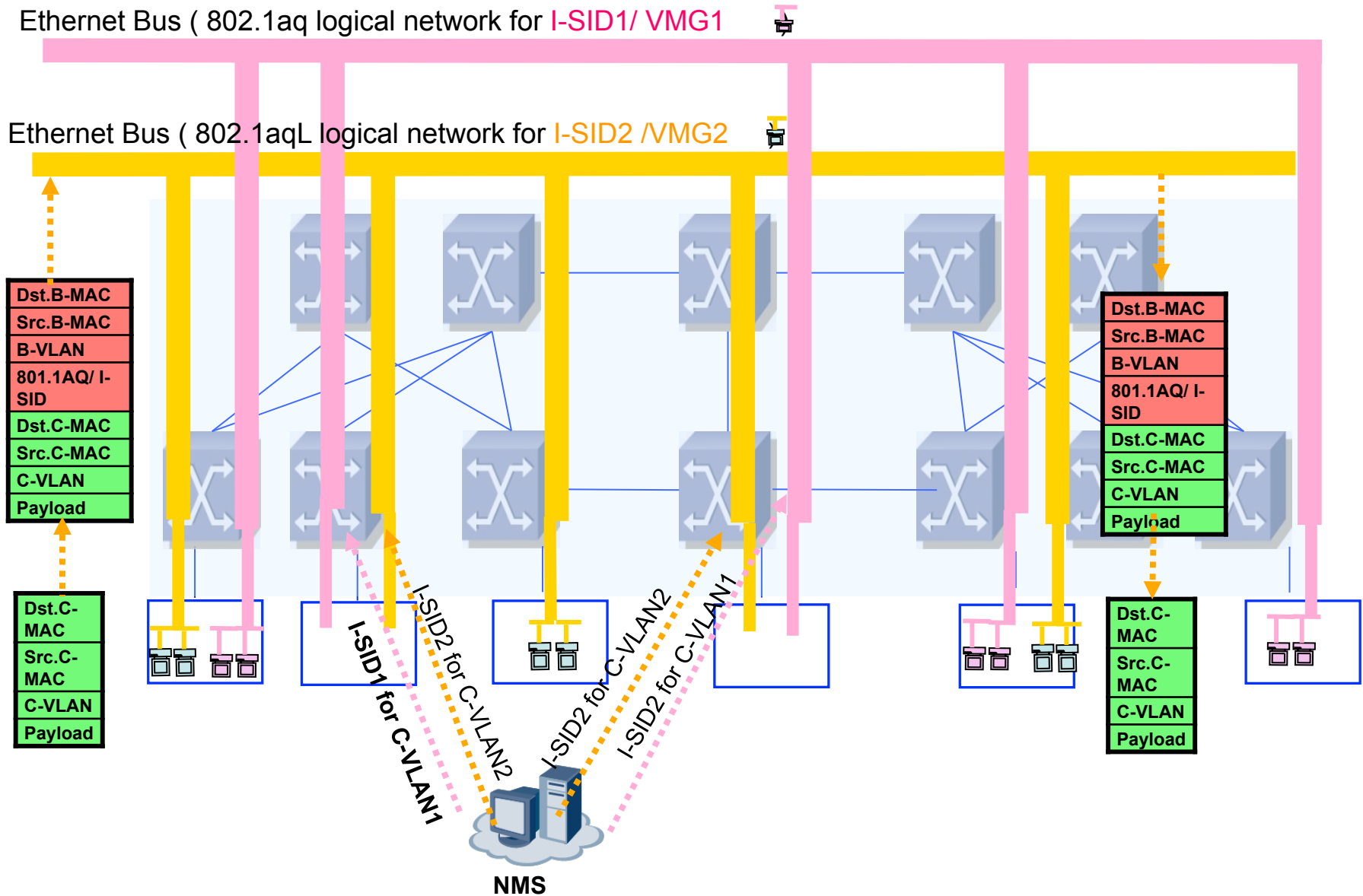
- Any server anywhere.
- Any router anywhere.
- Any appliance anywhere.
- Any VM anywhere.
 - Any IP address anywhere.
 - Any subnet anywhere.
- Any storage anywhere.
- Minimal congestion issues.
- Total flexibility for power use

Application Data Center

- Multiple shortest path routing
 - inter server traffic
- Deterministic traffic flows.
- Flexible subnet – expand/shrink anywhere.
 - Virtualization operates in subnet.
- Fully compatible with all 802.1 Data Center Bridging protocols & OA&M.
- Address isolation through m-in-m
- Fast recovery
- No loops



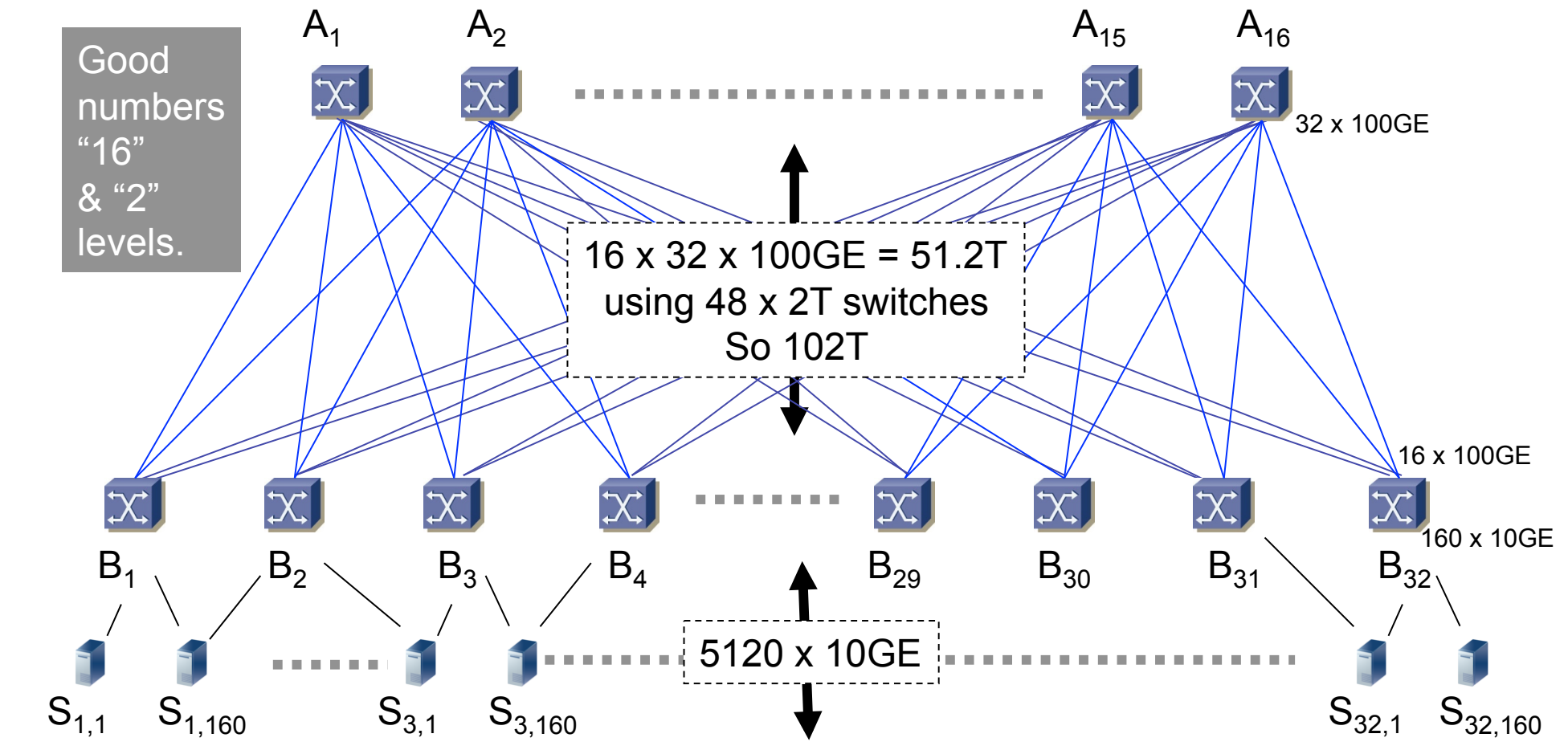
Application Data Center VM 'hot' migration (no interruption)



Application Data Center (cont'd)

- Totally compatible with Vmware server functions:
 - OA&M, motion, backup etc.
 - Apps that sit on Vmware 'just work'.
- Totally compatible with Microsoft load balancing (multicast over the L2)
- VRRP transparent (primary/stdby rtr per subnet) or proprietary variations on same protocol.
- It just makes the L2 part of the DC larger and better utilized.
- Compatible with emerging Inter DC overlay work.

Application High Performance switching cluster – assume 100GE NNI links



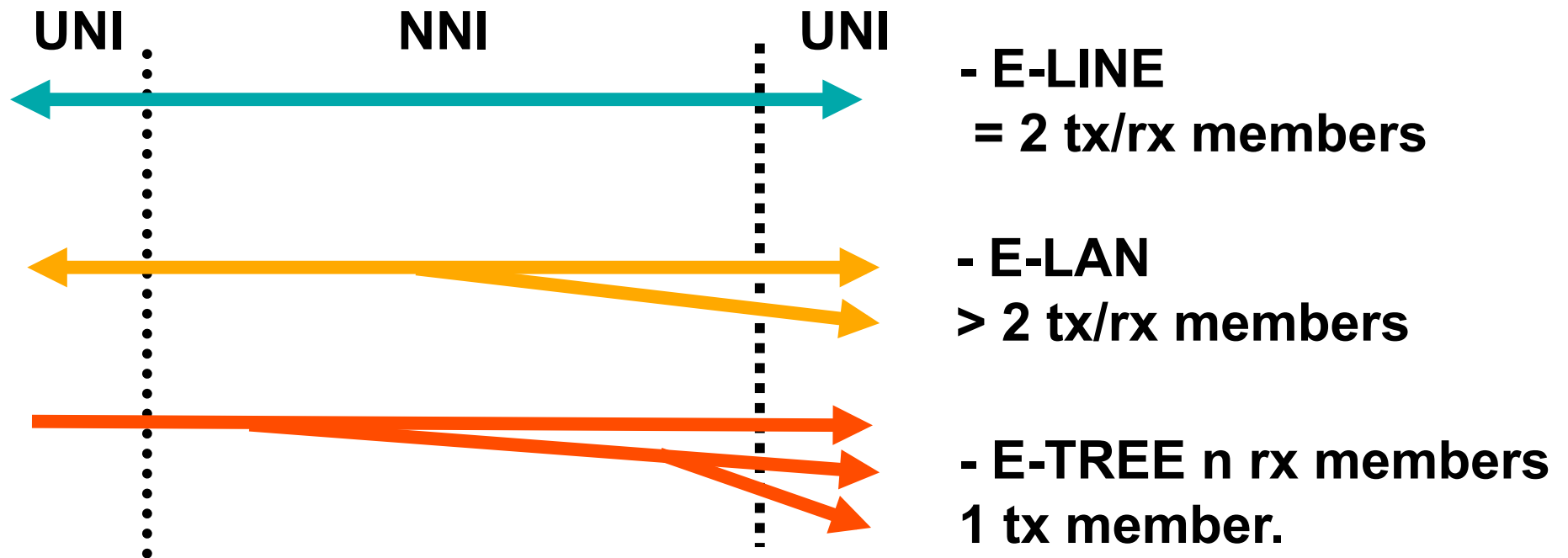
- 48 switch non blocking 2 layer L2 fabric
- 16 at “upper” layer $A_1..A_{16}$
- 32 at “lower” layer $B_1..B_{32}$
- 16 uplinks per B_n , & 160 UNI links per B_n
- 32 downlinks per A_n

- $(16 \times 100\text{GE per } B_n) \times 32 = 512 \times 100\text{GE} = 51.2\text{T}$
- 160 x 10GE server links (UNI) per B_n
- $(32 \times 160)/2 = \mathbf{2560 \text{ servers @ } 2 \times 10\text{GE}}$ per

