



An Introduction to
Bidirectional
Forwarding
Detection (BFD)
NANOG 39



Aamer Akhter / aa@cisco.com
ECMD, cisco Systems

Why BFD?

Methods needed to quickly determine forwarding failure

- Not everything is POS/SONET
 - Ethernet needs a solution for failure detection
 - Layer 3 Data Forwarding plane needs a check
 - Checking should not be bound to single hop
 - Fast Hello needed for LDP, OSPF, ISIS, PIM, RSVP, BGP etc to catch same types of issues.
- BFD is a single Layer 3 protocol for detecting forwarding failures
 - Other protocol timers can now be left at defaults

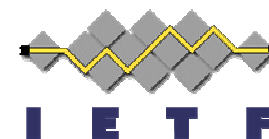
What is BFD?



Bi-directional Forwarding Detection:

- Extremely lightweight hello protocol
IPv4, IPv6, MPLS, P2MP
- 10s of milliseconds (technically, microsecond resolution) forwarding plane failure detection mechanism.
- Single mechanism, common and standardized
Multiple modes: Async (echo/non-echo), Demand
- Independent of Routing Protocols
- Levels of security, to match conditions and needs
- Facilitates close alignment with hardware

IETF Status

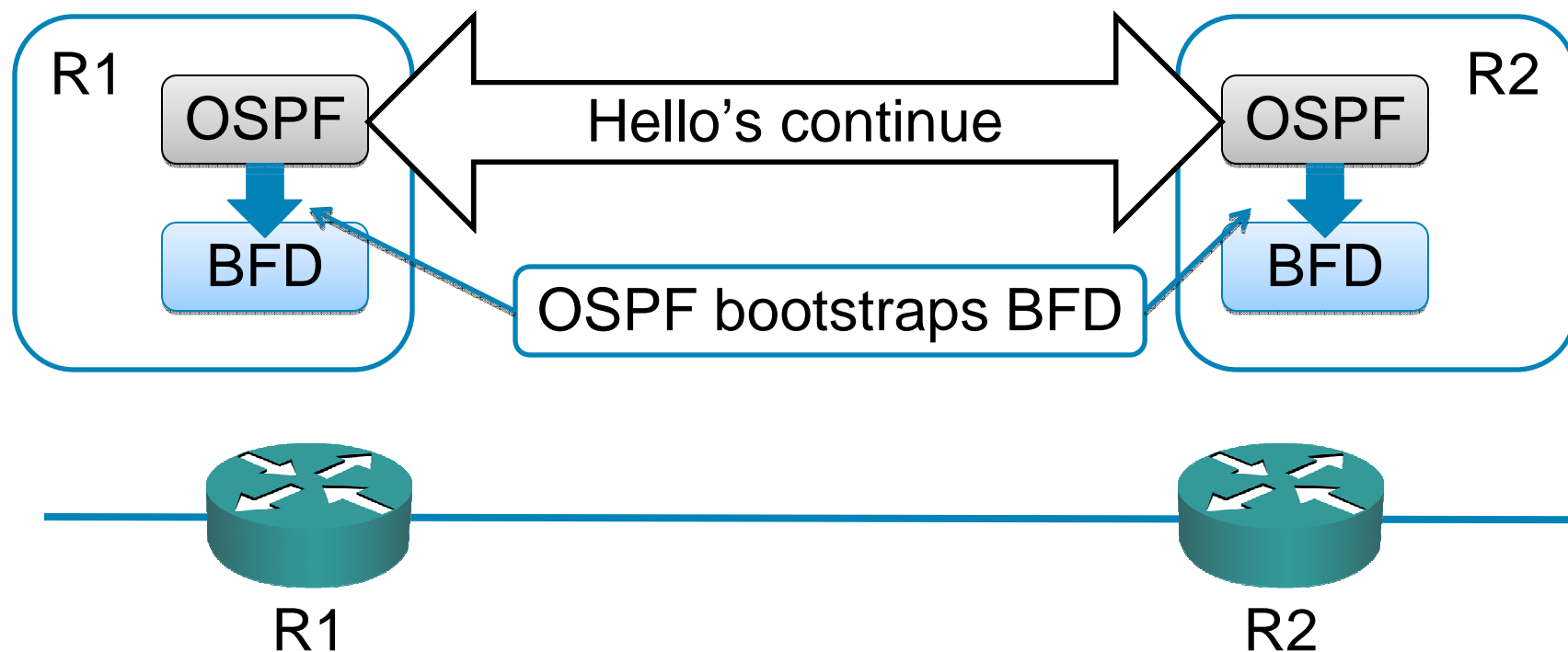


- [BFD workgroup](#)
- Base Draft (Ward and Katz) [draft-ietf-bfd-base-xx](#)
- Generic Application of BFD [draft-ietf-bfd-generic](#)
- BFD for IPv4 and IPv6 (Single Hop) [draft-ietf-bfd-v4v6](#)
- BFD for Multihop Paths [draft-ietf-bfd-multihop](#)
- BFD For MPLS LSPs [draft-ietf-bfd-mpls](#)
- BFD MIB [draft-ietf-bfd-mib](#)
- Additional BFD clients may not require standardization (eg statics client)

Basics of BFD Operation

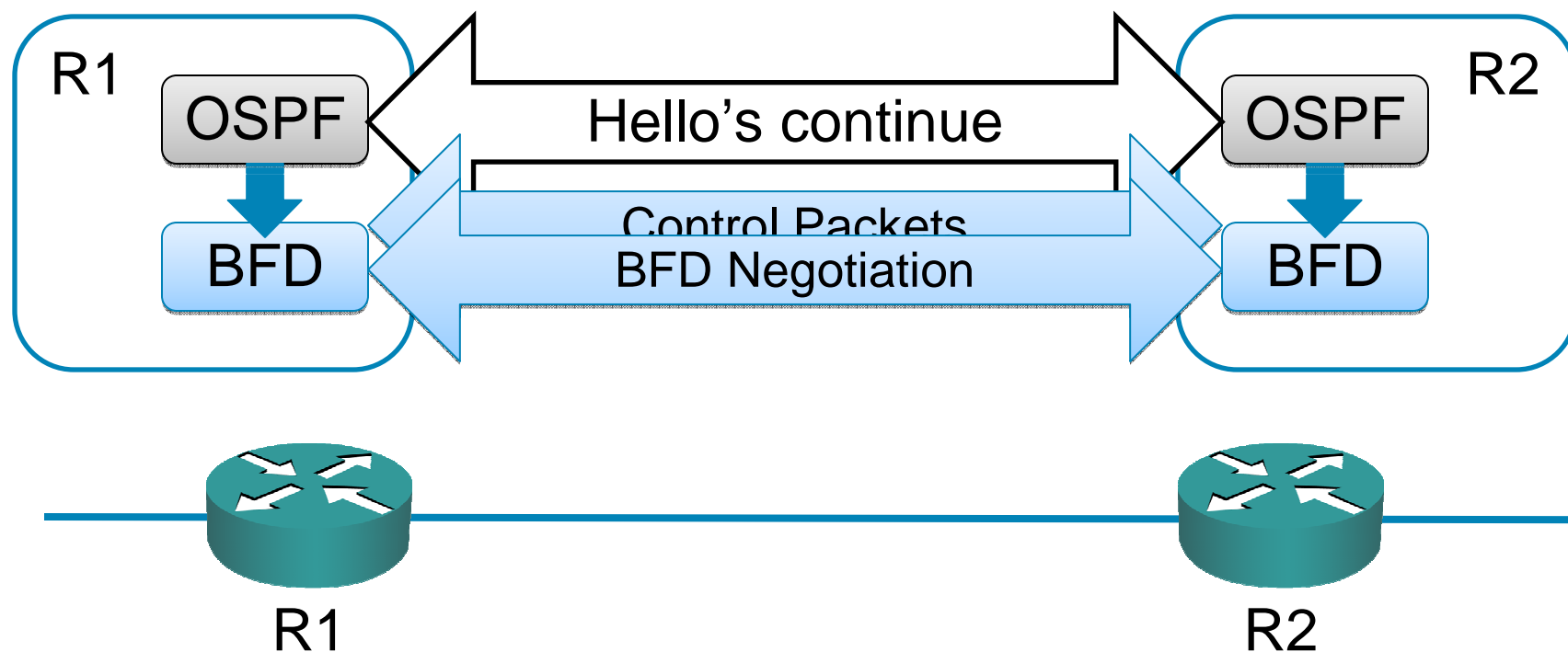
- Routing Protocol (BFD client) bootstraps BFD to create BFD session to a neighbor,
and to receive link status change notification.
- Receive and Transmit intervals are negotiated and configurable
- Two systems agree on method to detect failure
Via sending packets, watching counters etc
- In case of failure, BFD notifies BFD client
- BFD Client independently decides on action (if any)

BFD in Pictures (OSPF Example)



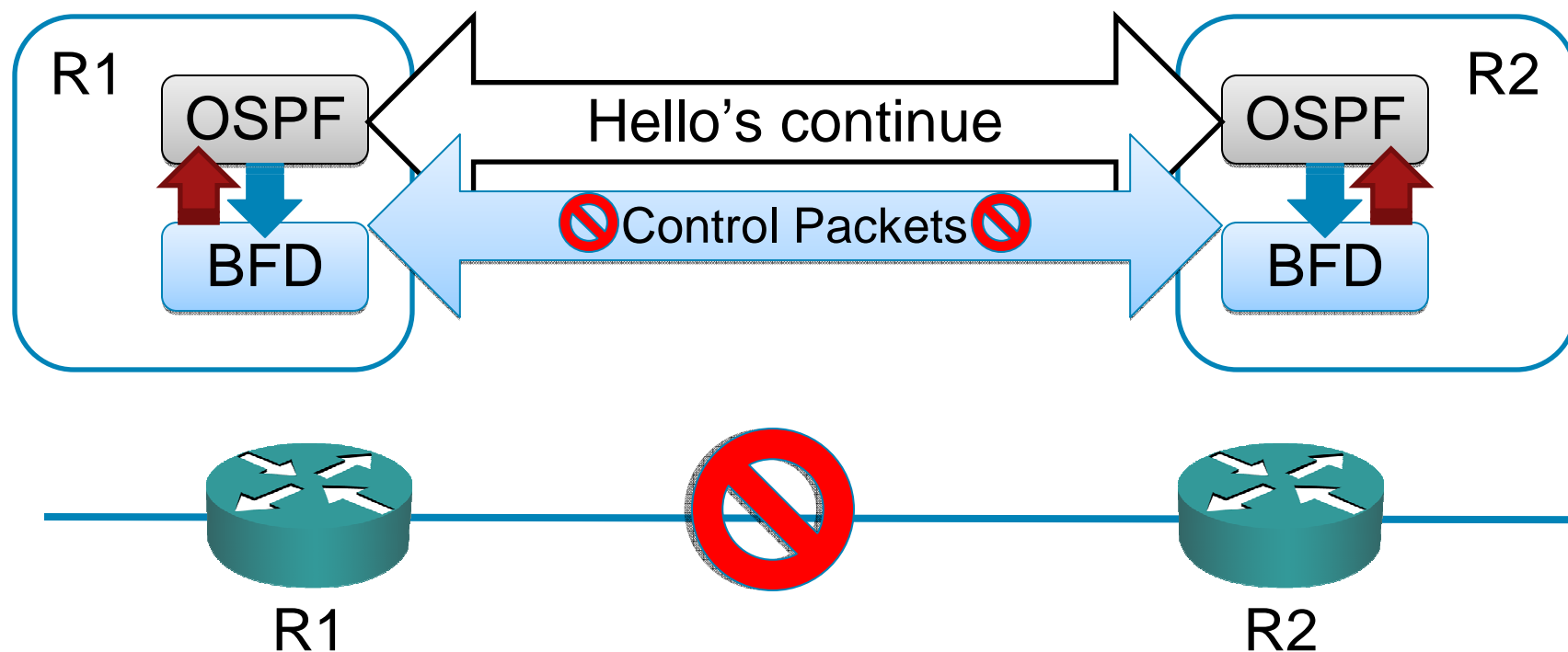
BFD in Pictures (async mode)