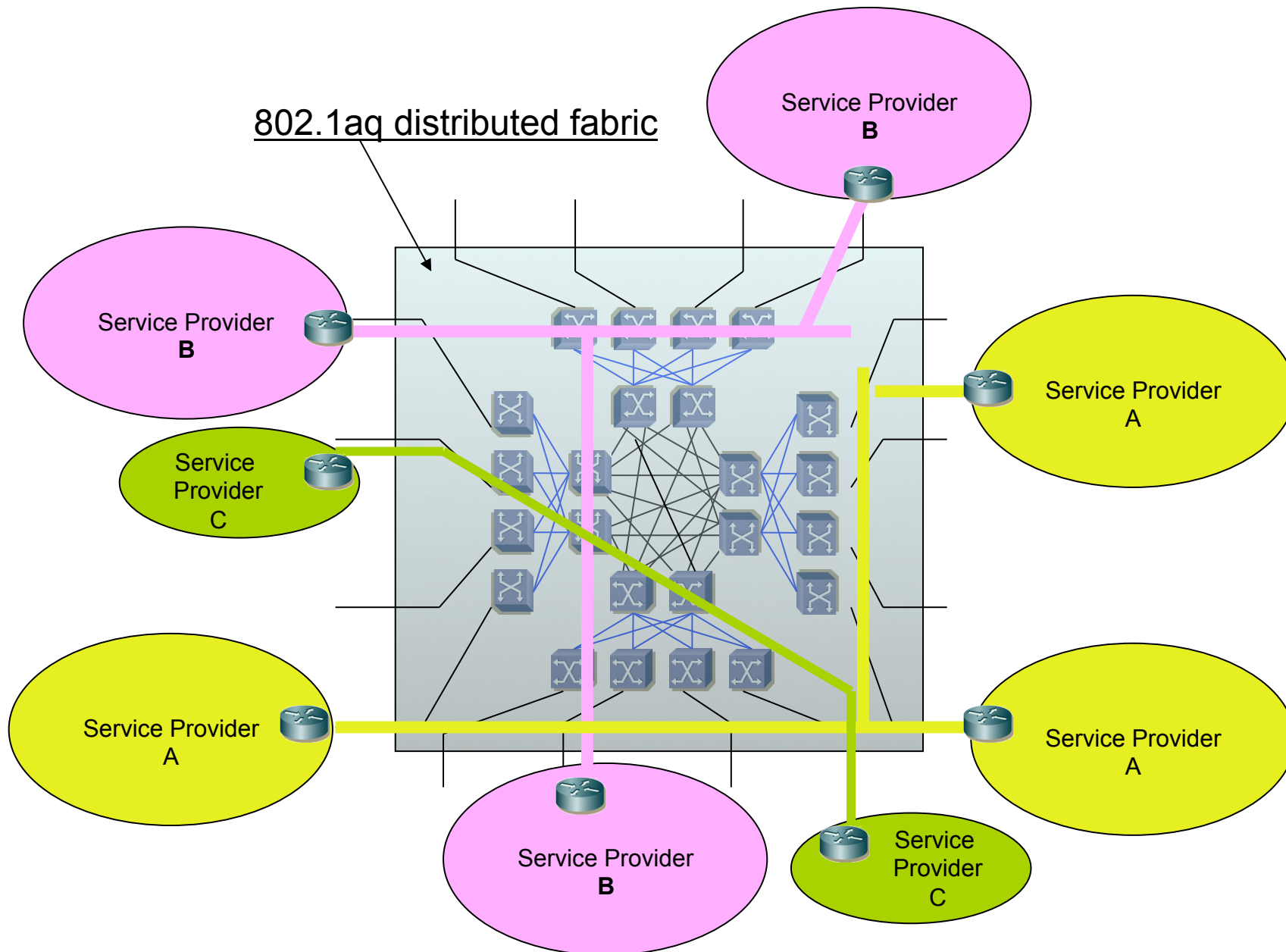
100+ Terra non blocking interconnection
fabric (if switches non blocking)

Application Metro/L2VPN

- Very light weight L2VPNs (2^{24} data path) of:
E-LAN, E-LINE, E-TREE flavors (a very cheap VPLS)
- Can do VPLS style head end replication
- Can do p2mp style transit replication (just one tx flag).
- Can support receive only membership (E-TREE)



L2 Internet Exchange Point



Outline

- Challenges
- What is 802.1aq/SPB
- Applications
- **How does it work**

How does it work?

- **From Operators Perspective**

- Plug NNI's together
- Group ports/c-vlan/s-vlan at UNIs that you want to bridge (2^{24} groups='services' m-in-m mode.)
- Assign an I-SID to each group..

- **Internally**

- IS-IS reads box MAC, forms NNI adjacencies
- IS-IS advertises box MACs (so no config).
- IS-IS reads UNI port services and advertises.
- Computations produce FIBs that bridge service members.

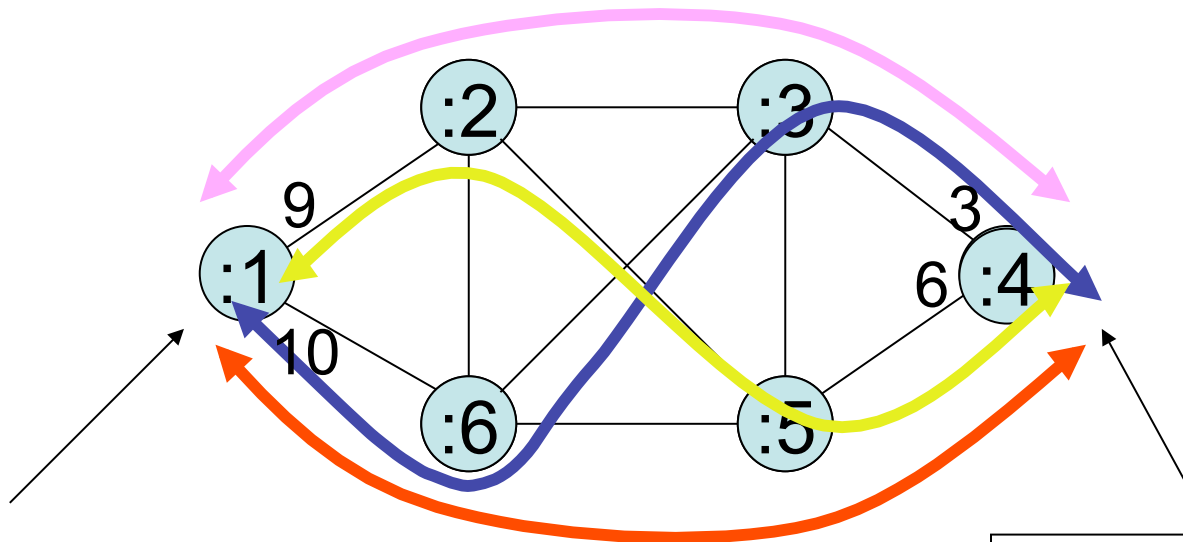
Data Path (M-in-M mode)

- C-vlan/S-vlan or untagged traffic arrives at UNI
- Its encapsulated with B-SA of bridge
- Its encapsulated with I-SID configured for group
- Its encapsulated with B-VID chosen for route
- C-DA is looked up, if found B-DA is set
- C-DA not found, B-DA is multicast that says:
 - Multicast to all other members of this I-SID group from 'me'. Or can head-end replicate over unicast.
 - C addresses to B address association learned at UNI only.

FDB (unicast M-in-M mode)

- A unique shortest path from node to all others is computed.
- BMAC of other nodes installed in FIB pointing to appropriate out interface.
- Above is repeated for 16+ shortest paths each causes a different B-VID to be used.
- Symmetry is assured through special tie-breaking logic. 16+ different tie-breaking algorithms permit 16+ different shortest paths.

FDB visually: ucast m-in-m mode



MAC	BVID	IF
:4	1	9
:4	2	9
:4	3	10
:4	4	10

MAC	BVID	IF
:1	1	3
:1	2	6
:1	3	3
:1	4	6

FDB (mcast M-in-M mode)

If no services require tandem replication
there is no tandem FDB:

Very VPLS like .. Pretty boring....head
replication over unicast paths .. Yawn..

Else (mp2mp like but without signaling)

If my node is on a unique shortest path
between node **A** , which transmits for a
group **I**, and node **B** which receives on group
I, then:

merge into the FDB an entry for traffic from
{ **A**/Group **I** } to the interface towards **B**.

How does it work – transit multicast format (n/a for head replication)

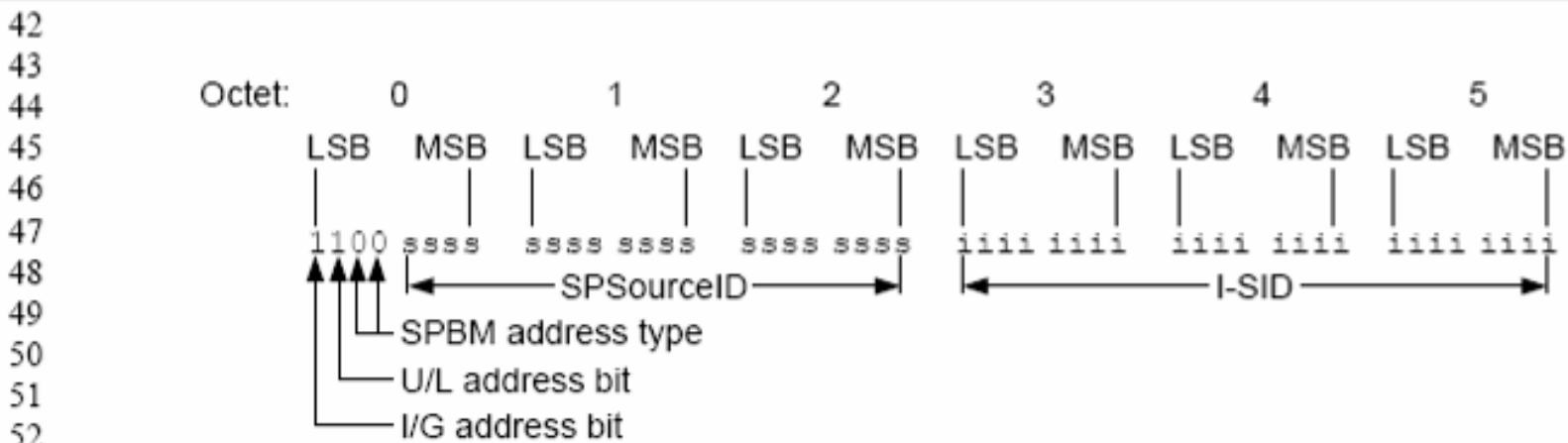


Figure 27-3—SPBM Group MAC Addresses used by this standard

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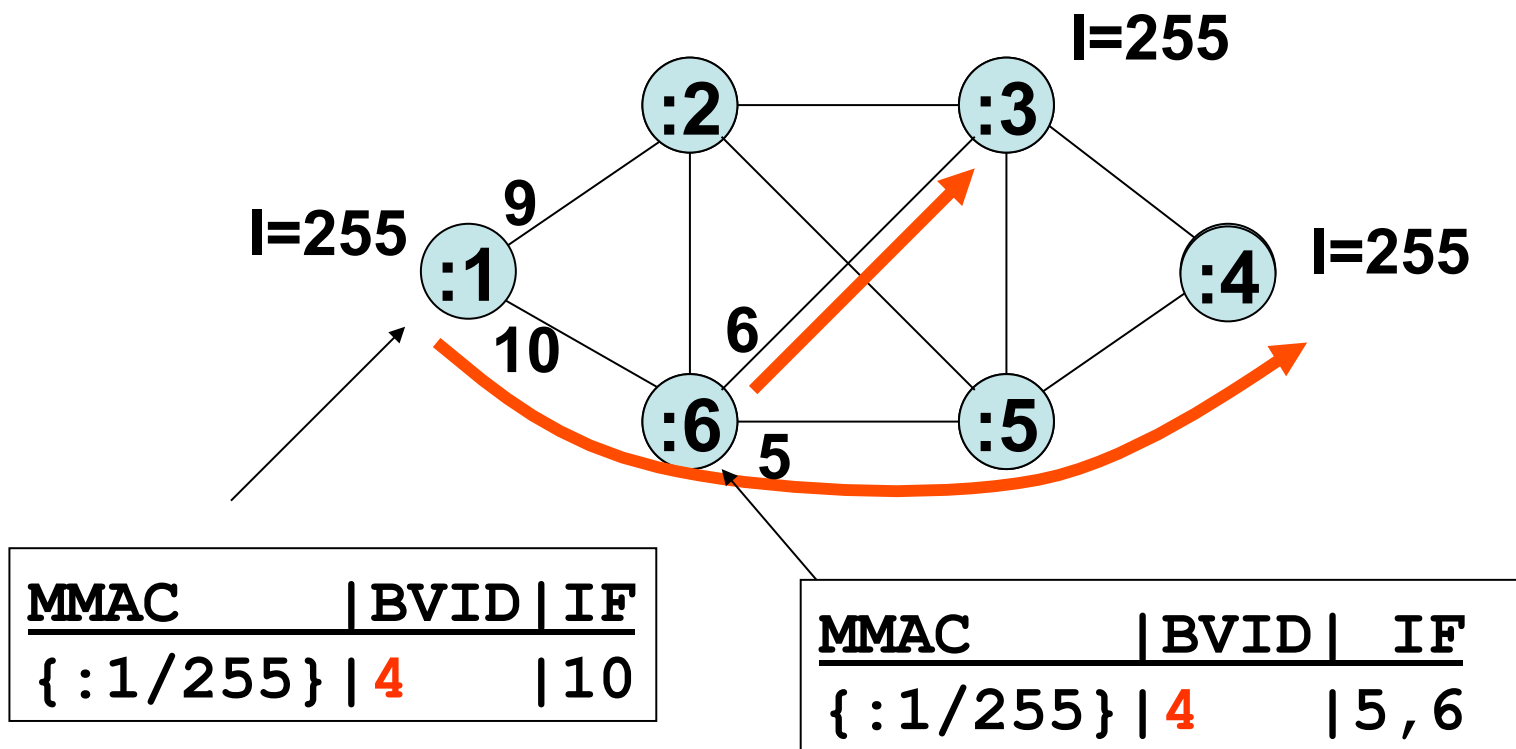
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Example: { SOURCE: 0A-BC-DE / ISID: fe-dc-ba }