- OSPF Hello's at slow rate
- BFD Control packets maintain state and test forwarding plane



BFD in Pictures (async mode)

- BFD notifies OSPF of failure
- OSPF declares neighbor dead
 - Other protocols (ISIS, BGP) may take more granular action



BFD: More Details

Operating Modes

Control & Echo

Timers

Operating Modes

- Asynchronous mode

 - Echo Mode

 - Non-Echo Mode

- Demand Mode

 - Both systems set D bit in control packet

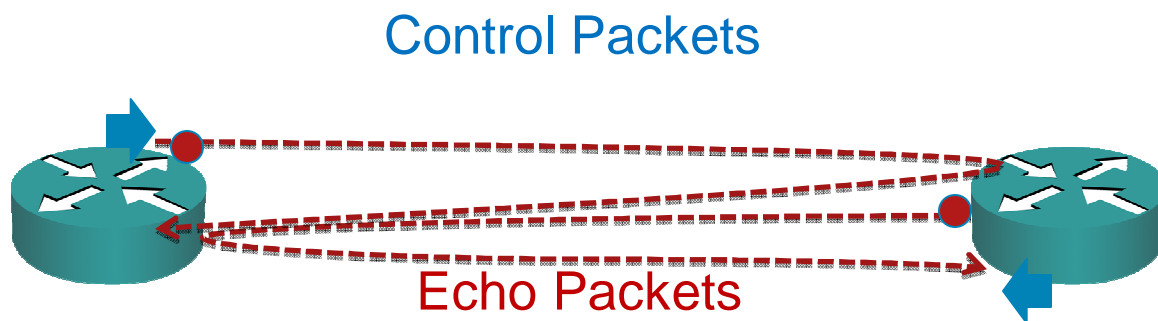
 - Implies that some other method is being used to check forwarding plane

 - Eg, looking at RX/TX counters on interface

 - BFD Polling mechanism will be used to verify liveness as a secondary measure.

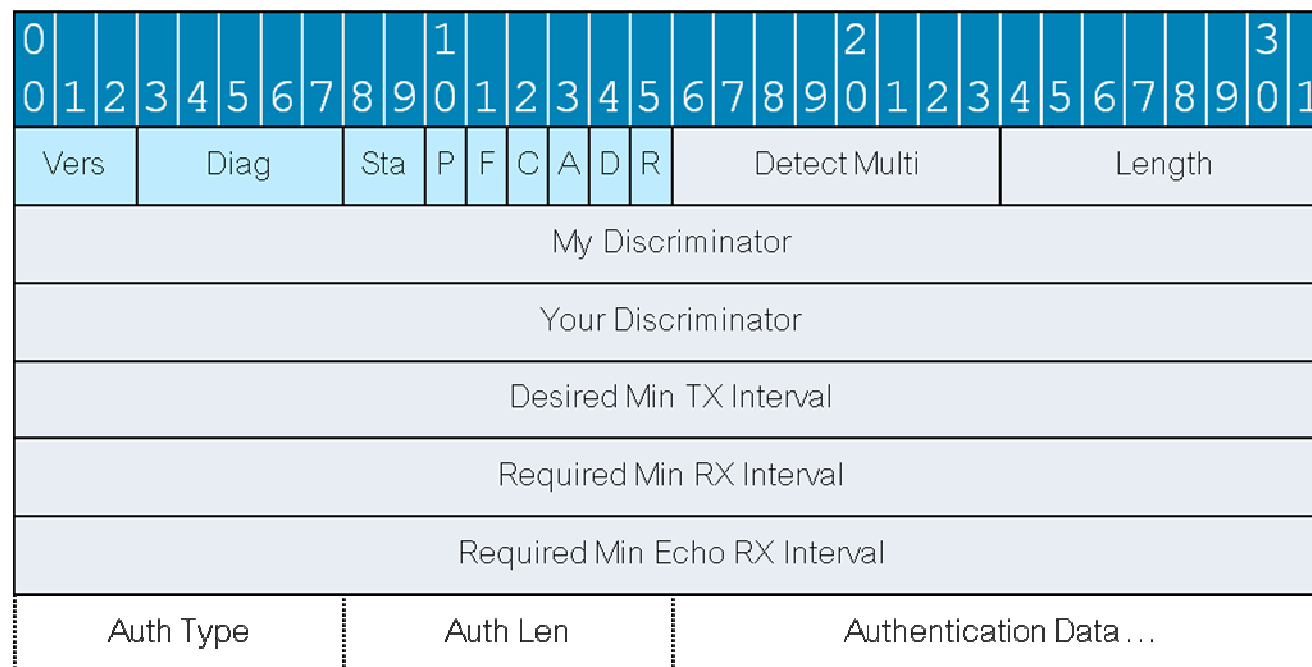
Control Packets and Echo Mode

- If echo function is not negotiated
control packets sent at high rate to achieve Detection Time
- If echo function is negotiated
control packets sent at a slow rate (Negotiated Rate)
self directed echo packets sent at high rate (Min Echo Rx Interval)



The BFD Control Packet

- Control Packets – control session state and parameters
 - Unicast Directed to BFD peer (IP address learnt via BFD client)
 - Are consumed by each BFD recipient
- Single hop: UDP sent to port 3784, source port (49152-65535)



Echo Packets

- Echo Packets

- Self directed

- Low overhead check of fwding plane

- Can be applied Asymmetrically

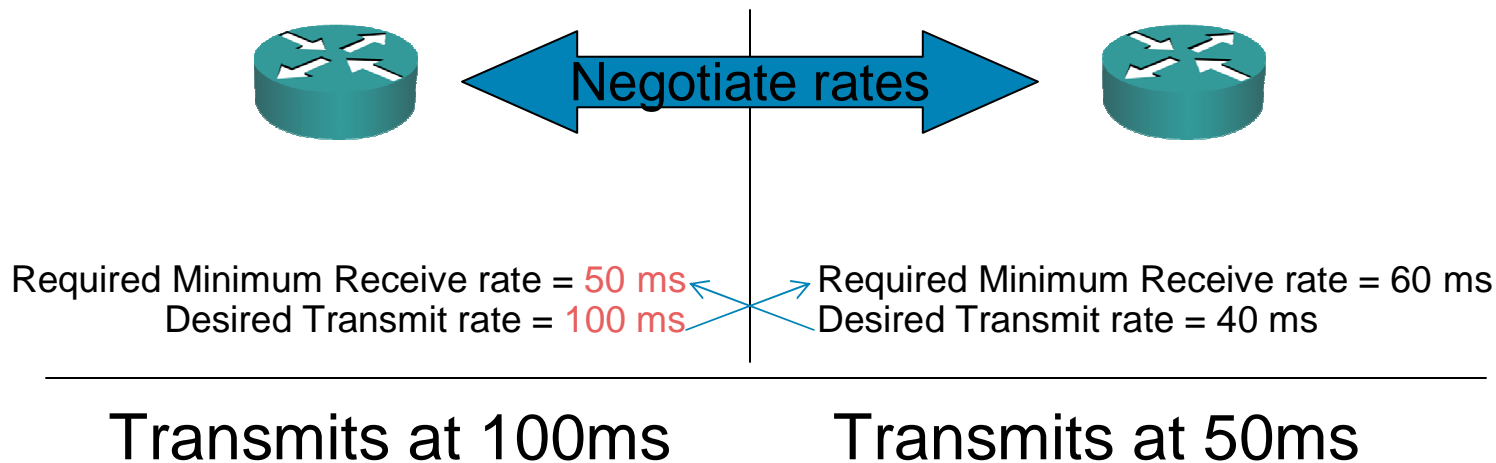
- Format and content of packet determined by sender (implementation dependent)