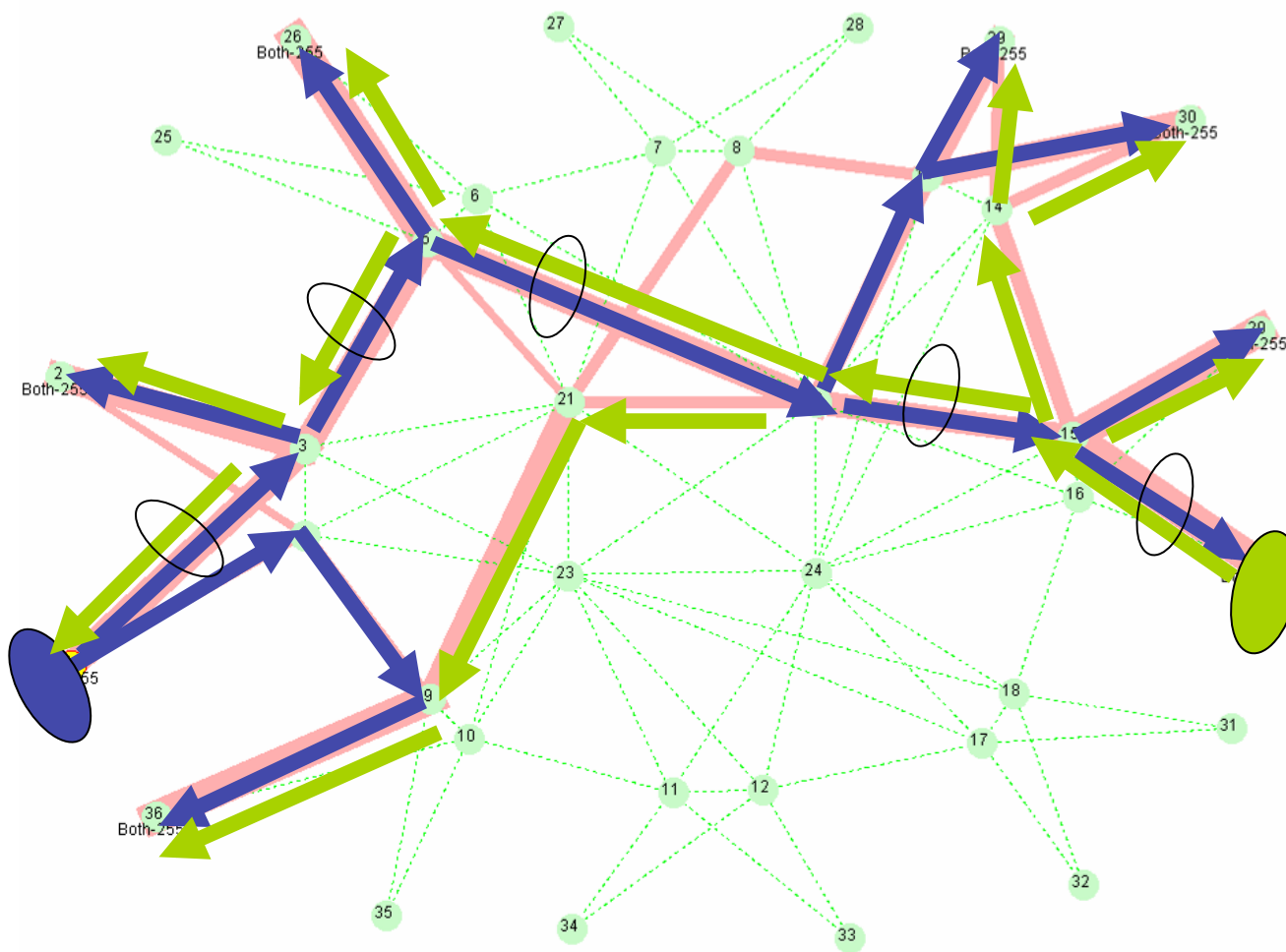
MMAC-DA: **A3-BC-DE-FE-DC-BA**

0011

FDB visually: mcast m-in-m mode



Animation for 8 member E-LAN '255'



I-SID 255 has 8 members

Shown are all routes used by this I-SID in pink.

Two trees shown blue/green.

Note symmetry of trees between source/dest

If transit multicast selected fork points in trees are replication points.

The Control Plane (m-in-m mode)

- Industry standard IS-IS Link State Protocol is basis for 802.1aq.
- Does not require any IP to operate.
- Does not preclude IPV4 or IPV6 being present in same IS-IS instance.
- SYSID carries B-MAC address
- Introduces no new PDU's to IS-IS.
- Hello TLVs augmented to pass Equal Cost Algorithm / Vid information and new NLPID.
- Update TLV's augmented to advertise SPB specific link costs.
- Update TLVs augmented to advertise ISID information.
- Update TLVs augmented to advertise nodal 'short form' name SPSOURCEID (transit mcast only).

802.1aq ISIS LSP extensions at a glance

LSPID	Seq Num	Checksum ...
-------	---------	--------------

4455.6677.0001.00-00 0x0000007a 0xc01f ...

SOURCE 4455.6677.0001.00 (1)

HOST NAME	Instance 1
-----------	------------

NLPID SPB (0xC1) (2)

AREA ADDR 22.3344

NBR	ID	4455.6677.0004.00	COST: 10
-----	----	-------------------	----------

NBR ID 4455.6677.0003.00 COST: 10

SPSOURCEID 07-00-01

SPB ECT-ALGORITHM 1 ECT-VID 101 (4)

SPB ECT-ALGORITHM 0 ECT-VID 100

SPB BMAC 44-55-66-77-00-01

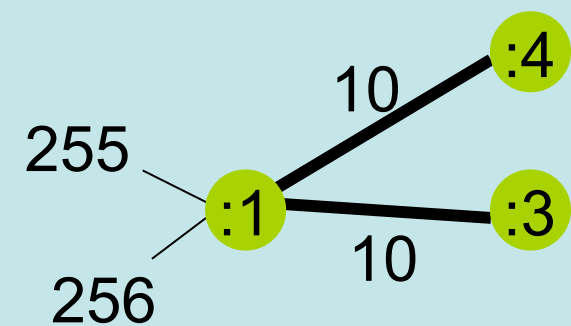
ECT-VID 100

SPB ISID 255T&R (5)

SPB BMAC 44-55-66-77-00-01

ECT-VID 101

SPB ISID 256T&R



LSP fragment for
node :1 with 2 peers
:4 and :3 and two
services 255, 256

TIE BREAKING/ECT

Symmetric Tie Breaking Requires that:

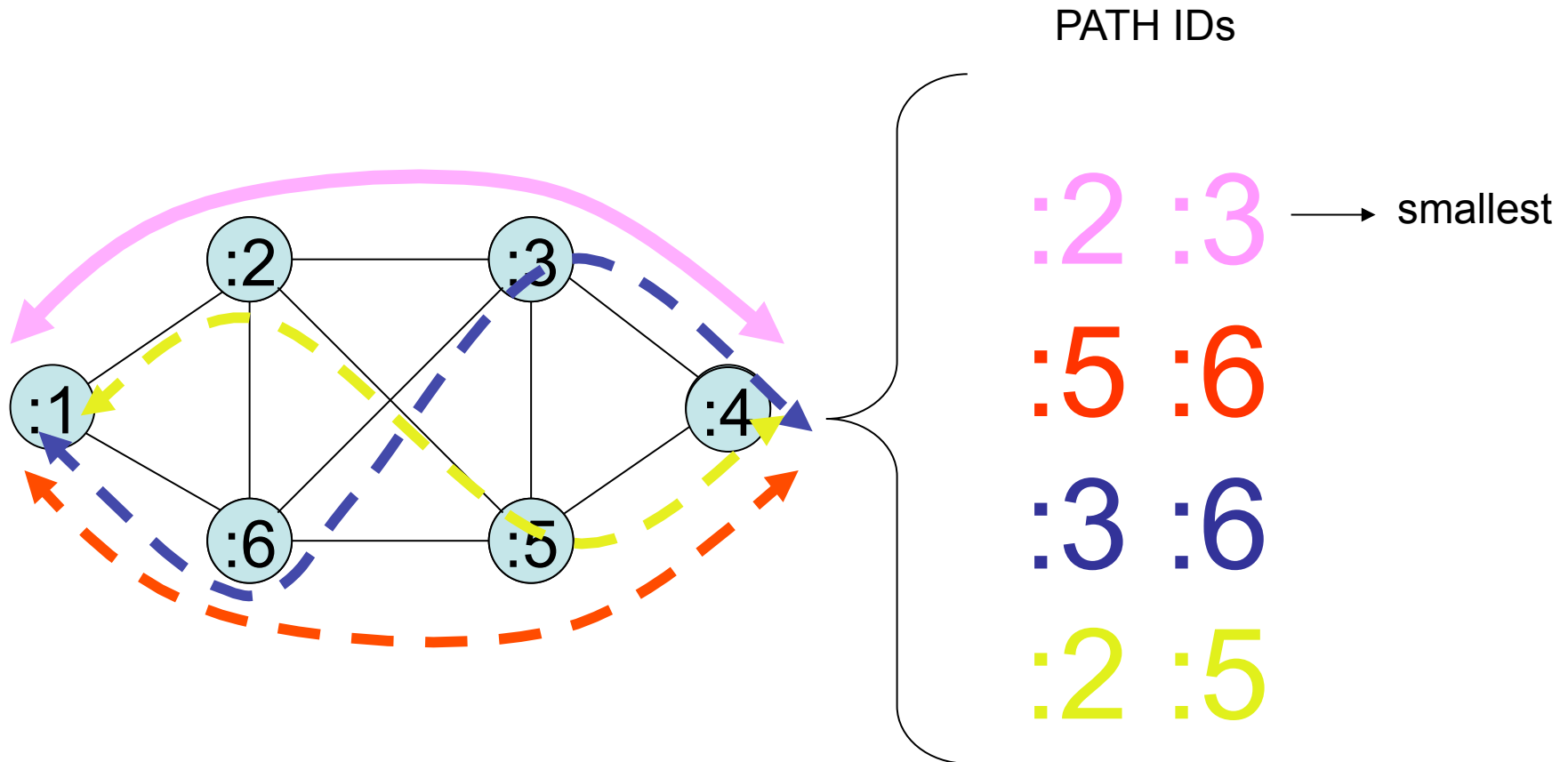
- Forward and Reverse path are the same
- Every subset of the end to end shortest path must also be a shortest path under the same tie breaking rules (transitivity): I.e.

$\text{SPF}(A,C) = \text{SPF}(A,B) \parallel \text{SPF}(B,C)$ <- required for hop by hop fwding

Solution – assign to a path a ‘Path Identifier’

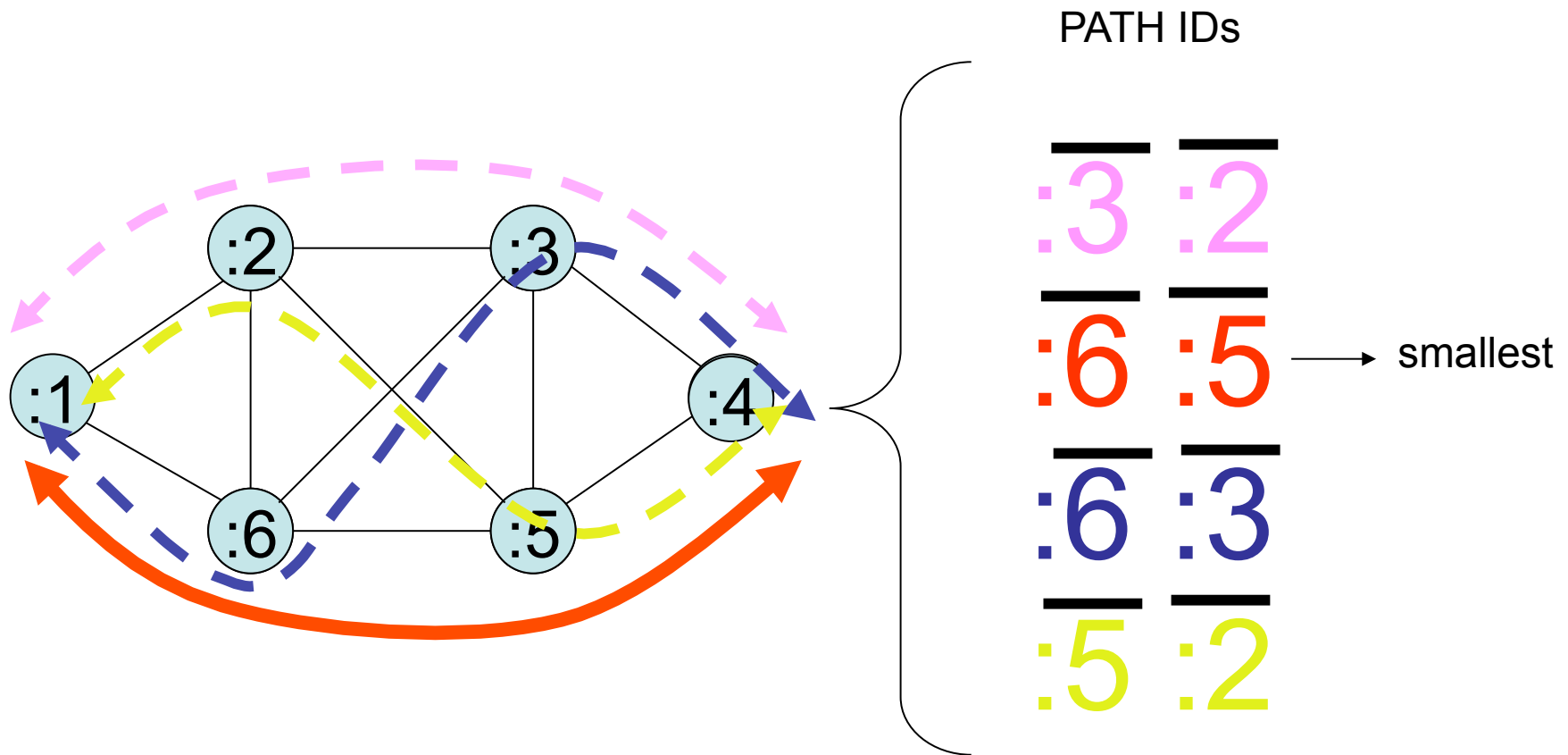
- ***Path Identifier*** := Sorted List of Nodes between SRC/DST
- Given a choice between two paths. Pick path with smallest ***Path Identifier***.

TIE BREAKING ECT – CONT'D



A path is symmetrically/deterministically chosen by picking the path with the Lowest path identifier (i.e. path with lowest Bridge Identifier on it). Transitivity proof left as exercise for the interested reader.

TIE BREAKING ECT – CONT'D



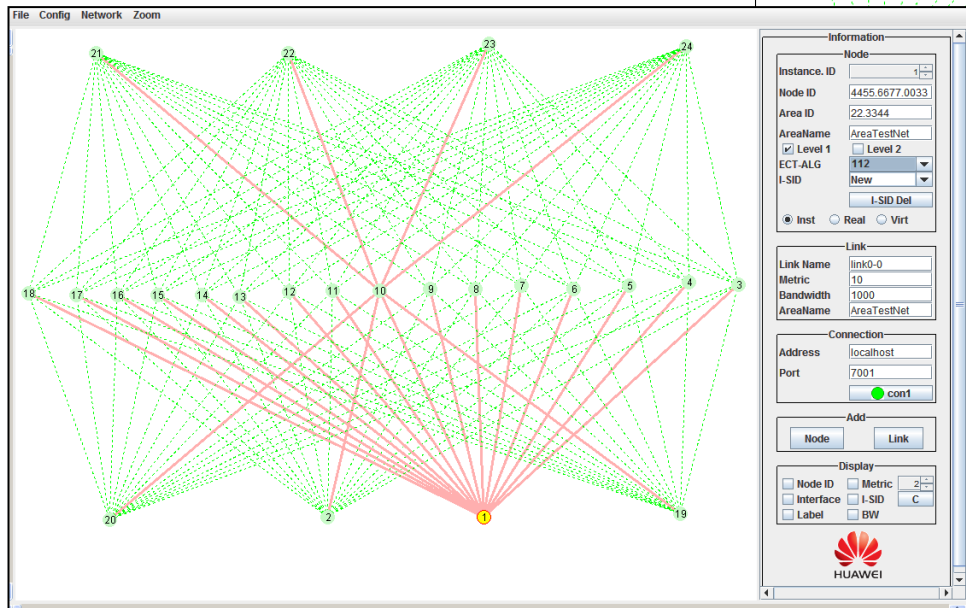
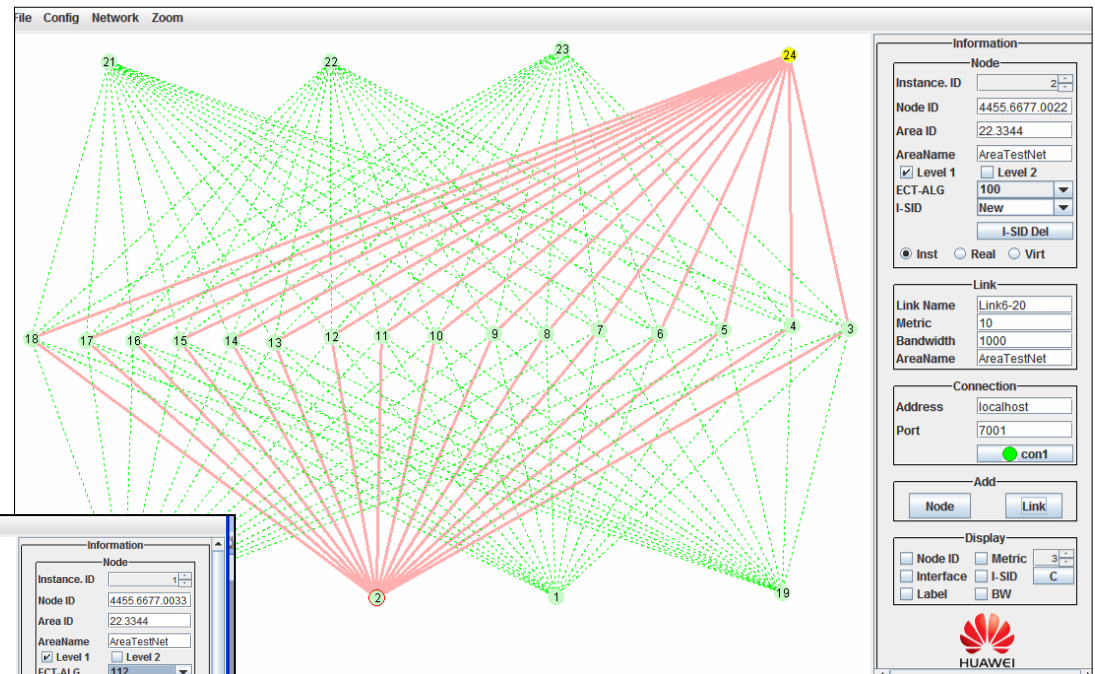
Applying a known bit mask to Bridge ID's and running the same algorithm will cause different selection. Inversion (shown) is equivalent to picking largest Path ID.

TIE BREAKING ECT – CONT'D

- In total 16 bit masks. 0x00..0x11,0x22 ... 0xee,0xff are provided.
- This yields 16 possible shortest paths.
- Paths can be tuned since Bridge Priority is included in Bridge Identifier as high 2 bytes. *Bridge Identifier = Bridge Priority:16 || SYSID*
- Adjusting Bridge Priority up or down allows tuning of ECT.
- Picking 16 priorities and assigning down network 'cut' yields perfect spread through cut. Eg: spine of a DC fabric.
- The 802.1aq TLV's support Opaque ECT behavior to allow further work on ECT.
- The Opaque ECT behaviors are designed to allow ECMP (hash based) as next step.

TIE BREAKING ECT – CONT'D

Can get perfect balance
down spine of a two layer
16 ECT L2 Fabric. Shown
Are all 16 SPF's from 2->24



16 different SPF trees
Each use different spine
as replication point.
Shown is one of the 16
SPF's from/to node 1.

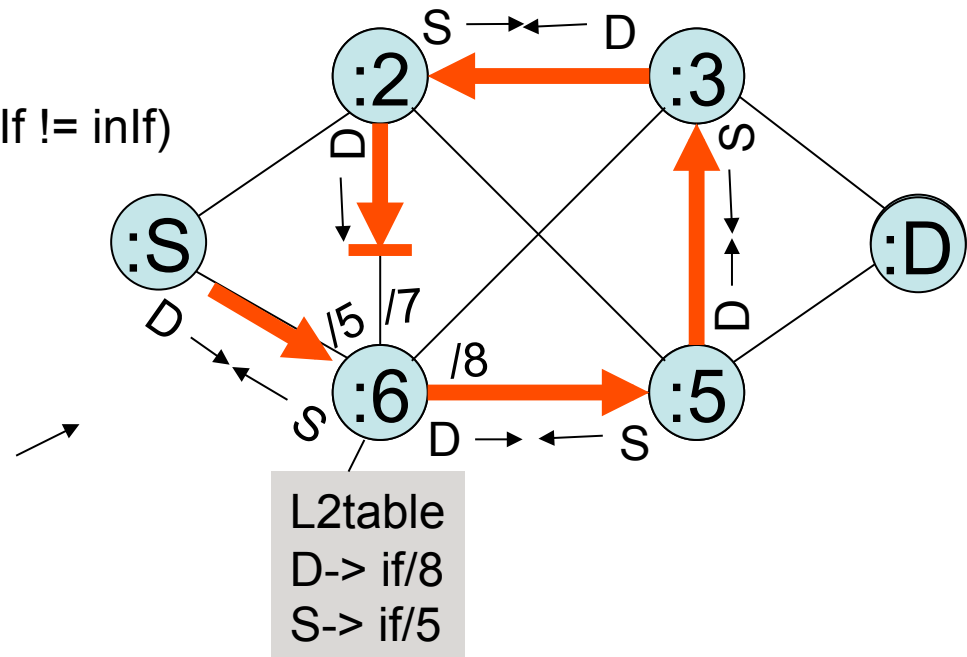
Loop Suppression & Avoidance

Suppression

- done on the data path using an SA check.
- prevents 99.99% loops if FDB's create one.
- no impact on convergence rates.
- exploits symmetric/congruence properties of routing.
- uses reverse learning options of most h/w to discard.

```
Forward(pkt, inIf) {  
  if (l2table[pkt.sa,pkt.vid].outIf != inIf)  
    discard (pkt);  
  else ...  
}
```

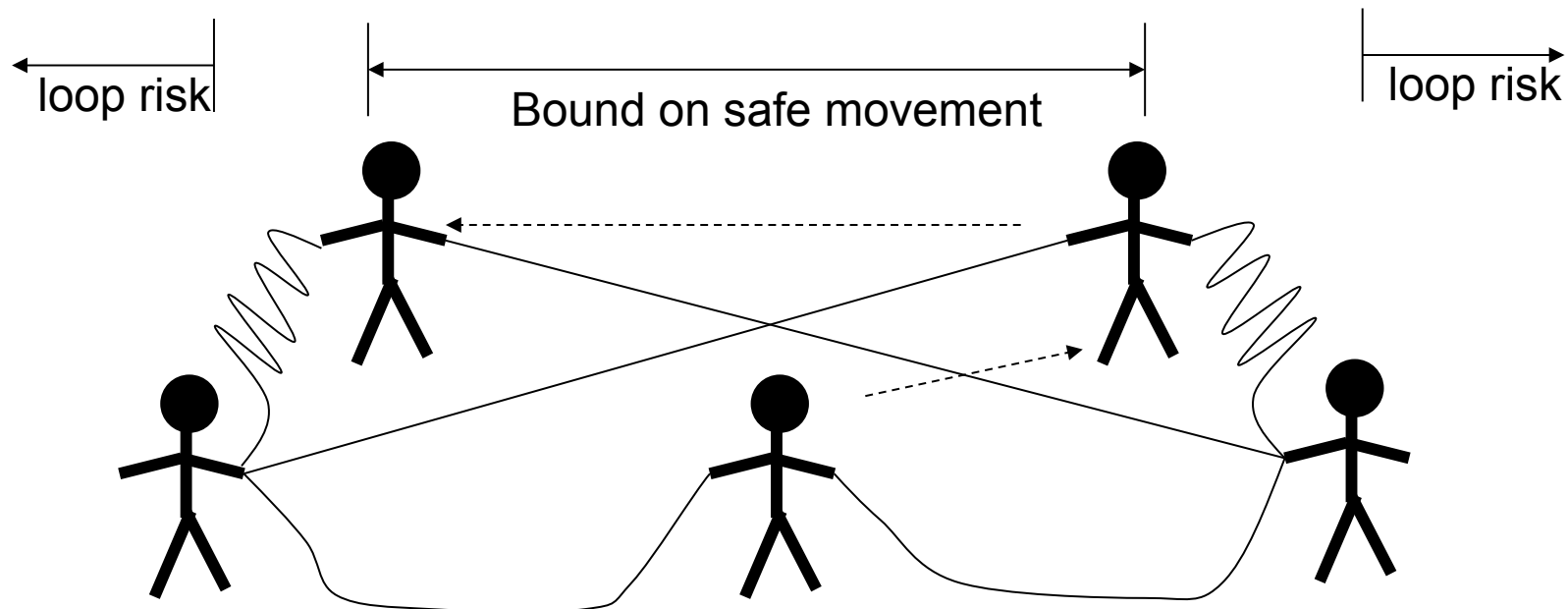
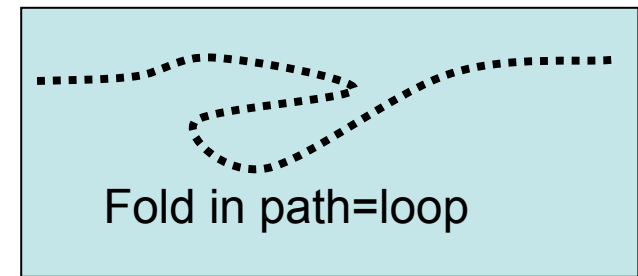
On :6
/7 != /5 so discard



Loop Suppression & Avoidance

Avoidance

- hellos augmented with topology 'digests'
- mismatched topology => some forwarding entries unsafe.
- blocks 'unsafe' entries.
- Unsafe => distance changed outside "safe zone"



802.1aq OA&M (inherited *by design*)

Service/Network Layer – 802.1ag Connectivity Fault mgmt

- Hierarchy (honors maintenance levels/abstraction)
 - Continuity Check
 - L2 traceroute
 - L2 ping

Link Layer – 802.3ah

- Link Monitoring (logical/physical)
- Remote Failure Indication
- Remote Loopback

Service Layer - Y.1731

- Multicast Loopback – depends on congruency/symmetry
- Performance Measurements (Loss/Delay etc.)
- One way/two way delays – symmetry important