- Sent to UDP port 3785, for IPv4/v6 single hops

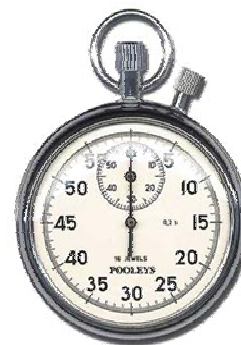
0										1											2										3
0	1	2	3	4	5	6	7	8	9	0	1	2	3	4	5	6	7	8	9	0	1	2	3	4	5	6	7	8	9	0	1
Version																															
My Discriminator																															
Sequence Number																															

Timer negotiation

- Neighbors continuously negotiate their desired transmit and receive rates in terms of microseconds.
- The system reporting the slower rate determines the transmission rate.



Detection Time



- Time to detect failure 😊
Not transmitted on the wire
- Adaptive Timers
Less Restrictive/Tight Time Intervals
- Asynchronous mode (non-echo)
Calculated by $(\text{remote Detect Mult}) * (\text{TX interval})$
Detect Mult is how many sequential packets can be missed before declaring down
- Asynchronous mode (echo)
Detection Time Calculated by $(\text{local Detect Multiplier}) * (\text{local Echo RX interval})$
Loss of 'local Detect Multiplier' # of sequential packets causes failure

Detection Time

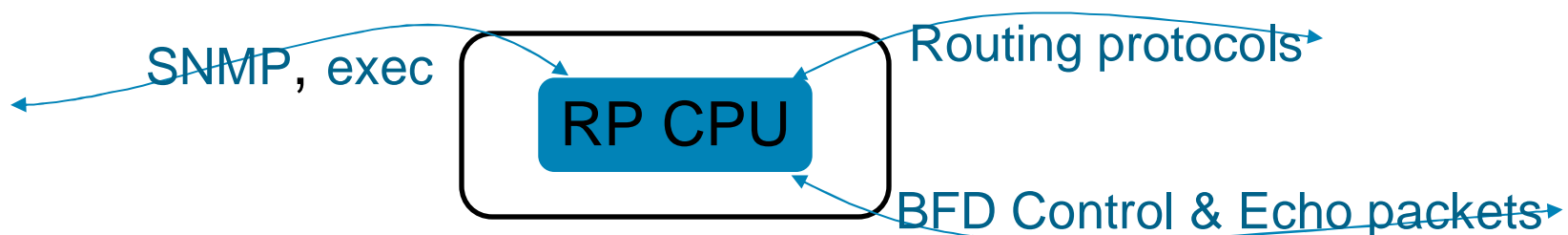
- Demand mode
 - Calculated only during Poll Sequence
 - Calculated by
 $(\text{Detect Multiplier}) * (\text{local TX interval})$

BFD: Implementation Models

- Centralized
- Distributed
- Dedicated Hardware

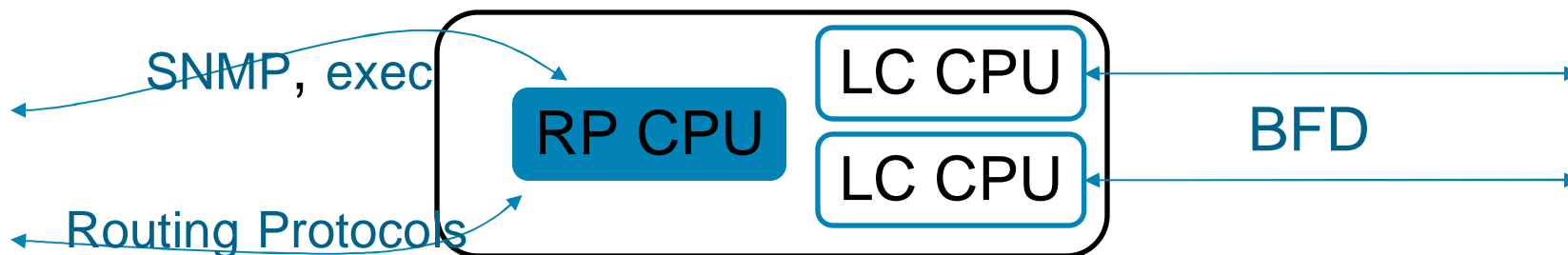
Centralized RP

- Shared CPU for all control plane (may also be shared for data plane)
BGP, OSPF, SNMP, exec
- Contention for CPU cycles ie:
BFD echo generation every 50ms
BGP UPDATE processing
IGP SPF runs
- Issue has always existed, but BFD may aggravate due to low timer values