802.1aq OAM capabilities

1. Continuity Check (CC)

- a) Multicast/unidirectional heartbeat
- b) Usage: Fault detection

2. Loopback – Connectivity Check

- a) Unicast bi-directional request/response
- b) Usage: Fault verification

3. Traceroute (i.e., Link trace)

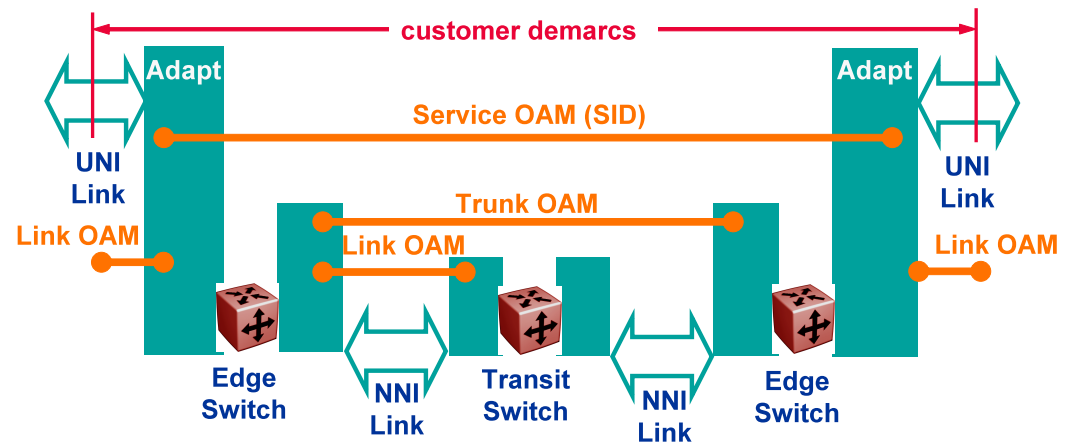
- a) Trace nodes in path to a specified target node
- b) Usage: Fault Isolation

4. Discovery (not specifically supported by .1ag however Y.1731 and 802.1ab support it)

- a) Service (e.g. discover all nodes supporting common service instance)
- b) Network (e.g. discover all devices common to a domain)

5. Performance Monitoring (MEF10 and 12 - Y.1731 for pt-pt now extending to pt-mpt and mpt-mpt)

- a) Frame Delay, Frame Loss, Frame Delay Variation (derived)
- b) Usage: Capacity planning, SLA reporting



MIBs (also all OA&M etc. inherited)

SpbSys	Basic config data
SpbSysDynamic	Basic run time data
SpbEctStatic	Desired Equal cost behaviors
SpbEctDynamic	Actual Equal cost behaviors
SpbAdjStatic	Desired Adjacency data (metric etc.)
SpbAdjDynamic	Actual Adjacencies
SpbTopNode	All other nodes in the SPB region and basic data
SpbTopEct	All other ECT behavior in the SPB region
SpbTopEdge	All other Edge data in the SPB region
SpbTopSrv	All services in the SPB region, who hosts etc.

Rapid Recovery

- ISIS augmented with multicast LSP flood to all 802.1aq nodes.
- Every 802.1aq node joins default service ISID 0xfffff.
- This E-LAN is just for control plane.
- LSPs can be advertised over this E-LAN
- Very fast distribution protocol (h/w multicast).
- On failure each end of link advertises over this 'default' E-LAN (in addition to normal updates).
- Reaches all 802.1aq participants at h/w multicast speed with no CPU involvement transit.
- Conceptually like having a shared LAN joining all nodes with a physical port but no DR election etc. is done, only used as unreliable very fast distribution mechanism backed up by normal IS-IS hop by hop LSPs.

TLV Summary

PDU	TLV	SUB-TLV	TYPE	TYPE	#OCCURRENCE
IIH					
		MT-Port-Capability	143		
		SPB-B-VID		5	1
		SPB-Digest		6	1
LSP					
		MT-Capability	144		
		SPB-Inst		1	1
		SPB-I-OALG		2	>=0
		SPBM-SI		3	>=0
		SPBV-ADDR		4	>=0
		MT-Intermediate-System	222		
		SPB-Metric		12	1
		SPB-A-OALG		13	>=0

No new PDU, no changes to IS-IS state machine logic

TLV – SPB-Digest (in Hello)

All TLV's in context of an MTID TLV

[illegible]

ALSO in Hello the **NLPID value 0xC1** is used and advertised.

```
HELLO padding unchanged ..
```

TLV – SPB-B-VID (in Hello)

[illegible]

TLV – SPB-Instance

```

+---+---+---+---+---+---+
|Type = SPB-Inst| = 1
+---+---+---+---+---+---+
|   Length       |   (1 byte)
+---+---+---+---+---+---+
|               CIST Root Identifier   (4 bytes)               |
+---+---+---+---+---+---+
|               CIST Root Identifier (cont)   (4 bytes)         |
+---+---+---+---+---+---+
|               CIST External ROOT Path Cost   (4 bytes)       |
+---+---+---+---+---+---+
|      Bridge Priority                |   (2 bytes)
+---+---+---+---+---+---+
|R R R R R R R R R R R R|V|               SPSOURCEID       |
+---+---+---+---+---+---+
| Num of Trees   |   (1 bytes)
+---+---+---+---+---+---+
|               VLAN-ID (1) Tuples       (8 bytes)           |
+---+---+---+---+---+---+
|               VLAN-ID (N) Tuples       (8 bytes)           |
+---+---+---+---+---+---+

```

TLV – SPB-Metric

```
+---+---+---+---+---+
|Type=SPB-Metric| = 12
+---+---+---+---+---+
|   Length       |      (1 byte)
+---+---+---+---+---+---+---+---+---+---+---+---+---+---+---+
|      SPB-LINK-METRIC                               | (3 bytes)
+---+---+---+---+---+---+---+---+---+---+---+---+---+---+---+
| Num of ports   |      (1 byte)
+---+---+---+---+---+---+---+---+---+---+---+---+---+---+---+
|   Port Identifier      |      ( 2 bytes)
+---+---+---+---+---+---+---+---+---+---+---+---+---+---+---+
```

Future Expanded ECMP behaviors ...

```
+---+---+---+---+---+
|Type=SPB-A-OALG| = 13
+---+---+---+---+---+
|   Length       |      (1 byte)
+---+---+---+---+---+---+---+---+---+---+---+---+---+---+---+
|                                     Opaque ECT Algorithm      (4 bytes) |
+---+---+---+---+---+---+---+---+---+---+---+---+---+---+---+
|                                     Opaque ECT Information (variable) |
+---+---+---+---+---+---+---+---+---+---+---+---+---+---+---+
```

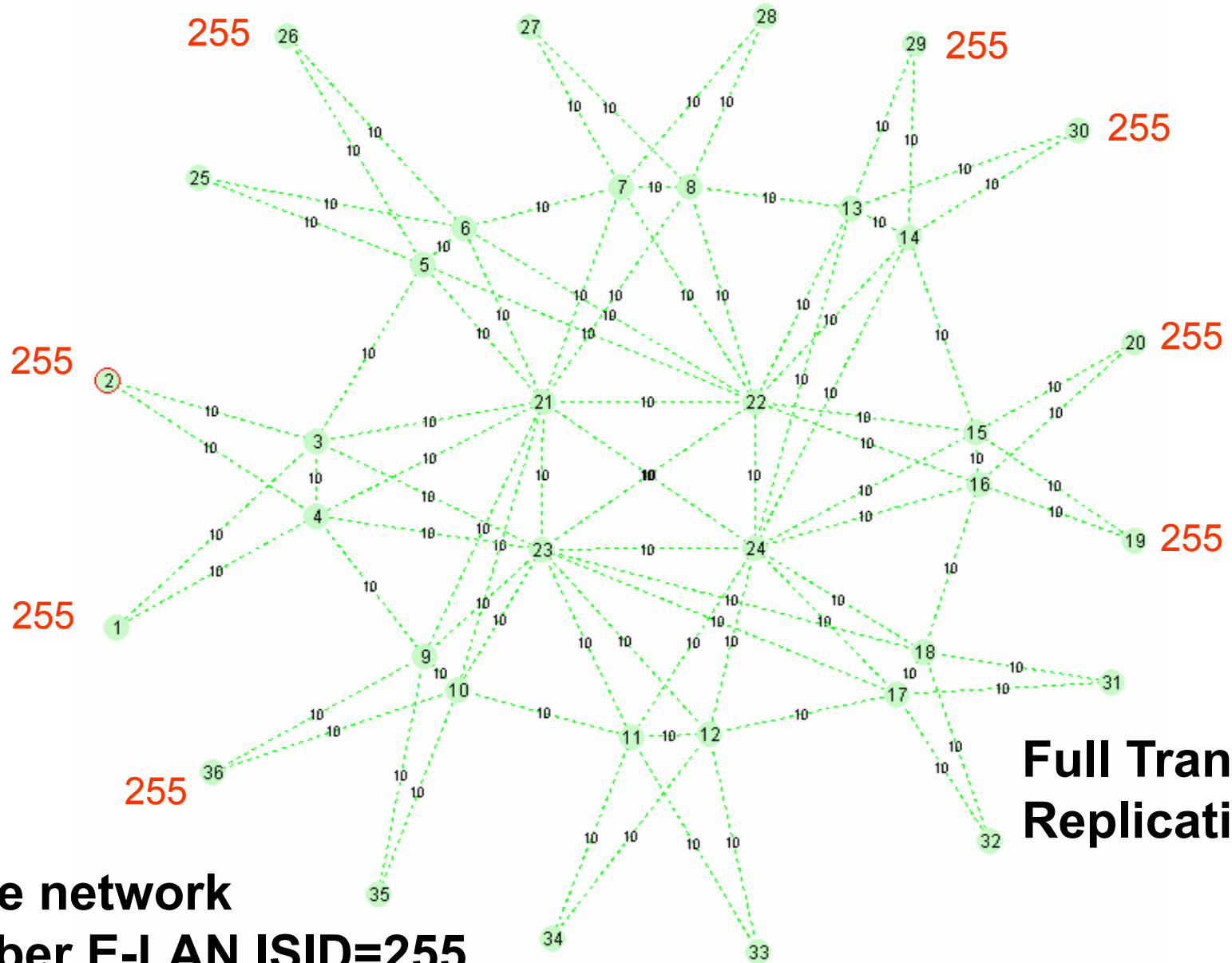
TLV – SPB-Service Instance

```
+--+--+--+--+--+--+--+--+--+--+
|Type = SPBM-SI | = 3
+--+--+--+--+--+--+--+--+--+--+
|   Length       |   (1 byte)
+--+--+--+--+--+--+--+--+--+--+
|                                     B-MAC ADDRESS                                     |
+--+--+--+--+--+--+--+--+--+--+
|   B-MAC ADDRESS (6 bytes)   | Res. |   Base-VID (12 bits)   |
+--+--+--+--+--+--+--+--+--+--+
|T|R| Reserved |                                     ISID  #1                                     |
+--+--+--+--+--+--+--+--+--+--+
|T|R| Reserved |                                     ISID  #2                                     |
+--+--+--+--+--+--+--+--+--+--+
|T|R| Reserved |                                     ISID  #n                                     |
+--+--+--+--+--+--+--+--+--+--+
```

Outline

- Challenges
- What is 802.1aq/SPB
- Applications
- How does it work
- **Short Demo (remote switch if possible)**

EXAMPLE NETWORK :



36 node network

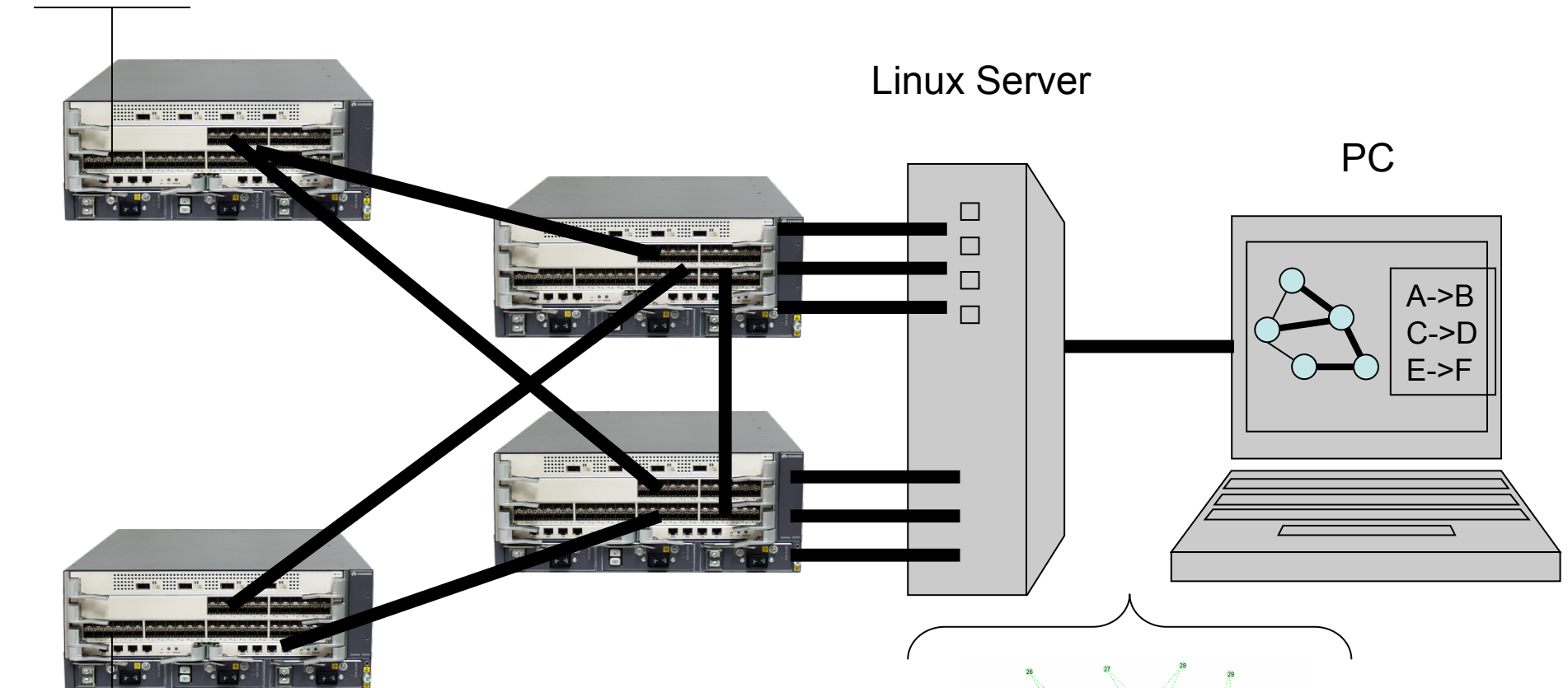
8 member E-LAN ISID=255

DEMO BASED ON OUR S9300 SERIES TERABIT SWITCHES S9303/1T, S9306/2T, S9312/4T

I-SID/255

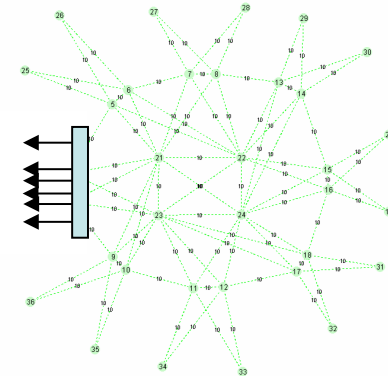
Linux Server

PC



I-SID/255

- 36 nodes
- 172 links
- 2 E-LAN:
16 members



EXAMPLE – ISIS PEERS AT NODE :3

```
<ottawa-9300-3>d spb
The current global spb information is :
Device HMAC is 44-55-66-77-00-03
Spsid is 07-00-03
Ect vlan amount is 2
Ect vlan sequence number [1] is: vlan 100 !
Ect vlan sequence number [2] is: vlan 101 !
<ottawa-9300-3>
```

```
<ottawa-9300-3>d isis peer
```

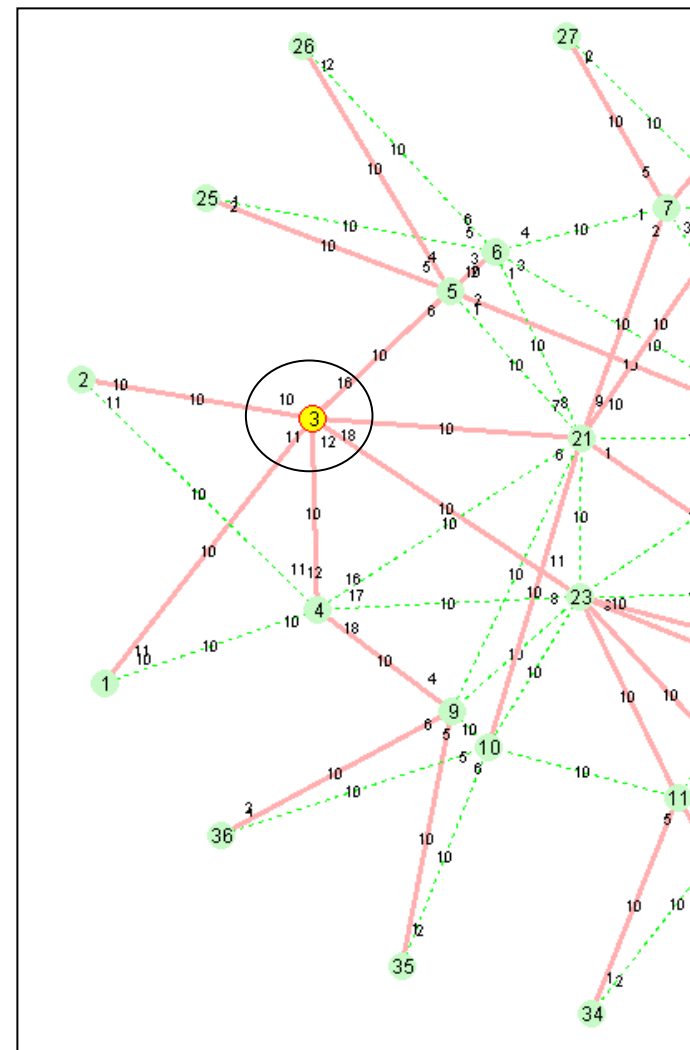
Peer information for ISIS(1)

System Id	Interface	Circuit Id	State	HoldTime	Type
4455.6677.0001	Vlanif211	00000000002	Up	26s	L1
4455.6677.0004	Vlanif212	00000000003	Up	23s	L1
4455.6677.0005	Vlanif216	004	Up	27s	L1
4455.6677.0015	Vlanif217	005	Up	27s	L1
4455.6677.0017	Vlanif218	006	Up	25s	L1

Total Peer(s) : 5

```
<ottawa-9300-3>
```

Logging on to node :3
We can see the basic SPB info
and the ISIS peers....

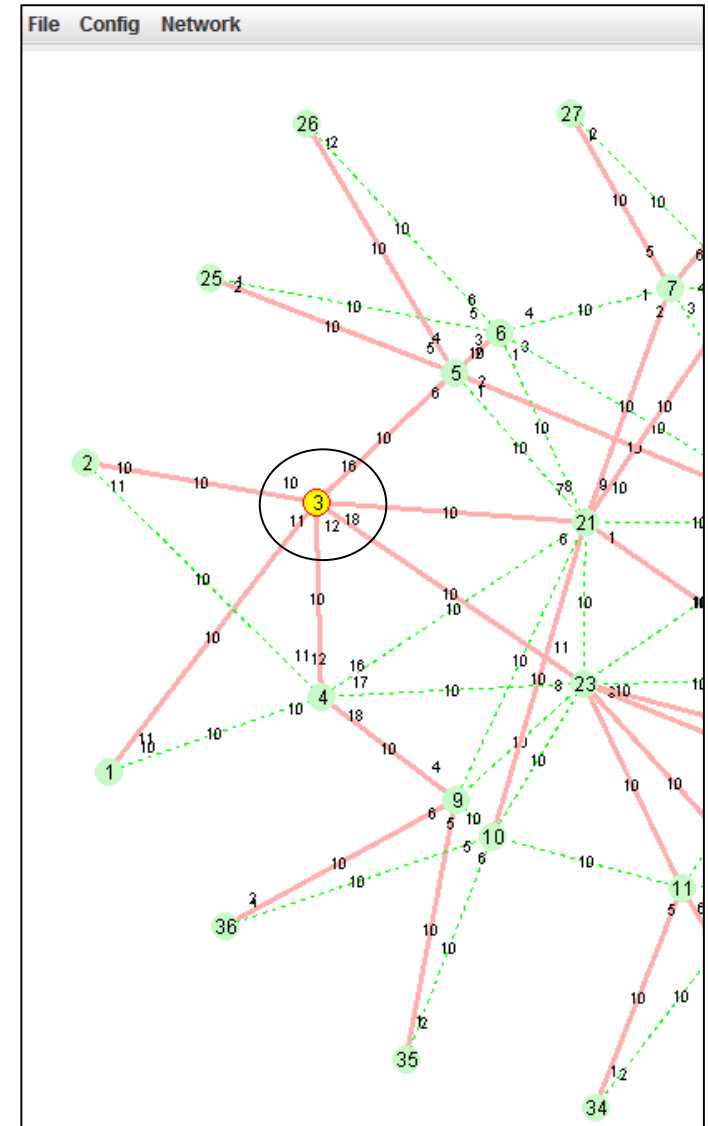


EXAMPLE – LSDB at node :3

Database information for ISIS(1)

LSPID	Seq Num	Checksum	Holdtime	Length	ATT/P/OL
4455.6677.0001.00-00	0x00000fd2	0x1cea	1044	236	0/0/0
4455.6677.0003.00-00*	0x00001448	0x3d27	683	323	0/0/0
4455.6677.0004.00-00	0x00000ff8	0xd1d9	1090	323	0/0/0
4455.6677.0005.00-00	0x00000b3b	0x9ba7	586	317	0/0/0
4455.6677.0006.00-00	0x00000b3b	0xbc31	819	293	0/0/0
4455.6677.0007.00-00	0x00000b3b	0xce10	1075	293	0/0/0
4455.6677.0008.00-00	0x00000b3e	0xebe0	288	293	0/0/0
4455.6677.0009.00-00	0x00000b3b	0x8b77	355	317	0/0/0
4455.6677.000a.00-00	0x00000b3b	0x57a	840	294	0/0/0
4455.6677.000b.00-00	0x00000b3a	0x1665	608	294	0/0/0
4455.6677.000c.00-00	0x00000b3a	0x6903	764	294	0/0/0
4455.6677.000d.00-00	0x00000b3a	0x89fd	431	294	0/0/0
4455.6677.000e.00-00	0x00000b39	0x3445	611	294	0/0/0
4455.6677.000f.00-00	0x00000b3a	0xdabe	616	294	0/0/0
4455.6677.0010.00-00	0x00000b3b	0x810e	452	294	0/0/0
4455.6677.0011.00-00	0x00000b3a	0x1b46	645	294	0/0/0
4455.6677.0012.00-00	0x00000b39	0x1b3e	447	294	0/0/0
4455.6677.0013.00-00	0x00000b3a	0x943a	419	176	0/0/0
4455.6677.0014.00-00	0x00000b3b	0xdff2	693	176	0/0/0
4455.6677.0015.00-00	0x00000b41	0xdade	1141	508	0/0/0
4455.6677.0016.00-00	0x00000b3e	0xa832	1011	464	0/0/0
4455.6677.0017.00-00	0x00000b40	0x1563	640	508	0/0/0
4455.6677.0018.00-00	0x00000b3a	0xadee	417	464	0/0/0
4455.6677.0019.00-00	0x00000b3a	0xcff	291	158	0/0/0
4455.6677.001a.00-00	0x00000b3b	0xb131	794	176	0/0/0
4455.6677.001b.00-00	0x00000b3b	0x7062	822	176	0/0/0
4455.6677.001c.00-00	0x00000b3b	0x5876	463	176	0/0/0
4455.6677.001d.00-00	0x00000b3b	0xa610	460	176	0/0/0
4455.6677.001e.00-00	0x00000b3a	0xf5c6	627	176	0/0/0
4455.6677.001f.00-00	0x00000b3b	0x8a19	825	176	0/0/0
4455.6677.0020.00-00	0x00000b3a	0x960a	584	176	0/0/0
4455.6677.0021.00-00	0x00000b3b	0x5b58	1033	176	0/0/0
4455.6677.0022.00-00	0x00000b3a	0x8927	693	176	0/0/0
4455.6677.0023.00-00	0x00000b3b	0x9352	664	158	0/0/0
4455.6677.0024.00-00	0x00000b39	0xc8ec	736	176	0/0/0

*(In TLV)-Leaking Route, *(By LSPID)-Self LSP, +-Self LSP(Extended),
ATT-Attached, P-Partition, OL-Overload



EXAMPLE LSP VERBOSE OF NODE :1 at NODE :3

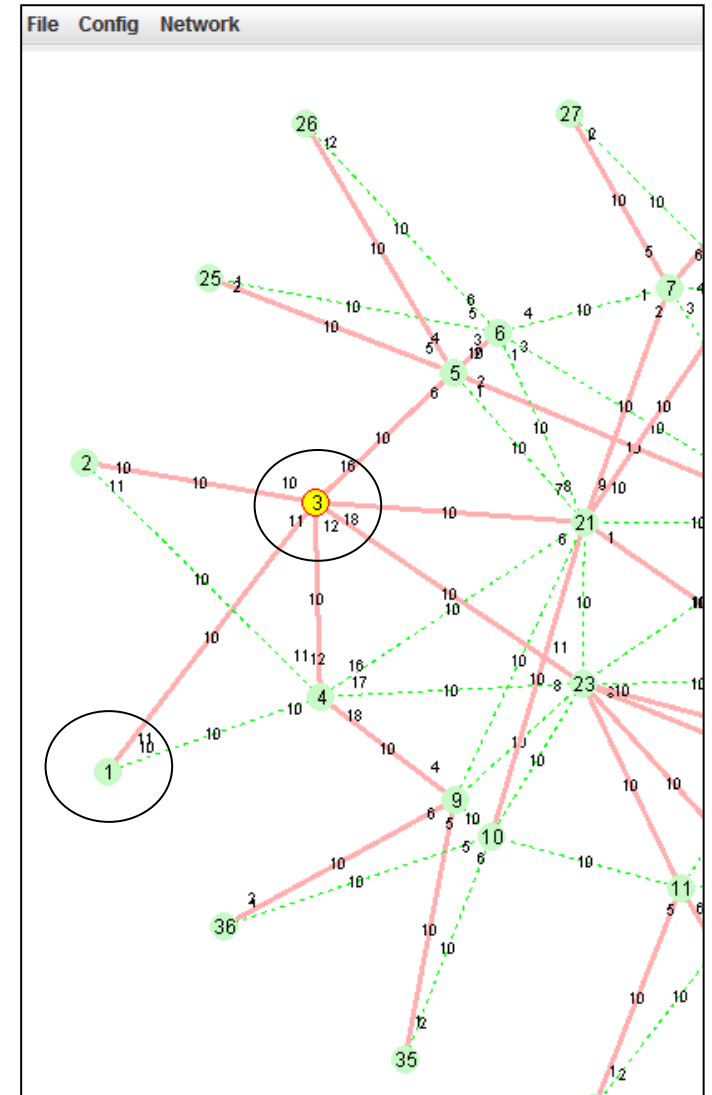
```
<ottawa-9300-3>d isis lsdb 4455.6677.0001.00-00 verbose
```