Distributed CPU

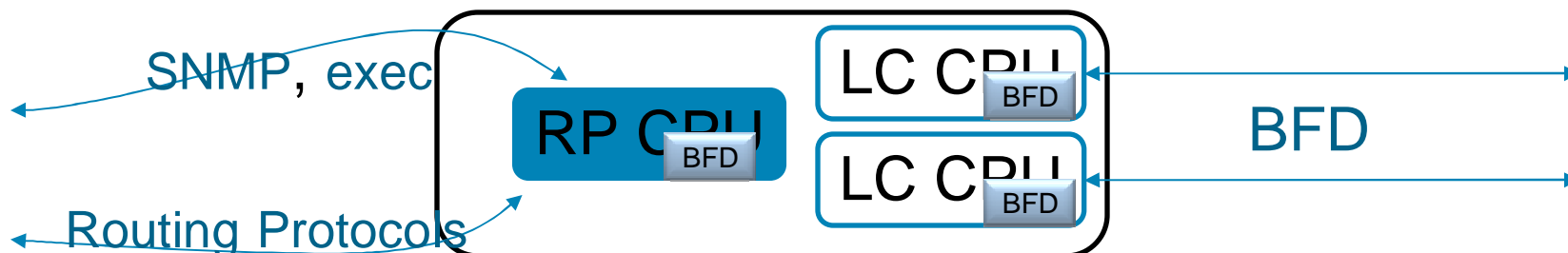
- BFD session maintenance implemented on distributed CPU (eg on line cards)
- LC CPUs generally lightly loaded
- RP can switchover (RP High Availability) w/o affecting BFD



(Semi) Dedicated Hardware

- BFD session maintenance implemented on dedicated or semi dedicated hardware (ie not GPU).
- May still be distributed for additional scalability
- Provides highest level of performance and timing precision

Allows more deterministic BFD performance



Control Plane Independent Bit (Graceful Restart handling)

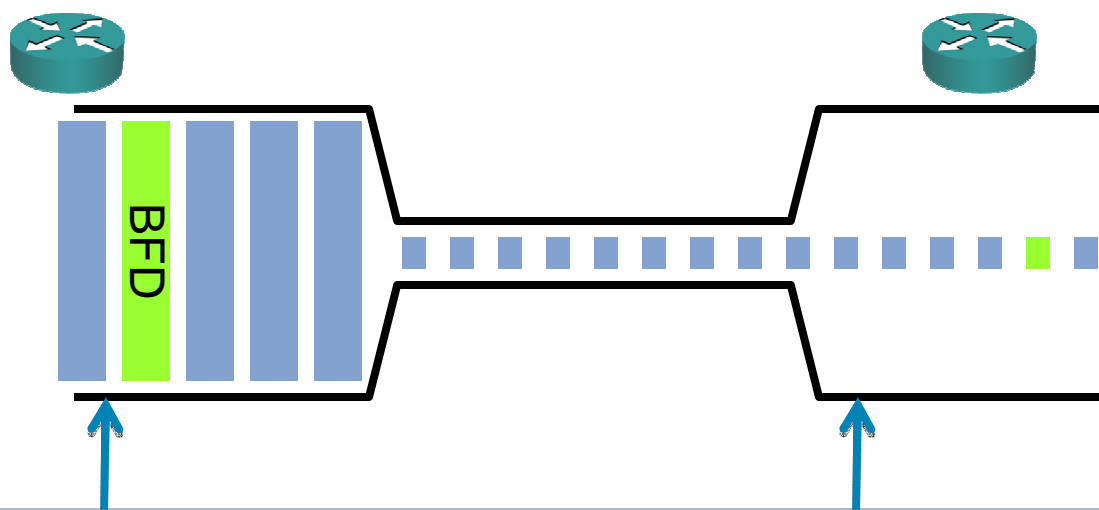
- BFD can inform peer via (c) bit of
Control Plane Independence in BFD implementation
- Makes graceful restart system aware of BFD capability
- Eg: On distributed system, BFD will not be affected by rebooting control plane (ie grace-full restarts, etc)
NSF can be checked by BFD!
 - Eg: In case BFD is hosted by RP and will be affected by GR, (ie C bit clear):
- BFD can temporarily ignore/disable/dilate timers during GR