Database information for ISIS(1)

Level-1 Link State Database

LSPID	Seq Num	Checksum	Holdtime
4455.6677.0001.00-00	0x00000fd3	0x1aeb	1194
SOURCE	4455.6677.0001.00		
NLPID	SPB (0xC1)		
AREA ADDR	22.3344		
+NBR ID	4455.6677.0003.00	COST: 10	
+NBR ID	4455.6677.0004.00	COST: 10	
SPB ECT-ALGORITHM 0	ECT-VID 100		
SPB ECT-ALGORITHM 1	ECT-VID 101		
SPB ECT-ALGORITHM 2	ECT-VID 0		
.....			
SPB ECT-ALGORITHM 15	ECT-VID 0		
SPSID	07-00-01		
SPB BMAC	44-55-66-77-00-01		
ECT-VID	100		
SPB ISID	255T&R		

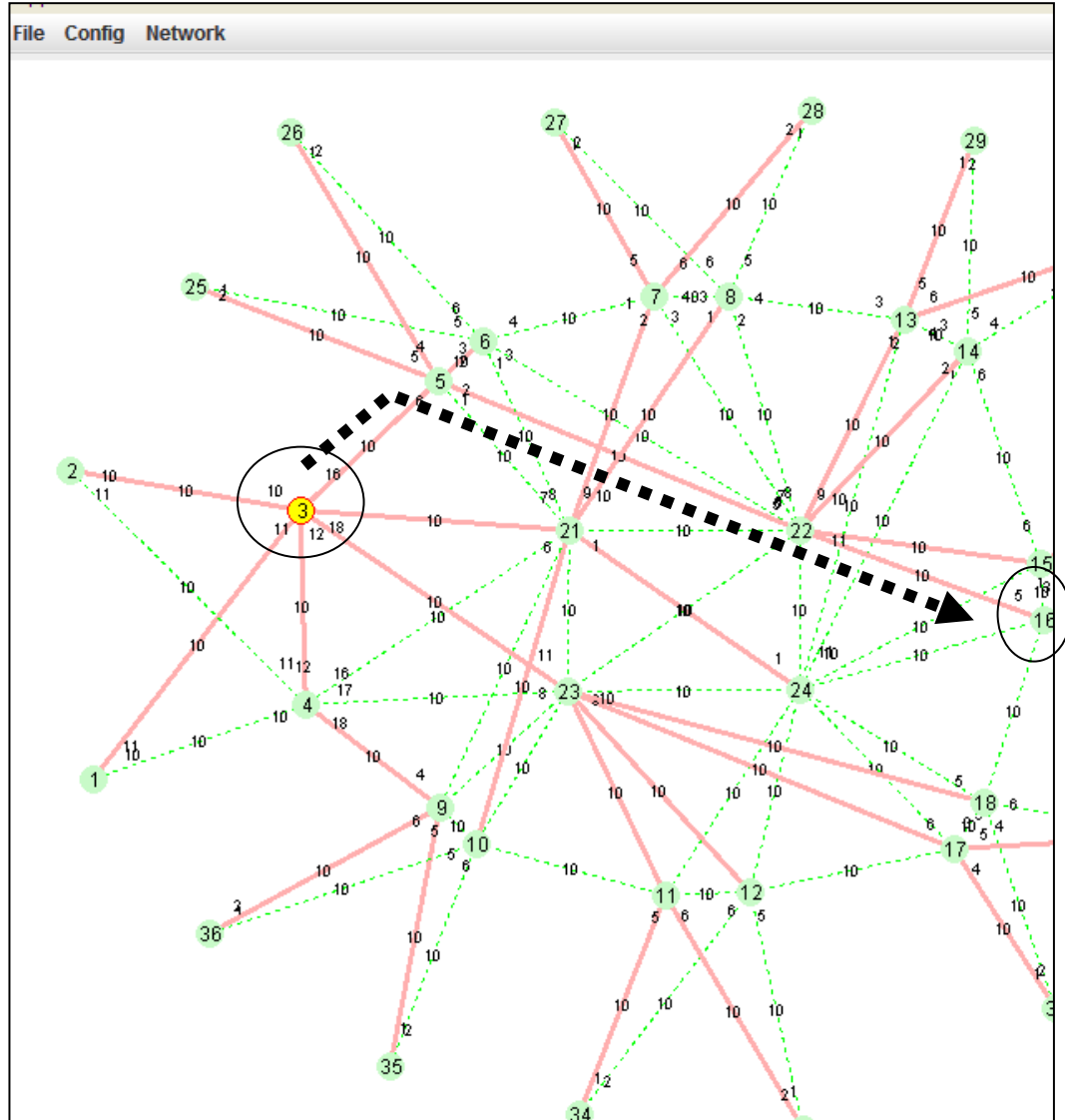
```
<ottawa-9300-3>
```



(first equal cost path)

<ottawa-9300-3>d spb umac

BMAC	BVLAN	IF NAME
4455-6677-0001	100	GE2/0/11
4455-6677-0001	101	GE2/0/11
4455-6677-0004	100	GE2/0/12
4455-6677-0004	101	GE2/0/12
4455-6677-0005	100	GE2/0/16
4455-6677-0005	101	GE2/0/16
4455-6677-0006	100	GE2/0/16
4455-6677-0006	101	GE2/0/17
4455-6677-0007	100	GE2/0/17
4455-6677-0007	101	GE2/0/17
4455-6677-0008	100	GE2/0/17
4455-6677-0008	101	GE2/0/17
4455-6677-0009	100	GE2/0/12
4455-6677-0009	101	GE2/0/18
4455-6677-000a	100	GE2/0/17
4455-6677-000a	101	GE2/0/18
4455-6677-000b	100	GE2/0/18
4455-6677-000b	101	GE2/0/18
4455-6677-000c	100	GE2/0/18
4455-6677-000c	101	GE2/0/18
4455-6677-000d	100	GE2/0/16
4455-6677-000d	101	GE2/0/18
4455-6677-000e	100	GE2/0/16
4455-6677-000e	101	GE2/0/18
4455-6677-000f	100	GE2/0/16
4455-6677-000f	101	GE2/0/18
4455-6677-0010	100	GE2/0/16
4455-6677-0010	101	GE2/0/18
...		
4455-6677-0024	100	GE2/0/12
4455-6677-0024	101	GE2/0/18
Total unicast fib entries is 68		
<ottawa-9300-3>		



EXAMPLE – NODE :3 ROUTE TO :10 (second equal cost path)

<ottawa-9300-3>d spb umac

BMAC	BVLAN	IF NAME
4455-6677-0001	100	GE2/0/11
4455-6677-0001	101	GE2/0/11
4455-6677-0004	100	GE2/0/12
4455-6677-0004	101	GE2/0/12
4455-6677-0005	100	GE2/0/16
4455-6677-0005	101	GE2/0/16
4455-6677-0006	100	GE2/0/16
4455-6677-0006	101	GE2/0/17
4455-6677-0007	100	GE2/0/17
4455-6677-0007	101	GE2/0/17
4455-6677-0008	100	GE2/0/17
4455-6677-0008	101	GE2/0/17
4455-6677-0009	100	GE2/0/12
4455-6677-0009	101	GE2/0/18
4455-6677-000a	100	GE2/0/17
4455-6677-000a	101	GE2/0/18
4455-6677-000b	100	GE2/0/18
4455-6677-000b	101	GE2/0/18
4455-6677-000c	100	GE2/0/18
4455-6677-000c	101	GE2/0/18
4455-6677-000d	100	GE2/0/16
4455-6677-000d	101	GE2/0/18
4455-6677-000e	100	GE2/0/16
4455-6677-000e	101	GE2/0/18
4455-6677-000f	100	GE2/0/16
4455-6677-000f	101	GE2/0/18
4455-6677-0010	100	GE2/0/16
4455-6677-0010	101	GE2/0/18

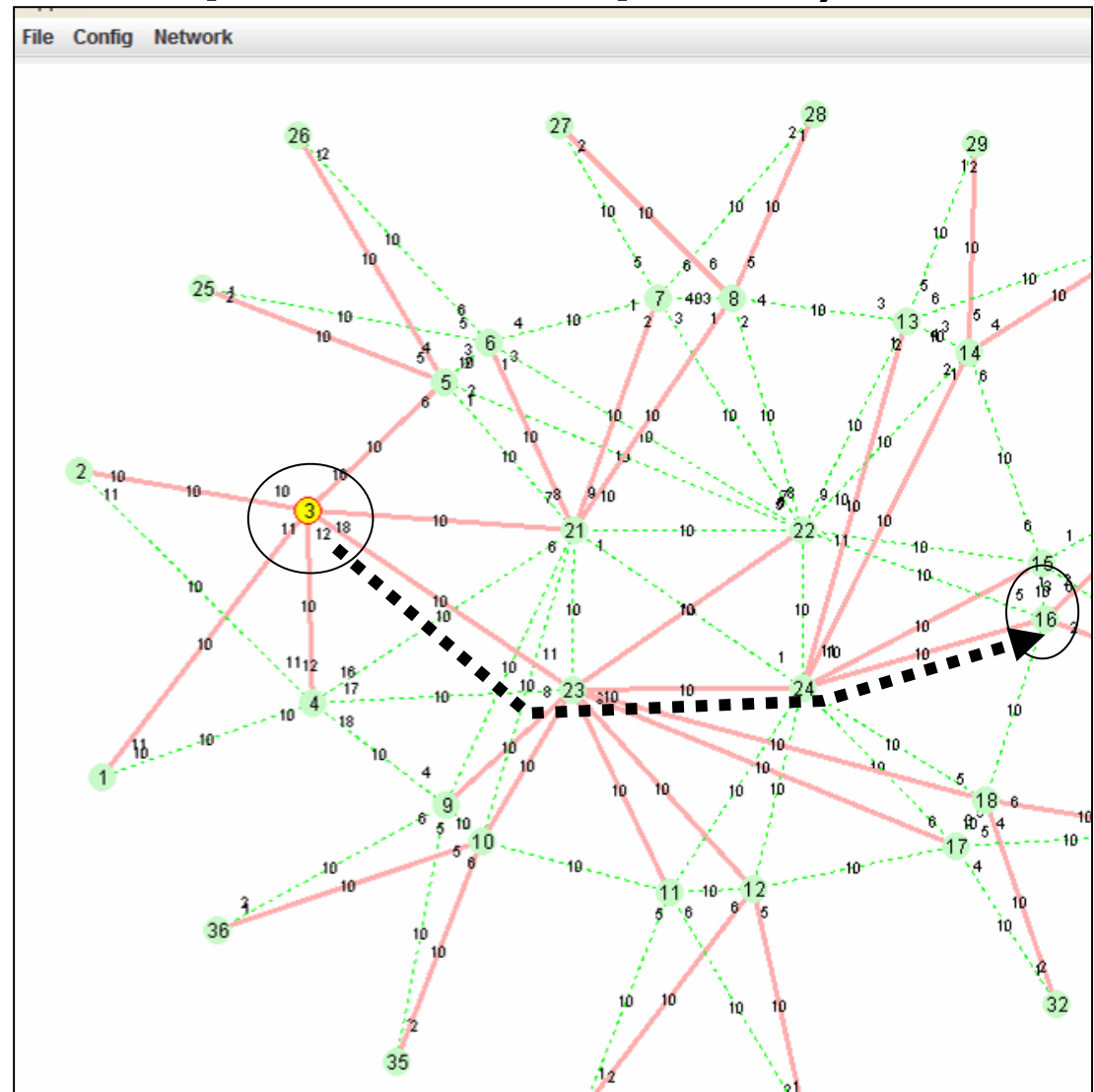
...

4455-6677-0024 100 GE2/0/12

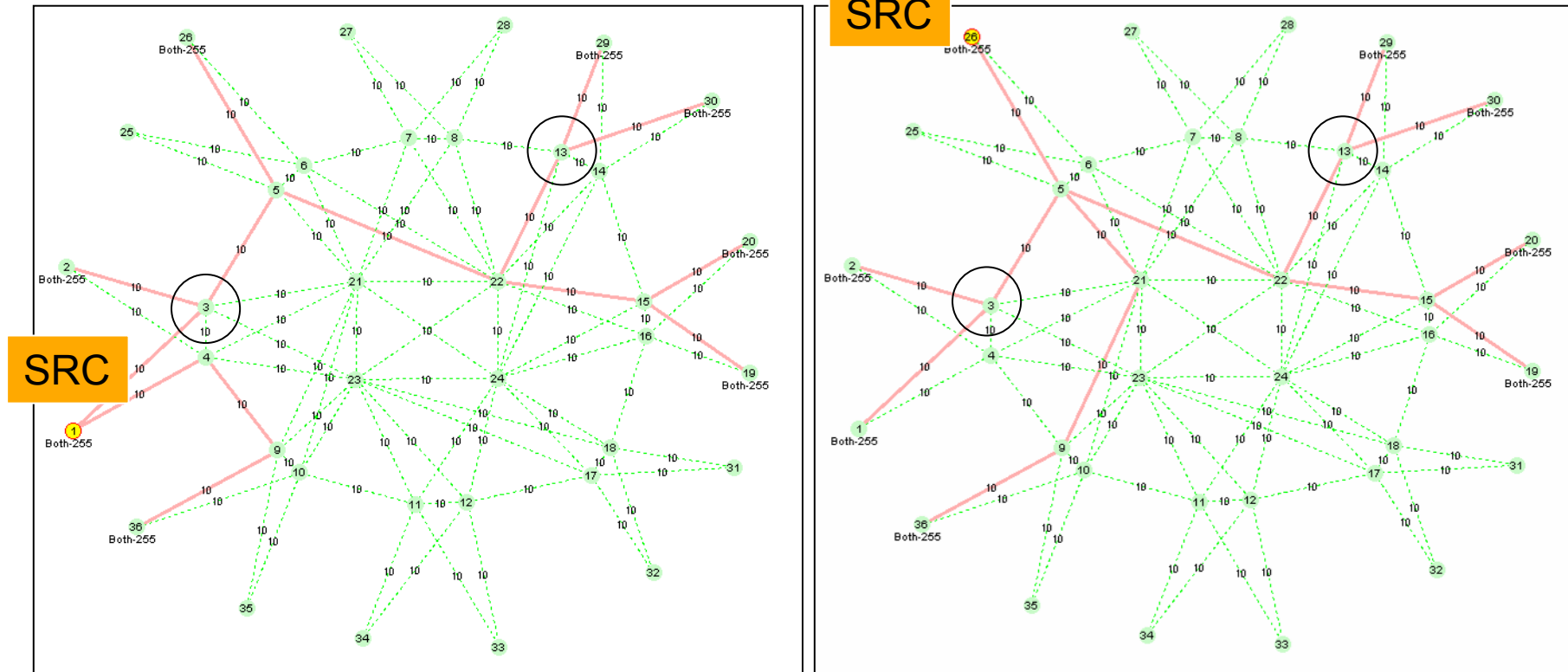
4455-6677-0024 101 GE2/0/18

Total unicast fib entries is 68

<ottawa-9300-3>

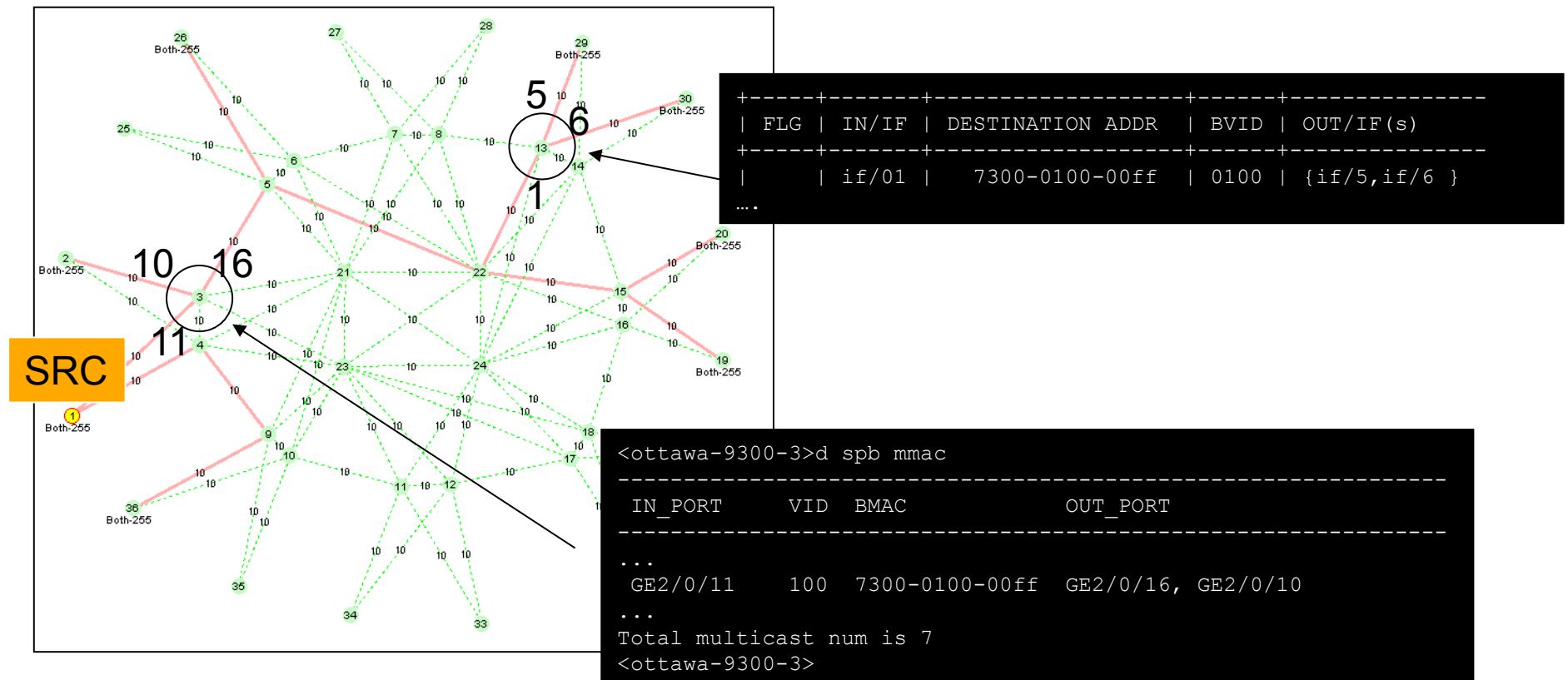


EXAMPLE: E-LAN MCAST ROUTES FROM :1 (left) and :26 (right)



Here are the multicast routes from node 1 for service 255 and also from node 26 for service 255. Note the symmetry in the route between the two multicast trees. The unicast route between :1 and :26 is also along that same path for the chosen B-VID. Since we've asked for transit replication for all members of the E-LAN we install MCAST ...

EXAMPLE: E-LAN MCAST ROUTES FROM :1 (left) and :26 (right)



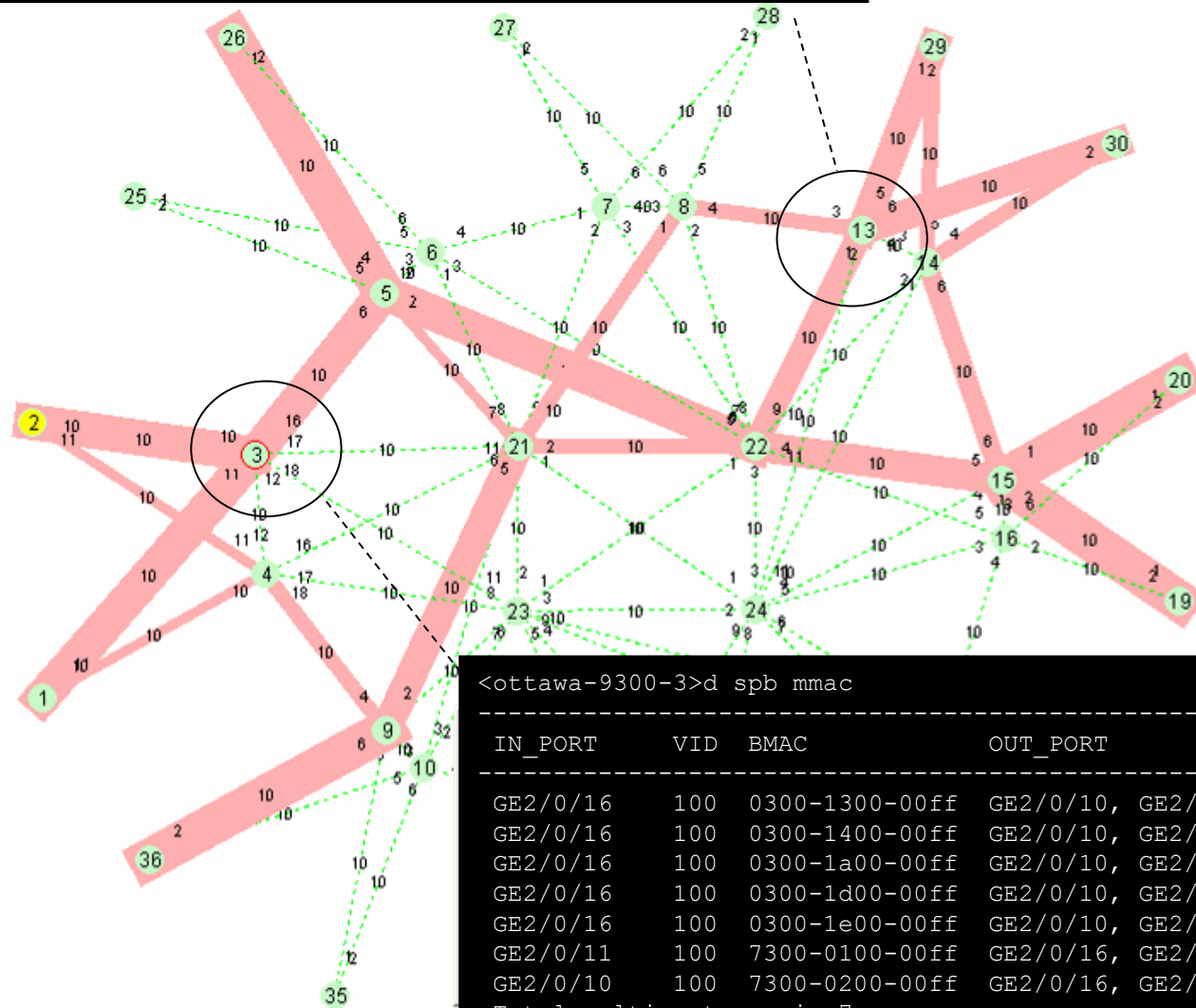
MULTICAST ADDRESS IS: [SOURCE = 07-00-01 | ISID=00-00-ff]

We only get this state if we configure transmit membership in the E-LAN.

Transmit still possible without multicast state but uses serial replication at head end.
Operator chooses trade-off between state/bandwidth usage.

FLG	IN/IF	DESTINATION ADDR	BVID	OUT/IF(s)
	if/01	7300-0100-00ff	0100	{if/5,if/6 }
	if/01	7300-0200-00ff	0100	{if/5,if/6 }
	if/01	0300-1a00-00ff	0100	{if/5,if/6 }
	if/05	0300-1d00-00ff	0100	{if/1,if/3,if/6 }
	if/06	0300-1e00-00ff	0100	{if/1,if/3,if/5 }
	if/03	0300-2400-00ff	0100	{if/5,if/6 }

Here are all mFIBs on nodes :3 and :13 related to this E-LAN.



Information

Node

Instance. ID 3

Node ID 4455.6677.0003

Area ID 22.3344

AreaName AreaTestNet

☒ Level 1 ☒ Level 2

VLAN ID 100

I-SID New

☒ Inst ☐ Real ☐ Virt

Link

Link Name link0-0

Metric 10

Bandwidth 1000

AreaName AreaTestNet

Connection

Address localhost

Port 7001

con1

```
<ottawa-9300-3>d spb mmac
```

IN_PORT	VID	BMAC	OUT_PORT
GE2/0/16	100	0300-1300-00ff	GE2/0/10, GE2/0/11
GE2/0/16	100	0300-1400-00ff	GE2/0/10, GE2/0/11
GE2/0/16	100	0300-1a00-00ff	GE2/0/10, GE2/0/11
GE2/0/16	100	0300-1d00-00ff	GE2/0/10, GE2/0/11
GE2/0/16	100	0300-1e00-00ff	GE2/0/10, GE2/0/11
GE2/0/11	100	7300-0100-00ff	GE2/0/16, GE2/0/10
GE2/0/10	100	7300-0200-00ff	GE2/0/16, GE2/0/11
Total multicast num is 7			

```
<ottawa-9300-3>
```

References

“IEEE 802.1aq” : www.wikipedia.org:

http://en.wikipedia.org/wiki/IEEE_802.1aq

<http://www.ietf.org/internet-drafts/draft-ietf-isis-ieee-aq-00.txt> The IETF IS-IS draft (check for later version 01.. etc).

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“Shortest Path Bridging – Efficient Control of Larger Ethernet Networks” :
upcoming IEEE Communications Magazine – Oct 2010

“Provider Link State Bridging” :

IEEE Communications Magazine V46/N9– Sept 2008

<http://locuhome.com/wp-content/uploads/2009/02/>

[ieeecomcommunicationsmagazinevol46no9sep2008-carrierscaleethernet.pdf](http://locuhome.com/wp-content/uploads/2009/02/ieeecomcommunicationsmagazinevol46no9sep2008-carrierscaleethernet.pdf)

Thank-You