Queuing (implementation)

- BFD packets can get stuck behind regular packets
- Host IP/UDP protocol stack can be bypassed.

Requires BFD to have direct access to HW queues (no output features)

Might be difficult on distributed architecture where BFD session is on RP & interface on LC ☺



Is internal to router and should give priority to BFD

BFD: Considerations

Queuing

Scalability vs. Performance

Security

Queuing / Latency (not implementation architecture)

- Be realistic: Don't set detection time to 50ms for transatlantic links!
- Multi-hop (and echo mode) BFD still requires queue management on transit routers
 - Marked with Precedence 6
 - Prec 6 may need QoS policy configured on transit routers
 - Depending on link speed, BFD may need to be in LLQ
 - Verify that the correct data path is validated (ie DS-TE, DSCP PBR, MTR can put packets on different paths)

Scalability and Low Timer Values

- Possible to move resources around
 - BFD to LC, dedicated BFD CPU, relaxing timers etc
- Practical and finite limits on resources
- Lower BFD timer values and strict detection timing will decrease scalability (# of sessions) for that resource
 - Adaptive timers take care of situations w/o strict detection timing needs
 - Implementation may protect you by disallowing enabling of sessions that can't be maintained.



Scalability and Low Timer Values

- Testing under your ‘real-world’ conditions is essential
- Aggregate BFD pps per LC/CPU creates a composite between sessions and timer values
- Spreading BFD amongst multiple resources (line cards)
- BFD does not invalidate operator experience with low BGP KA or IGP hello times— just changes the game



False Positives / Oscillations

- Use adaptive timers / echo mode / demand mode
- Why they can happen:
- Generally, implementation issues
- Conditions change in network, but nothing really wrong
 - When testing, account for **stress** conditions, not best conditions
 - BGP updates
 - IGP recalculations
 - SNMP polls
 - Traffic bursts
- Stress can be transient
 - or related to new services that cross perf. threshold

Tuning Timers

- BFD allows timer renegotiation during session
- Adaptive Timers (all modes have this)
 - Less restrictive, can automatically adapt for slow local/remote system
 - Puts actual fault detection time into grey area
- Examples:
 - Control Mode: (slow remote) avg of last few rxvd pkts is used
 - Echo Mode: (slow self)
 - RX Count is done on packets actually transmitted, not what *should* have been transmitted
 - Detection Time = losing 'detect multi' # of packets (regardless of when they were sent out)

Tuning Timers



- Need to monitor state of router
 - SNMP traps on consistently slow RX/TX, etc
- Are BFD packets being sent/rcvd at the configured value?
 - Increase or decrease interval's accordingly (if needed)

Security



- TTL checking
 - single hop
 - 255 on sending, check for 255 on receipt
- Authentication (more work for router)
 - Multi-hop applicability
 - Key ID's allow for rollover
 - Simple Password
 - MD5
 - SHA1
 - Meticulous Keyed

BFD Summary

- Solves a real issue with fast forwarding plane checks
- Extremely lightweight hello protocol
IPv4, IPv6, MPLS, P2MP
- 10s of milliseconds (technically, microsecond resolution) forwarding plane failure detection mechanism.
- Single mechanism, common and standardized
Multiple modes: Async (echo/non-echo), Demand
- Independent of Routing Protocols
- Facilitates close alignment with hardware

Q and A

